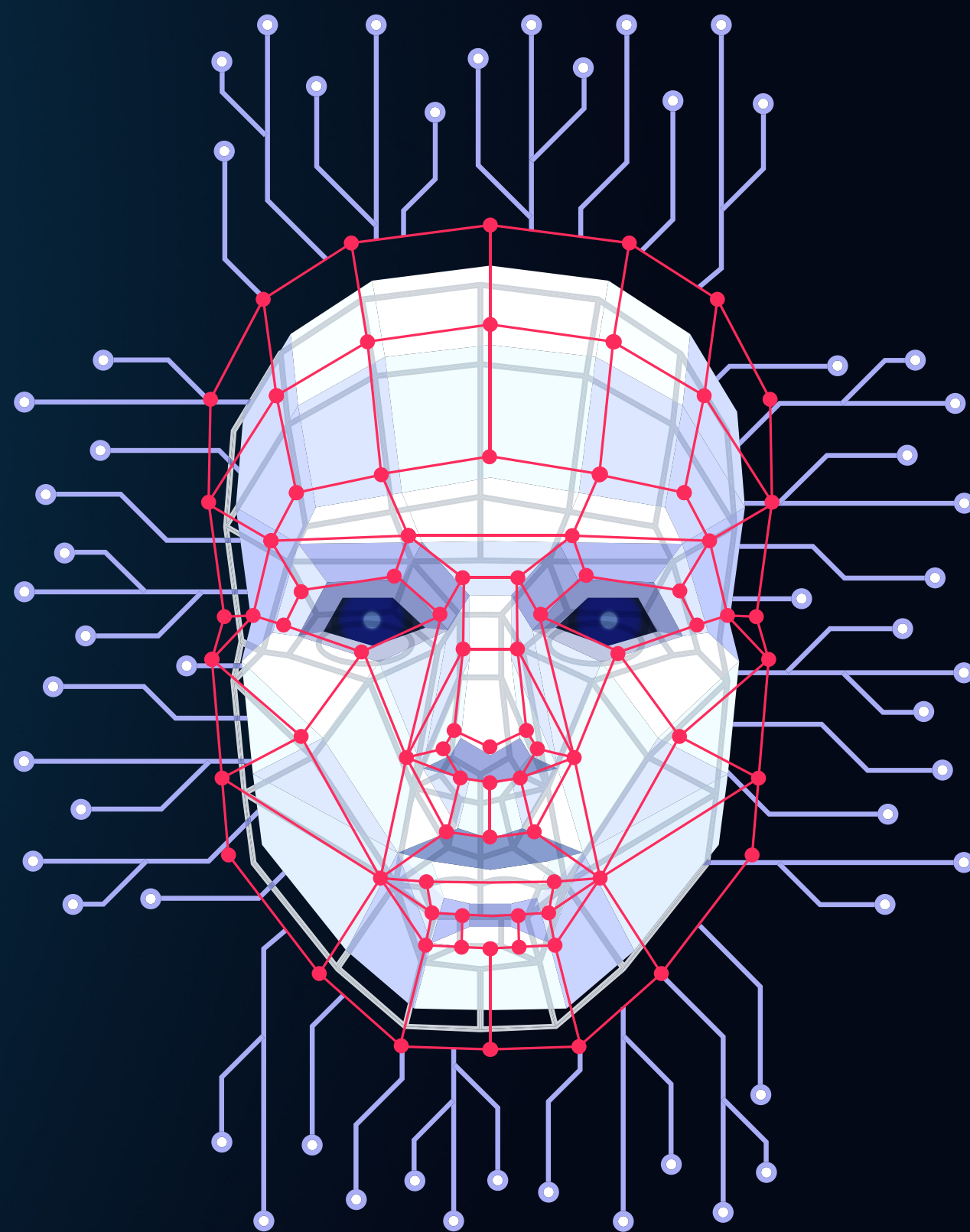




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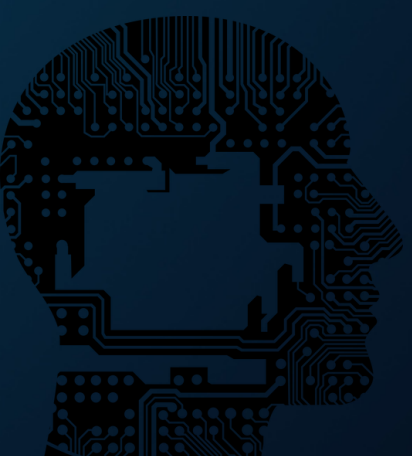
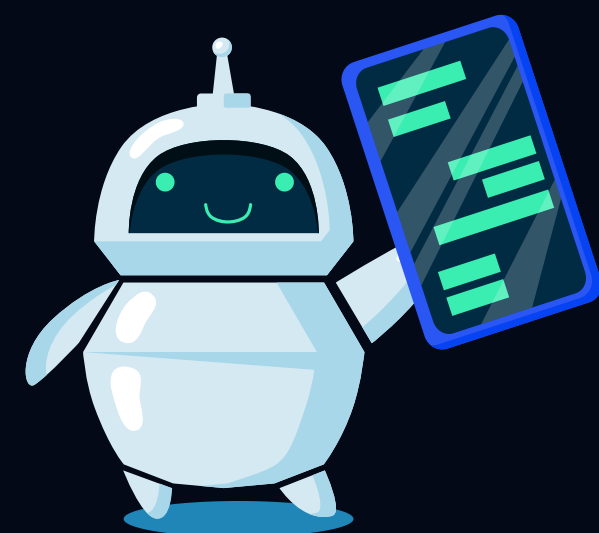
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INTELLIGENCE, RESEARCH & TECHNOLOGY**

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An Edited Book



Chief Editor

Dr. Vipin Mittal

Director, IJRTS Takshila Foundation, Jind, Haryana, India

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Intelligence, Research & Technology**
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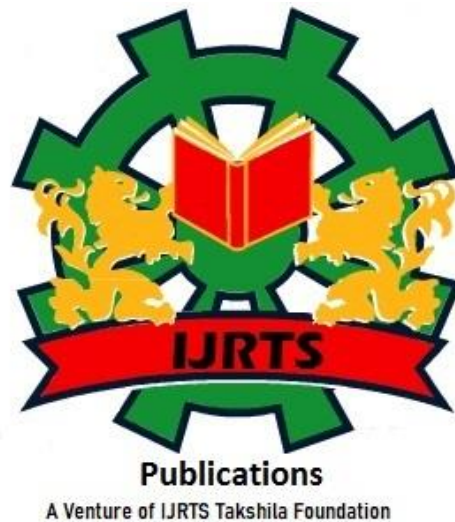
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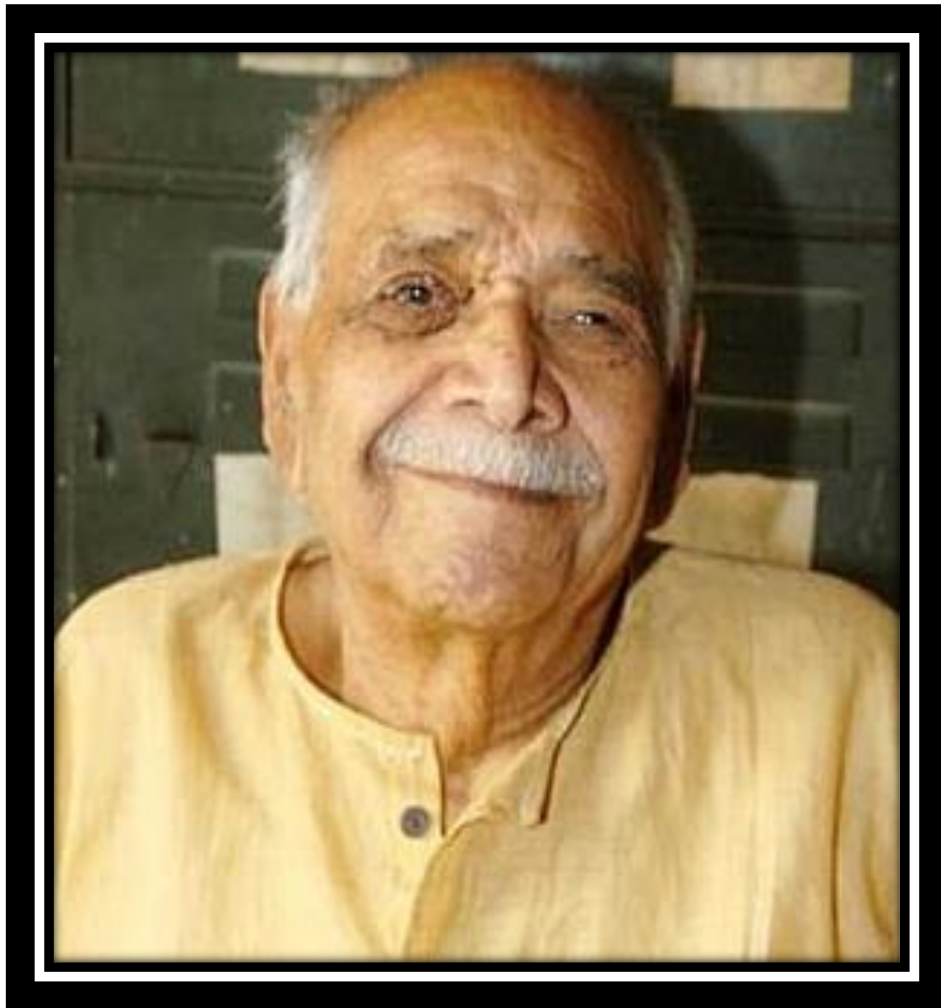
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Ram Sharan Sharma was an Indian Marxist historian and Indologist who specialised in the history of Ancient and early Medieval India. He taught at Patna University and Delhi University (1973–85) and was visiting faculty at University of Toronto (1965–1966). He also was a senior fellow at the School of Oriental and African Studies, University of London. He was a University Grants Commission National Fellow (1958–81) and the president of Indian History Congress in 1975. It was during his tenure as the dean of Delhi University's History Department that major expansion of the department took place in the 1970s. The creation of most of the positions in the department were the results of his efforts. He was the founding Chairman of the Indian Council of Historical Research (ICHR) and a historian of international repute. His pioneering effort resolved the border dispute forever as recorded by Sachchinand Sinha in a letter to Rajendra Prasad.

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Mathematics behind the Evolution of Calendar

Dr. Govil Jindal

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Abstract

This chapter deals the evolution of our calendar system. We passed to many great civilisations and each civilisation had their own system of calendar. We encountered with many system of calendar such as Babylonians, Hindu, Chinese, Egyptians, Greek, Hebrew, Roman, Julian, Gregorian and so on. Every system used three types of calendar, Lunar, Solar and Lunisolar calendar for measuring the times and keeping the events in records. The Romans started with a 10-month calendar in 738 B.C., borrowing from the Greeks. The months in the original Roman calendar were Martius, Aprilis, Maius, Junius, Quintilis, Sextilis, September, October, November and December. Julius Caesar changed the calendar. Ignoring the moon but keeping the existing 12 month's names, the year was divided into 12 months having 30 or 31 days, except Februarius at the end with 29 days. In the Julian Calendar a year is leap if it is divisible by 4. In the Gregorian Calendar a year is leap if either (i) it is divisible by 4 but not by 100 or (ii) it is divisible by 400.

INTRODUCTION

The purpose of the calendar is to reckon past or future time, to show how many days until a certain event takes place—the harvest or a religious festival—or how long since something important happened. The earliest calendars must have been strongly influenced by the geographical location of the people who made them. In colder countries, the concept of the year was determined by the seasons, specifically by the end of winter. But in warmer countries, where the seasons are less pronounced, the Moon became the basic unit for time reckoning; an old Jewish book says that “the Moon was created for the counting of the days.” Most of the oldest calendars were lunar calendars, based on the time interval from one new moon to the next—a so-called lunation. But even in a warm climate there are annual events that pay no attention to the phases of the Moon. In some areas it was a rainy season; in Egypt

it was the annual flooding of the Nile River. The calendar had to account for these yearly events as well.

SOME INTERESTING FACTS NEED TO INVESTIGATE:

1. February me 28 din kyo hote h?
2. October is Eighth month?
3. Ramzan Id 11 din pehle kyu ati h?
4. July or August me 31 din kyo?
5. October 4 – Saint Teresa of Ávila dies. She is buried the next day, October 15?
6. Six and a half million Britons went to bed on September 2, 1752, and woke up on September 14. The reason?

TYPES OF CALENDAR

1. SOLAR CALENDAR

A *solar calendar* is designed to approximate the tropical year using days. In order to synchronize with the tropical year (solar year: 365 d, 5 h, 48 min, 46 sec (365.242216d; $0.007784d \cdot 1300y = 10.1192d$) and hence the seasons, days are sometimes added, forming leap years, to increase the average length of the calendar year. A solar calendar year can be divided into months but these months ignore the Moon. The Gregorian calendar is a solar calendar with a common year having 365 days and a leap year having 366 days. Every fourth year is a leap year unless it is a century year not divisible by 400.

2. LUNAR CALENDAR

An old Jewish book says that “the Moon was created for the counting of the days.”

A *lunar calendar* consists of a number of lunar months with each month covering the period between two successive new moons or full moons. We say that the lunar month follows, or depends on, the lunar cycle. Each calendar or lunar year has 12 lunar months. Each month has an average length of about 29.5 days. This amounts to about $12 \times 29.5 = 354$ days a year, around 11 days shorter than the tropical year. Hence a lunar calendar ignores the tropical year and does not keep in line with the seasons. The Muslim calendar is a lunar calendar. We can see that the Ramzan Id/Eid-ul-Fitar festival always falls about 11 days earlier than a year ago in the Gregorian calendar.

RamzanId/Eid-ul-Fitar Occurance in last 05 years

- 2016: 06 July
- 2017: 26 June
- 2018: 16 June

- 2019: 05 June
- 2020: 25 May

3. LUNISOLAR CALENDAR

A *lunisolar calendar* is designed to keep in phase with the tropical year using lunar months. A whole lunar month is occasionally added at every few years interval to help the calendar keep up with the tropical year. This additional month is known as the *leap month* or the *intercalary month*. The Chinese calendar is a lunisolar calendar, consisting of 12 lunar months, each beginning at new moon. A normal calendar year has 12 months and a 13th month is added according to certain rules to synchronize with the tropical year.

KALI PUJA / DEEPAVALI / DIWALI DATE LIST

from
2000 to 2019 :

Year :	Date :	Weekday :	Tithi :
2000	October 27, 2000	Friday	(Amavasya - Krishna Paksha)
2001	November 15, 2001	Thursday	(Amavasya - Krishna Paksha)
2002	November 04, 2002	Monday	(Amavasya - Krishna Paksha)
2003	October 25, 2003	Saturday	(Amavasya - Krishna Paksha)
2004	November 12, 2004	Friday	(Amavasya - Krishna Paksha)
2005	November 02, 2005	Wednesday	(Amavasya - Krishna Paksha)
2006	October 22, 2006	Sunday	(Amavasya - Krishna Paksha)
2007	November 09, 2007	Friday	(Amavasya - Krishna Paksha)
2008	October 28, 2008	Tuesday	(Amavasya - Krishna Paksha)
2009	October 18, 2009	Sunday	(Amavasya - Krishna Paksha)
2010	November 06, 2010	Saturday	(Amavasya - Krishna Paksha)
2011	October 26, 2011	Wednesday	(Amavasya - Krishna Paksha)
2012	November 13, 2012	Tuesday	(Chaturdashi - Krishna Paksha)
2013	November 03, 2013	Sunday	(Amavasya - Krishna Paksha)
2014	October 23, 2014	Thursday	(Amavasya - Krishna Paksha)
2015	November 11, 2015	Wednesday	(Amavasya - Krishna Paksha)
2016	October 30, 2016	Sunday	(Amavasya - Krishna Paksha)
2017	October 19, 2017	Thursday	(Amavasya - Krishna Paksha)
2018	November 07, 2018	Wednesday	(Amavasya - Krishna Paksha)
2019	October 28, 2019	Monday	(Amavasya - Krishna Paksha)

THE JOURNEY OF CALENDAR: FROM ROMAN CALENDAR TO GREGORIAN CALENDAR

1. ROMAN CALENDAR

When Rome emerged as a world power, the difficulties of making a calendar were well known. The Romans started with a 10-month calendar in 738 B.C., borrowing from the Greeks. The months in the original Roman calendar were Martius, Aprilis, Maius, Junius, Quintilis, Sextilis, September, October, November and December. The names Quintilis through December come from the Roman names for five, six, seven, eight, nine and 10. This calendar left 60 or so days unaccounted for. The months Januarius and Februarius were later added to the end of the year to account for the 60 spare days. However, four months of 31 days, seven months of 29 days, and one month of 28 days added up to only 355 days. Therefore the Romans invented an extra month called Mercedonius of 22 or 23 days. It was added every second year.

Before 47 B.C.			45 B.C.			08 A.D.			1583 onwards		
Num ber	Month	Days	Num ber	Month	Days	Num ber	Month	Days	Num ber	Month	Days
1	Martius	29 or 30	1	January	31	1	January	31	1	January	31
2	Aprilus	29 or 30	2	February	29	2	February	28	2	February	28
3	Maius	29 or 30	3	March	31	3	March	31	3	March	31
4	Junius	29 or 30	4	April	30	4	April	30	4	April	30
5	Quintilis	29 or 30	5	May	31	5	May	31	5	May	31
6	Sextilis	29 or 30	6	June	30	6	June	30	6	June	30
7	September	29 or 30	7	Quintilis	31	7	July	31	7	July	31
8	October	29 or 30	8	Sextilis	30	8	August	31	8	August	31
9	November	29 or 30	9	September	31	9	September	30	9	September	30
10	December	29 or 30	10	October	30	10	October	31	10	October	31
11	Januarius	29 or 30	11	November	31	11	November	30	11	November	30
12	Februarius	29 or 30	12	December	30	12	December	31	12	December	31

2. JULIAN CALENDAR

Following his conquest of Egypt in 48 B.C. Roman Emperor, Julius Caesar consulted the Alexandrian astronomer Sosigenes about calendar reform. Caesar adopted the calendar

identical to the Alexandrian Aristarchus' calendar of 239 B.C., that consisted of a solar year of 12 months and of 365 days with an extra day every fourth year. This was a truly solar calendar and was later called, "Julian Calendar" named after him. Month lengths were extended to 30 or 31 for different months to bring the calendar's total to 365 with an addition of an extra day every fourth (leap) year to account for the true length of the solar year more close to 365.25 days.

By the year 46 B.C. it was noticed that the spring equinox had shifted by about 2 months coming in May. Caesar wanted to bring spring equinox into correct position of Martius (March) 21. The year 46 B.C. is regarded as the "year of confusion" by modern authors, because in that year two intercalations were done; one to correct for spring equinox and the other to change lunisolar calendar to purely solar calendar resulting in the length of that year to 445 days. Then the solar year (with the value of 365 days and 6 hours) was made the basis of the calendar.

The months were 30 or 31 days in length, except Februarius at the end with 29 days. Every fourth year, Februarius gained an extra day. Later, he decided to make Januarius the first month instead of Martius, making Februarius the second month, which explains why leap day is at such a funny point in the year.

Caesar was assassinated in 44 B.C. and the month Quintilis was later renamed as Julius (July) to honor him. Soon after Caesar's assassination, his nephew and adopted son Octavian became the emperor that we refer to as "Augustus," an honorary title that had been bestowed upon him. Similarly, Sextilis was renamed to honor Augustus, hence August. But the way the story goes, Sextilis in Octavian's time had only 30 days. Octavian's advisors suggested that the month named after him should not be deficient compared to Julius's month of 31 days. Octavian (Augustus) thus looked around for a month from which he could remove a day. Romans had religious festivals that were held on specific days of each month; it would have been difficult for him to delete any day that remembered a revered event, so he took a day from Februarius which was already an odd month and added it to Sextilis, which was then renamed Augustus (August). So February now had only 28 days in common years but had 29 in leap years.

3. GREGORIAN CALENDAR

In 1545, the Council of Trent authorized Pope Paul III to reform the calendar once more. However, the Julian year still differs from the true year and by 1582, the error had accumulated to 10 days. So, 10 days were dropped from the year 1582 so that October 4,

1582 was followed by October 15, 1582. Most of the mathematical and astronomical work was done by Father Christopher Clavius, S.J. The immediate correction, advised by Father Clavius and ordered by Pope Gregory XIII, was that Thursday, Oct. 4, 1582, was to be the last day of the Julian calendar. The next day would be Friday, Oct. 15. Pope Gregory XIII implements the Gregorian calendar. The Gregorian reform was not adopted throughout the West immediately. Most Catholic countries quickly changed to the pope's new calendar in 1582. . So on October 15, 1582, Pope Gregory XIII released a papal bull—a declaration from the leader of the Catholic church—decreeing that those under the dominionship of his church would have to skip some days. Spain, large parts of Italy (which was not yet unified), the Netherlands, France, Portugal, Luxembourg, and Poland and Lithuania (who were at the time tied under a commonwealth) all adopted Gregory's bull that year. But Europe's Protestant princes chose to ignore the papal bull and continued with the Julian calendar. It was not until 1700 that the Protestant rulers of Germany and the Netherlands changed to the new calendar.

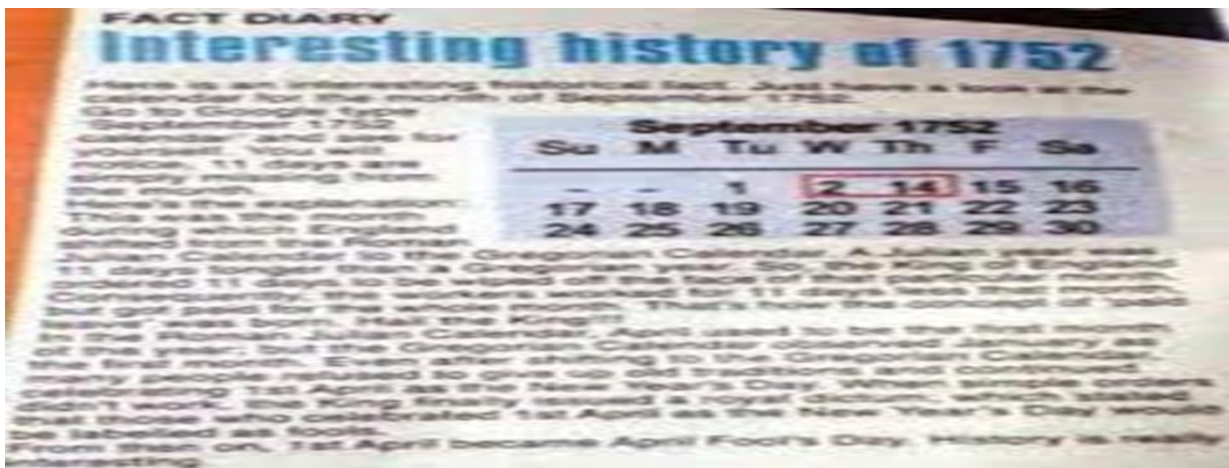
October 4 – Saint Teresa of Ávila dies. She is buried the next day, October 15.

Six and a half million Britons went to bed on September 2, 1752, and woke up on September 14. The reason?

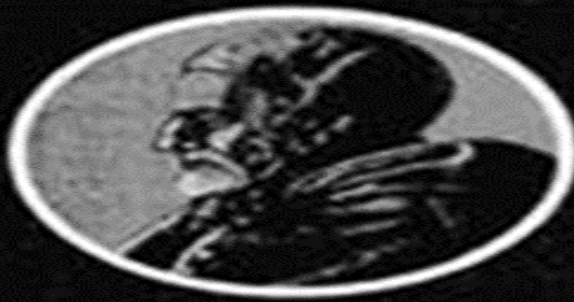
This was the month during which England shifted from the Roman Julian Calendar to the Gregorian Calendar. A Julian year was 11 days longer than the Gregorian year. So, the King of England ordered 11 days to be wiped off the face of that particular month. Consequently, the workers worked for 11 days less that month, but got paid for the whole month. That's how the concept of Paid Leave was born. Hail the King !!!

Publications

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Julius Caesar



Gregory XIII

1582		OCTOBER					1582
SUN	MON	TUE	WED	THU	FRI	SAT	
	1	2	3	4	15	16	
17	18	19	20	21	22	23	
24	25	26	27	28	29	30	
31							

Lastly, Lets talk about Which one of this is a Leap year??

- A. 1922
- B. 1800
- C. 1954
- D. 2010

Which is explain clearly by below mentioned picture:

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February 1922						
Sunday	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday
			1	2	3	4
5	6	7	8	9	10	11
12	13	14	15	16	17	18
19	20	21	22	23	24	25
26	27	28				

February 1800						
Sunday	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday
						1
2	3	4	5	6	7	8
9	10	11	12	13	14	15
16	17	18	19	20	21	22
23	24	25	26	27	28	

<http://luing.altervista.org/calendar/index.php>

<http://luing.altervista.org/calendar/index.php>

February 1954						
Sunday	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday
	1	2	3	4	5	6
7	8	9	10	11	12	13
14	15	16	17	18	19	20
21	22	23	24	25	26	27
28						

February 2010						
Sunday	Monday	Tuesday	Wednesday	Thursday	Friday	Saturday
	1	2	3	4	5	6
7	8	9	10	11	12	13
14	15	16	17	18	19	20
21	22	23	24	25	26	27
28						

<http://luing.altervista.org/calendar/index.php>

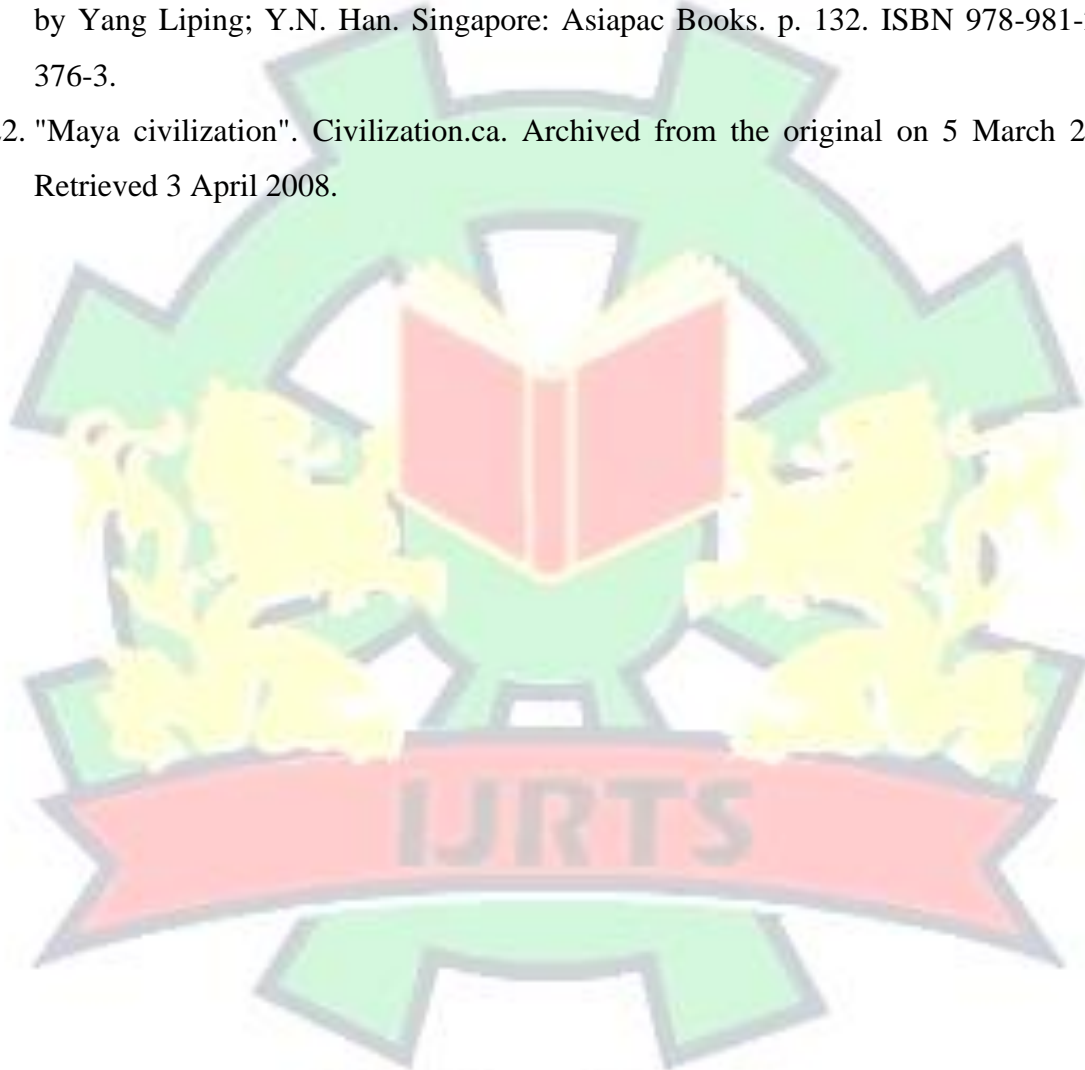
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Publications

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The Study of Computer Applications as Essential Part of Mathematics

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Abstract

Computer science originated from mathematics and has advanced significantly, but is still deeply rooted in mathematics. Many programmers create math tools, while engineers use science for practical applications. Logic is the mathematical foundation of computers. Other mathematical disciplines, such as calculus, probability theory, and set theory, are applied in computer programmes but are not as significant. Russell and Whitehead demonstrated, however, that logic can serve as the foundation for all branches of mathematics. Therefore, even though a computer can implement other types of mathematics, those other types cannot operate a computer. Academically, computer science majors engage in extensive Linear Algebra. In practice, a strong grasp of mathematical concepts assists students in comprehending computers. People with a strong grasp of mathematical concepts can readily comprehend how logic controllers' function, how to compose more efficient algorithms, and how encryption operates.

Content

It is impossible to envision computer science without mathematics. Programming in the field of computer science requires mathematics. For abstract algebra, Sage's software system is indispensable. It is necessary to have an independent, open-source software system for advanced mathematics. The Sage software system is readily accessible on a user's personal computer, a local server, or Sage Math Cloud. The most recent version of Sage is 6.8, which provides accurate results. In addition to Microsoft Excel, we can solve the important abstract algebra query. We employ programming to solve abstract algebra. The majority of algebra is utilized in programming. In this thesis, we will also address the crucial subject of the relationship between mathematics and computer science. In addition, it is used to teach

students throughout their academic careers that computer science is fundamentally about mathematics.

In introductory computer science, we study standard Boolean logic, arithmetic, and set theory, all of which are essential. All of these disciplines employ the "Finomet automata" principle. We determine whether the machine is capable of solving all of our problems. Even though we have an infinite amount of time to solve the problem, allocate resources, and run the programme, machines can do so in a matter of seconds. Functions and iterations, which have become a "prime time" topic for most of us in functional programming, are also very fascinating topics to study.

The primary function of a "electronic" computer is to efficiently perform even the most lengthy and complex calculations. This demonstrates that the first computer capable of encrypting British military communications was constructed by the British at Bletchley Park. Alan Turing is regarded as the first programmer, meaning that today we recognize a programmer, and he firmly believed that there was a great need for more mathematicians who could write programmes, as computer science was essentially a mathematical subject that was a computer. Without mathematics, the inventiveness of science was impossible.

Computer science has only ever originated from mathematics. The field of computer science has advanced significantly over the past few decades, becoming vastly superior to the first computer, but it is still deeply rooted in mathematics. Numerous mathematicians contributed to the development of modern computers, and mathematics is now widely employed in computer science. There are two primary explanations why we typically do not perceive it. We discuss the first one because so many programmers have spent a significant amount of time creating math tools and have shifted the primary implementation away from its mathematical origins. The second is that the majority of us are computer engineers rather than computer scientists. Engineers use science for practical applications.

The code is compiled using a compiler, and the compiler logic functions as the implementation of a mathematical treatise. We use cryptography to identify the complex arithmetic involved in the development of complex algorithms whose codes we are unable to address. We send keys via email and the internet without considering mathematical concepts, despite the fact that they are based on mathematical principles. In developing algorithms and enhancing the efficacy of diverse processes and processors, the application of all mathematics

becomes crucial. The terms buzz and statistical model are used because the use of artificial intelligence and data analysis has increased significantly in the current era, prompting their application. Which is necessary for computer engineers and mathematicians to consume. Since the advent of computers, it is crucial to pay close attention to the development of mathematics and logic. What tasks does the computer perform, or what tasks should the computer perform, that mechanical devices were unable to do? Pascal constructed an electric addition engine in 1642 to make the computer unique. This served as a vacuum tube for the valves, which were replaced by transistors in the future.

The formal logic initially required by Bole, Freeze, Pierce, Russell and Whitehead, Hilbert, Godel, Turing, and others who worked in computer science provided a mathematical foundation for this, and the inventors of calculus established the foundations of physics to solve such problems. In addition to calculus, set theory, and probability theory in mathematics, logic in computers serves as the foundation for computer programmes in other disciplines, such as calculus, set theory, and probability theory. Russell and Whitehead have stated that logic can be applied to all other branches of mathematics; therefore, computer science can be applied to other branches of mathematics. This algebra thesis will be inspired by three problems, which we trust you will find fascinating. Abstract algebra is a theoretical course in which we ponder why things are true rather than how to solve problems. Occasionally, algebra is more concerned with beauty and elegance than with utility and efficacy. You may be compelled to query yourself numerous times about this abstraction and theory.

In light of the preceding examples, algebra is not only practical but also essential. This demonstrates that its applications are vast. Finally, we see how the preceding algebra can be solved. Here, we will attempt to highlight some crucial examples. Results suggest that using this thesis, Excel can teach groups independently about rings. The fact that I will be able to construct it using the RSA algorithm indicates that it can be readily simplified. Previously, I would solve abstract algebra problems with the aid of cryptography, RSA encryption, error detection, and coding.

This thesis will cover the theoretical aspects of groups, rings, and regions. With the advancement of computation over the past several decades, applications of abstract algebra and discrete mathematics are becoming increasingly prevalent and their importance has

grown. Engineers, scientists, and computer scientists are currently searching for an alternative to abstract algebra. A theory of abstract algebra remains crucial to the study of abstract algebra. Utilizing mathematics is one of the most pertinent examples for software development. We cannot programme any software or application without using mathematics. There are numerous varieties of software that can be constructed using matrices, and in some instances, the software significantly relies on matrices for storing or manipulating data or for the user interface. This is required for understanding the inherent properties of matrix groups. Who should comprehend the concept of operating in groups? When creating software components, software developers rely on quaternary groups, which can be understood through group theory. When writing algorithms, knowledge of group theory in abstract algebra is required. When a possible subset of a set is applied, a set is formed, which is a type of group. And it contains components that make up a complete set. Using Lagrange's theorem, we will present a very simple example involving cyclic groups, generators, etc. All of these exist in group theory, which constitutes a substantial portion of abstract algebra. In numerous scientific and computer-related disciplines, abstract algebra is employed in order to determine its prevalence. Specifically, algebraic data consists of a type of combination. Algebraic types can be divided into two general divisions, which can be referred to as classified, dissatisfied union, or distinct types. In mathematics, a group is a collection of elements with a binary operation defined in such a way that the conjugation of any of its two components yields the third element, which satisfies the four restrictions. These restrictions are referred to as arithmetic and include modifier, associativity, instantiation, and inversiveness. The most prevalent example of a group is a collection of integers with addition binary operations; any two integers can also be added together to produce an integer. The abstract formulation of group identities is the segregation of a particular group or the observable manifestation of its operation.

Thus, abstract algebra and beyond have a significant mathematical impact. The pervasiveness of groups in numerous mathematical and nonmathematical disciplines has made them a central organising principle of modern mathematics. Multiple values of a product type typically share the same combination of field types. Set-theory product set of all conceivable values of a product type's field types. Variants can be categorized into a single form of value or, more commonly, several classifications. Therefore, these sorts of values are typically produced by a semi-functional device with its own manufacturer. Here, a predetermined

number of arguments with the predetermined categories are utilized. The set-theoretic sum of all possible values of these types is the sum of all possible values of the type. Algebra analyses the values of the categories of data and, along with them, a value known as its field, which contains the data to be eliminated. The University of Edinburgh demonstrated algebra data types. Which is a modest, 1970s-developed functional programming language? We use programming for the Boolean algebra subfield of abstract algebra. The fundamental algebra used by Boolean is standard algebra as opposed to float. It is, after all, used for domain-specific, which we employ repeatedly for the remainder function. Programming, a subfield of abstract algebra, employs the cyclic group. For instance, a zodiac mod has been adopted. When working in three dimensions, algebra actually works. When an object A is rotated in the direction of 3D and object B is rotated in the opposite direction, their ultimate positions are distinguished by reversing their order. For the 3D curve, AB does not equal BA . So in Algebra, $AB = BA$ cannot be taken. Several tangible algebraic pairs can be derived as follows: We can be derived through movement in three-dimensional space, the purest quadrilateral of which is algebra. Which is used extensively in 3D graphics.

Utilising mathematics is one of the most pertinent examples for software development. We cannot programme any software or application without using mathematics. There are numerous varieties of software that can be constructed using matrices, and in some instances, the software significantly relies on matrices for storing or manipulating data or for the user interface. This is required for understanding the inherent properties of matrix groups.

Who should comprehend the concept of operating in groups? When creating software components, software developers rely on quaternary groups, which can be understood through group theory. When writing algorithms, knowledge of group theory in abstract algebra is required. When a possible subset of a set is applied, a set is formed, which is a type of group. And it contains components that make up a complete set. Using Lagrange's theorem, we will present a very simple example involving cyclic groups, generators, etc. All of these exist in group theory, which constitutes a substantial portion of abstract algebra. In numerous scientific and computer-related disciplines, abstract algebra is employed in order to determine its prevalence. In physics, for instance, groups are utilised to represent symmetry operations. Group theory is used to simplify differential equations. This form of system is utilised for equation reduction. In algebra, for instance, the number of axes of a theory is equal to its amplitude.

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A Study of Computer Programming as an Experimental Results in Mathematics

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Abstract

Because mathematics is a fundamental component of computer systems, every programmer and computer scientist must possess fundamental mathematical skills. The type and quality of mathematics required depends on the computer science field in which you wish to operate. Some career paths in computer science require minimal mathematical knowledge. For instance, if you know fundamental programming languages and can perform simple arithmetic, you can develop fundamental mobile applications. Nonetheless, if you wish to advance in your computer science career, perhaps by investigating advanced disciplines such as computer security or operating systems, you will need to have a firm grasp of advanced math concepts. No one expects computer scientists to be mathematicians, but learning algebra and calculus in high school and college can open doors. Computers are increasingly utilised in mathematics education as a result of consumer preferences and market trends. However, these implementations can result in the loss of algebraic computation skills and a lack of comprehension among students. Software can improve students' mathematical reasoning, problem-solving, and creativity, thereby enhancing their awareness.

Content

This century has frequently been referred to as the "age of information" due to the fact that communication and computer technology now play a significant role in our daily existence. There is a desire to incorporate new technological characteristics into academic programmes to foster excellence in education and to facilitate new developments. In this new era, the development of information and communication technologies is of the utmost importance. In the ultramodern era, a PC is a multifunctional device that has provided us with unique opportunities for teaching and learning. When coaching science is utilised, it creates an environment that enhances

students' educational experiences. As a result, the contemporary development of computer technology has impacted all global, cultural, and economic lifestyles as well as educational settings, and computers no longer play a significant role in the roles of students and teachers. Identical requirements are utilised in the provided instructions. In a constructivist manner, education software can equate and motivate pupil mastery and significant mastery for all students. When a few talent environments are introduced using computers during the academic cycle, teacher-centered guidance transforms into student-centered guidance.

The use of computers in arithmetic instruction is gaining popularity. However, consumer characteristics and numerous other variables define the method's application in mathematics education. Previously, instructors and academic researchers viewed computer systems as supplementary tools, comparable to an overhead projector, a slide, or a television. Computers were initially utilised for simple calculations as opposed to delivering vibrant digital pages with colour photographs and instructing college students in the formation of their knowledge. It failed to provide a solid foundation for substantial education and excellent learning practices. Some research also reveals that these common applications of computers in education have negative effects. The use of computers as calculators, where students can see the result of a simple calculation, is an example of the widespread application of computers to the instruction of algebraic principles and procedures. According to Norcliffe (1996), pupils will lose their algebraic computation skills if they exclusively use computers for simple calculations. Mackie (2002) stated that students can reach a point where they lack comprehension of what they are doing when typing certain keys required to rectify system problems in computer software. Huh. Historically, some summary mathematical standards have been difficult to comprehend. Due to the abstract nature of arithmetic, it can be difficult for many students to understand summary concepts. Nonetheless, this concern can be eliminated or mitigated by using concrete equipment supplied by a software programme (Bucky, 2002). It is possible to use technological tools to enhance mathematical reasoning, problem-solving, and creativity in order to build student awareness skills.

This software discourages memorization and prevents college students from performing tedious procedures. In conjunction with software-based algebraic knowledge, the qualitative application of data processing science in college mathematics education is

one of the most frequently discussed topics among researchers. The implementation of new arithmetic teaching methods can influence the development of instructional policy, strategy, and structure (Ersoy, 2003). The high-tech industries of the twenty-first century have made abstract reasoning and creative abstraction skills crucial for global economic competition. After establishing the logical foundation of mathematics, students can only develop their capacity for abstract thought. Students will be able to acquire problem-solving and fundamental inquiry skills as a result of the enhancement of data processing instruments used in mathematics education. In this procedure, the social constructivist approach must be utilised to develop skills such as exploration, analysis, and algorithms for the underlying logical relationships of ideas (Hacısalihoglu et al., 2003). Extensive research has been conducted on the efficacy of computer-assisted instruction. These studies have demonstrated that student achievement is significantly higher in computer-equipped classrooms than in those without computers (Krishnamani and Kimmins, 2001; Rainbolt, 2002; van & Dempsey). In addition, some researchers assert that computer-assisted practise led to the development of higher-order thinking skills and that students analysed the concept rather than memorization (Rainshaw & Taylor, 2000; Ubuz, 2002).

To develop by comprehending concepts based predominantly on constructivist philosophy. Utilising fact processing can result in a more fruitful and effective method of acquiring environmental knowledge. In such an environment, the student is able to analyse complex query problems, hypothesis viable solutions, formulate hypotheses, and generalize. Through the use of software programmes, the student is able to create their own mathematical compositions. The software programme functions as a tool for the student to explore information, ideas, and statistics in the teacher-designed work environment. Participation in the project is possible if the student employs all of their abilities to manipulate their own learning. Derive TM and Mathematica TM are two software applications that facilitate this type of education. The vast majority of software applications available on the market require minimal programming expertise. Excel, for instance, is a programme that can be utilised with a pen, paper, or calculator.

In an introduction to summary algebra, it is particularly difficult for students to comprehend and maintain positive attitudes towards concepts such as the team concept (Campbell & Zuckis, 2002). It is possible to implement the majority of these concepts

and make them accessible to college students for similar learning through the use of computer technology. The majority of abstract ideas can be condensed using this method, making them simple for students to comprehend (Bucky, 2002). Today, numerous software applications are utilised for the instruction and acquisition of summary algebra knowledge. The Magma TM PC software programme is one of the products (Quinlan, 2007). The application language Magma TM is based on elementary algebraic principles and meaning. In addition, Magma TM is preferable to Maple TM and Mathematica TM for reading algebra, in my opinion (Kulich, 2000). It is developed by a team of computer algebra researchers at the University of Sydney with support from the Australian Research Association. Between 1975 and 1985, the group also developed the Kelly TM software to teach crew principles in related theoretical domains. Kelly TM served as the inspiration for the development of Magma TM. Kelly TM focuses on a specific area of Algebra. Magma TM, which encompasses the standard algebra system. Magma TM is one of the PC software programmes that can be used for algebra, a variety of theories, and geometry in its entirety. Common classes utilised to teach Magma TM include finite groups, half-groups, rings, spheres, and geometric structures. In addition, the software has been utilised in advanced PC algebra courses, introduction to team concept courses, and vector area instruction (Cannon and Playway, 1997). Matlab TM and Maple TM are two of the most popular packages used in summary algebra (Quinlan, 2007). Originally developed at the University of Waterloo in Canada, the Maple TM PC algebra processor is a product of Maplesoft. Charlwood (2002) examined the teaching of symmetric and matrix organisations in an introductory summary algebra course and found that the learning strategy augmented by Maple TM PC working pages contributed to the development of students' proofreading skills.

In their e-book titled "Learn Abstract Algebra with ISETL," Dubinsky and Leran (1994) state that their understanding is predicated on the optimistic belief that college students want to engage in intellectual activities. This imaginative and prescient understanding has been developed by other researchers in recent shared research (Seldon, 2005). Before students can comprehend any presentation of abstract mathematics, these computer-supported activities provide an experiential foundation for any future verbal explanations. In addition, they may examine these explanations through the lens of their cognitive experiences. This method is based on extensive theoretical and empirical research as well as the authors'

extensive experience teaching summary algebra. Computer construction, particularly minor applications written in the programming language ISETL, constitutes the foundational tasks in an algebraic orientation. The most essential tool for reflection is the discussion and debate of tasks in groups of two to four students. Due to the similarity between contemporary written mathematics and ISETL expressions, there are few preliminary data regarding the use of ISETL. ACE (Activities, Classroom Discussion Materials, Practise) is a learning circle developed by researchers utilising ISETL-based learning techniques. Curriculum is divided into weekly units by the instruction cycle. Each week, college students spend a few days and the remainder of the week in a computer center in a regular classroom. The textbook contains closely related definitions and evidence regarding constructions in activities, written in a casual, suggestive tone.

Numerous arithmetic teachers, including the first author, have utilised ISETL to teach algebra in summary form. Researchers have reached a consensus that ISETL is an excellent learning tool (Pessenen & Malvella, 2000; Khrishamani & Kimmins, 2001; Smith, 2002; Okur, 2006; Weller et al., 2002). FGB (Finite Group Behaviour), an enhanced Windows TM paradigm of ESG (the discovery of small groups), was once authored by Edward Keppelman. The programme shares some similarities with ESG, which operates on DOS (Kulich, 2000). It is used to teach the fundamentals of team theory, including the search order of groups, subgroups, frequent subgroups, issue agencies, and elements. The FGB examines whether an algebraic shape, including its factors, resources, and amenities, can be grouped using operation tables. The colour coding of crew tables allows students to observe subdivisions and issue agencies of a crew on their display screen (Perry, 2004). According to Kepelman and Webb (2002), the FGB can be used to rehearse and practise concrete examples of corporations with beginning students. The application operates on a library containing all finite organisations up to isomorphisms of order 16 or less and all non-Abelian organisations of order 40 or less. Tables for companies can be saved as text archives with a distinct format, and clients can make notes on the searchable attributes. Each inquiry is then stored in a distinct section of the personnel information file.

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A Critical Study of Relationship Between Core Programming in Computer Science and Mathematics

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Abstract

Most users are likely aware of the close connection between mathematics and computers. The most fundamental form of numbers is what brings entertainment, communication, and information into our residences and to our fingertips. In the realm of programming, logic and mathematics go hand in hand. Math is one of the most important skills for programmers to have in order to develop complex applications. If a programmer lacks mathematical knowledge, they are essentially handicapped. Therefore, it is essential for all programmers to have a firm grasp of certain mathematical concepts that are fundamental to programming. Not only is mathematics closely related to computer programming, but it also functions as its basis. Math is a requirement for college degrees in computer science, information technology, information systems, and other related disciplines. To completely comprehend computer programming and the science underlying its operation, it is necessary to first comprehend the mathematics underlying its structure. Although it is impossible to summarize all of the mathematical mechanics of computer science in a single article, a few of the most essential are highlighted below. The fundamental functionality of any computer is founded on binary arithmetic, and each integer in the computer is represented in binary. Essential to low-level hardware programming are binary reading and elementary mathematical operations. Numerous programming duties, such as altering the color of an object, require familiarity with the hexadecimal number system. Numerous programming duties necessitate the use of standard arithmetic, and virtually every programme ever written employs addition, subtraction, multiplication, and division. Numerous computer programmers will encounter problems that can be solved with algebra. Computer science studies computational approaches for solving abstract algebra, focusing on various subfields like artificial intelligence, security, and software package engineering. These algorithms address both theoretical and practical challenges, ensuring the efficiency of computer hardware and

software. Computer science is a department of mathematics, as abstract algebra plays a crucial role in programming.

Content

Computational approaches for solving abstract algebra are studied in computer science. The majority of electrical and laptop engineers don't have a vested interest in the theory, design, development, or applications of the software packages and systems that laptop scientists rely on daily. Artificial intelligence, laptop systems and networks, security, information systems, human-computer interaction, vision and graphics, programming languages, and software package engineering are some of the most popular subfields within computing. Understanding the significance of the programmes to the study of computer science is just half the battle, however. Algorithms for solving programmes and researching the efficiency of computer hardware and software have been developed and examined by computer scientists. The challenges faced by computer scientists range from the theoretical, such as determining what problems may be addressed by improving computer hardware and software, to the practical, such as developing programmes for mobile devices. There are sophisticated ones that don't need a lot of technical know-how to use and keep secure.

You use your computer for everything from watching Netflix and playing video games to doing work and exploring websites like Reddit, but have you ever honestly considered how computers function? Also, what exactly is computer science? The simplest explanation may be found in computer science:

The field of computer science encompasses research into computing devices, computational theory, the development of software and hardware, algorithm design, and user interaction with computers. A high wage may sway some people to choose a profession in computer science since it is challenging and financially rewarding but also relatively invisible. Computer scientists strive to find the best possible solution to any given issue. In order to solve an issue, a computer scientist will collect relevant information, learn the appropriate programming language(s), and make sound judgements (including the usage of scripts). To tackle the issue, they write a set of rules or instructions that the machine can follow. To learn more about the field of computer science, its potential career paths, and the tools and resources available to those interested in entering it. In today's world, disconnecting from technology is a need. Even when we're not actively using any electronic devices, it might seem like our homes and communities are always buzzing with the activity of electronic

machinery. We definitely live in a technologically dependent culture. There are now two distinct groups of individuals in modern society: those who allow technology to control their daily activities, and those who create such revolutionary tools. On whose side will you be?

To what extent does mathematics play a role in computer science? No one has done a good job of describing it. To me, this reads like pure mathematics. Second, I was wondering if anybody knew of any PC science or programming books that were very dogmatic and axiomatic. Array's primary focus is on laying a solid groundwork for future computer scientists and programmers. Since mathematical imagination is necessary for computer science, computer science is a branch of mathematics. Programming heavily relies on abstract algebra.

Many students (myself included) educational horizons have been confined to the realm of mathematics once you consider that this has been proven. However, students are drawn to computer science in the hopes of one day becoming programming masterminds. "BIs it truly that simple, mam tai.n? In a word, no, my good buddy. Maths and computing go hand in hand fairly naturally. Are you familiar with the notations used in logic, set theory, combinatorics, sketch theory, probability, number theory, algebra, etc.? Don't worry, they're all just different levels of discrete mathematics that provide a solid basis for programming and computer science (and the desire to learn more about them is what motivated me to do so in the first place). Algebra is a good illustration of this. Relational algebra, in contrast to Boolean algebra, is employed in databases and logic gates. If you need further evidence, consider the various uses of quantity theory in cryptography and cryptanalysis.

You've probably all heard of algorithms by now (if you haven't, go back and study!). They're a crucial part of computer science. Simply said, they are a collection of guidelines for carrying out a programme or application [79]. An algorithm was originally employed in a math class rather than a computer science one. The expression " $2 + 3 = 5$ " will serve as an easy case to discuss now. This is a simple mathematical method for representing the sum of two and three. Thus, it is clear that mathematics is crucial for even the most fundamental understanding of the sophisticated algorithms used in modern computer science. To fully grasp this mathematical theorem, you'll recall that any two $n \times n$ metrics of the number were required. That B has the "largest common true divisor" is evident. In other words, A and B can be divided correctly by some number. That has reached the age of maturity as an integer A, where $A = AD$ and $B = B'D$ 'and B. D may be evenly divided by every right divisor of A and B. While many programmers are busy developing mathematical techniques, many

mathematicians are really secretly backing the computer science community. Many of the instances I've seen of mathematical memory that are certain of the exact algorithm have fascinated me. The mathematicians who wrote the thesis employ language that is both precise and significantly less natural than the comprehensive algorithm approach available to modern computer scientists. The majority of the 35-page content, for instance, may be delivered in around two pages via Abraham Wald's writing when the material is presented in algorithmic phrases, and numerous instances can be offered. However, "Ibis" is a problem for any alternative viewpoint.

While many programmers are busy developing mathematical techniques, many mathematicians are really secretly backing the computer science community. Many of the instances I've seen of mathematical memory that are certain of the exact algorithm have fascinated me. The mathematicians who wrote the thesis employ language that is both precise and significantly less natural than the comprehensive algorithm approach available to modern computer scientists.

Abstract algebra is used to teach most arithmetic ideas. Learning about different programming languages is an individual concern of computer science. These languages tend to be summary as well. Characteristics include the use of grammar, clearly defined procedures, symbols, single words, and even pictures. Since mathematics is abstract, it will provide you with a blank slate upon which to build your knowledge of a programming language. The skills of reading, perceiving, and observing a problem in order to arrive at a solution are all skills that a student will acquire via studying mathematics. When the fields of computer science and programming mature, these skills will be indispensable. Algorithm is a word often used in the fields of computer science and technology. It lays the groundwork for the development and deployment of any programme or application. Most human laptops discover a time period in science class, but they were really raised in a maths class, contrary to popular assumption. Do you aware that a basic algorithm is an equation like $7+3=10$? In college, students are introduced to complex algorithms using seemingly simple formulae.

No longer is coding done only by one person. You must double-check your work afterwards to be sure it is right. Reason being, when coding, you could make a mistake or two. College-level mathematics coursework often includes mandatory self-evaluation. There has been a decline in the reliability of the answers you provide. You'll want to go back and review your original data and formula. If you make a mistake, you may fix it before finding out the remainder of the solution. In other words, learning mathematics helps pupils become ready

for the work involved in preparing and revising them once they reach college. College students will not be interested in doing this task even if there are machines that can do it automatically.

In addition to the aforementioned abilities, abstract algebra is also very useful in the field of computer science. Here, you'll have the chance to apply what you've learned in abstract algebra to the real world by using a digital calculator. You will be sad if you don't pay attention in maths class. To plot the programme needed to control an autonomous vehicle, for instance, several mathematical formulae and formulations are employed. It would be quite challenging to construct such packages if they were no longer feasible, especially for someone who is absolutely new in the discipline of mathematics.

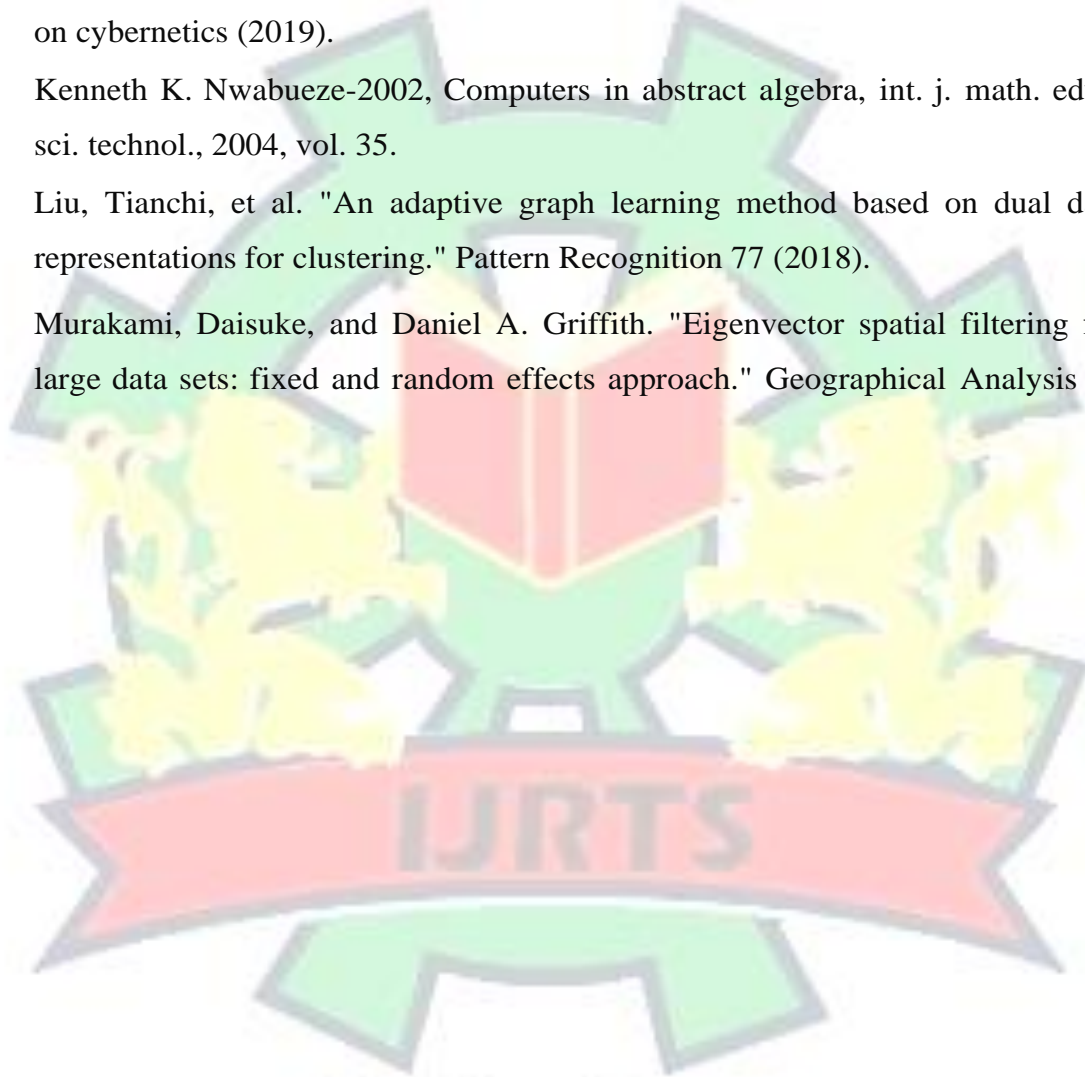
Learning the many forms of algebra required for computer science and programming is no longer a secret. In-depth understanding of algorithms, computability, and complexity is gained via this. What is Abstract Algebra? A problem with feature manipulation that arises in the field of discrete mathematics. Learning these mathematical concepts will make their practical application in computer science much easier.

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Publications

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A Study of the Safe Multiparty Equality Check with Homotopy Cryptosystem and the Design of Safe Multiparty Computing Protocols for Privacy

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Abstract

The subject of information security where each party wants to compute some function of their private inputs without disclosing their inputs to one another is called Secure Multi-Party Computation (SMC). For example two banks want to perform certain data mining operation over union of their individual databases but no bank wants to disclose its private database to other. In such scenario SMC becomes relevant. we present our research work which is aimed to designing novel SMC protocols for preserving privacy of an individual data while joint computation is taking place. Due to huge growth of the Internet and its easy access by common man opportunities are increasing for cooperative computations by many physically distributed parties. All these parties have their own data and for the sake of mutual benefit, they want to compute some function of these data. Traditionally, the computation is easy; every one supplies the data input to get the value of the function. But as this model is growing the need of confidentiality of data is also increasing. Multiple parties want to know the result of computation but they are also concerned about secrecy of their data inputs. The subject of SMC has been evolved from two party comparison problems to multiparty image template matching problems. Many specific SMC problems have been defined and analyzed by researchers like Private Information Retrieval, Selective Function Evaluation, Privacy-Preserving Database Query, Privacy-Preserving Geometric Computation, Privacy-Preserving Statistical Analysis, Privacy-Preserving Intrusion Detection and Privacy-Preserving Cooperative Scientific Computation. Based on these general SMC problems many real life applications emerged like Privacy-Preserving Electronic Voting, Privacy-Preserving Bidding

and Auctions, Privacy-Preserving Social Network Analysis, and Privacy-Preserving Signature and Face Detection. Many protocols assume the network to be secure while devising SMC algorithms. In this work we proposed novel and improved secure sum computation algorithms, and equality check algorithms which are application to insecure networks. We used public key cryptography techniques to hide the value of the data while it is being travelling on network lines. We also used additive homomorphic public key cryptosystem where the sum of encrypted values is same as the encryption of the sum. Another technique we used in our algorithms is the Shamir's Threshold encryption scheme which allows a value to be reconstructed out of certain minimum number of shares of the value. Our proposed secure sum algorithms and secure equality check algorithms are based on real model as well as Ideal model of SMC. In one of the proposed algorithms we coin the term Semi-Ideal model as it is a hybrid of ideal and real model of SMC. The reason behind working on the secure sum is that all other arithmetic operations can be implemented using addition operation. Secure equality is also used in many database searches as well as in the process of authentication. Our works are suitable for honest but curious parties. Further, the algorithms can be devised which are suitable for malicious parties.

INTRODUCTION

Present age is rapidly becoming the "electronic" or "online" age. Almost every transaction of our day-to-day life involves electronic transactions. May it be transferring money by the bank or individual from one account to another account. We are almost in the age of e-banking. Due to the huge growth of e-commerce, orders are being placed online, whether it is business-to-business or business-to-customer. These orders are being processed using Supply Chain Management (SCM) software and the database is maintained in association with Customer Relationship Management (CRM) software. The modules such as SCM and CRM are nowadays integrated in Enterprise Resource Planning (ERP) software. All these software are using the database distributed along multiple sites with the help of the Internet. Bidding and auction activities are also being performed online. The computation on the database may be done at a central site or it may be distributed at multiple physical sites in a distributed system. Owing to easy access of the Internet by common man, its fast transactions, and huge growth in every sector people are being dependent on it for almost all kinds of computations. A new paradigm is rapidly emerging where multiple parties jointly work and perform computation on their data but they are bothered about the privacy of their data while the

computation is being performed by them. In other words these joint parties do not have full trust on each other. But for the sake mutual benefit they have to work in cooperation. These cooperating parties wish to evaluate some common function of their sensitive data but do not want to disclose individual data to one another. In other words it is a paradigm of cooperative computation of multiple parties where all the parties are interested in knowing the result of joint computation on their data but they are unwilling to disclose actual value of their private data to other parties involved in the computation. All want that the privacy or confidentiality of their individual data must be preserved while the computation is being performed. Therefore it is a paradigm of privacy preserving computation. This area of information security where multiple cooperating parties wish to evaluate a common function of their data jointly such that the data is not disclosed to other parties is known as Secure Multiparty Computation (SMC). Consider a scenario where multiple banks jointly wish to find common defaulters from their databases so that the list of wilful defaulters as well a list of defaulters from multiple banks can be prepared. This is in common interest of all the banks involved in the computation. But all the banks are concerned about the privacy of their database. As they are competitors in the same sector, they don't want to disclose their database to other banks. In this scenario some protocol is needed where all the banks may know the list of common defaulters keeping their database secret.

Many specific protocols for SMC. Few of them are listed below:

1. Private Information Retrieval (PIR)
2. Selective Function Evaluation
3. Privacy-Preserving Database Query
4. Privacy-Preserving Geometric Computation
5. Privacy-Preserving Statistical Analysis
6. Privacy-Preserving Intrusion Detection
7. Privacy-Preserving Cooperative Scientific Computation

Since the inception of SMC many interesting and real life applications emerged. Few of them are listed below:

1. Privacy-Preserving Electronic Voting
2. Privacy-Preserving Bidding and Auctions
3. Privacy-Preserving Social Network Analysis
4. Privacy-Preserving Signature and Face Detection

Early research of the SMC focussed on circuit evaluation techniques using combinational logic circuits. But these techniques were complex and expensive. Nowadays there are three main techniques used to devise SMC protocols:

1. Cryptographic SMC methods
2. Randomization SMC Methods
3. Anonymization SMC methods

Other set of protocols in these are used to check the equality of data with multiple parties keeping the private data secret. These protocols can be used for matching or verification of a particular data without violating or hurting the privacy of actual data. We propose a novel multiparty equality protocol using homomorphic cryptosystem. The protocol is also suitable for insecure networks as the data flow in encrypted form. We have proposed Secure MultiEqualityCheck protocol based on homomorphic public-key cryptosystem [10]. In another protocol we used hashing function for generating hash code of the actual data. A hash function produces a fixed size value when a variable-size input is given. This is infeasible to produce the actual value when a hash code of that value is given. But if so happens it is called collision. We assume that there exists a collision-resistant hash function. Based on this two protocols EqualityHashCheck for ideal model and EqualityHashCheck for real model are proposed.

EXAMINING PROBLEMS

We have observed that many protocols for SMC assume the network to be secure. Therefore we propose novel protocols which are suitable for insecure networks. This is achieved by using some cryptographic technique and not sending the data in its actual form. For example the secure sum protocol proposed by Clifton et al. [4] uses a random number that is added to the actual data and the partial sum is sent on the network lines. The parties are arranged in a ring. If some intruder intercepts both incoming and outgoing line of a party, the difference of the partial sum he gets is the actual data of the party. Similarly, if two neighbours to a middle party share what they send and what they receive, the difference is the actual data of the middle party. In actual practice this is possible as the lines are always insecure. We use cryptographic methods to design the protocols for insecure networks. In our protocols the above cases of data breach are not possible. In our secure sum protocols we encrypt the data with either symmetric or asymmetric encryption before sending on the lines. Thus, the privacy of the data against intruders and colluding neighbours is preserved. We proposed secure protocols both for real model and ideal model of SMC. In ideal model the

trustworthiness of the third party is a concern. With homomorphic cryptosystem and threshold encryption scheme we ensure that the privacy of the data is also preserved against third party. We also proposed multiparty equality check protocols in ideal as well as real model. We use additive homomorphic encryption in one protocol and hash function in other two protocols. Thus, we have eight secure sum protocols and three multiparty equality check protocols proposed. All the protocols are suitable for semi-honest adversaries.

Novel secure sum protocols

Secure sum with symmetric key for insecure networks

This approach of getting secure sum is slight extension of the secure sum protocol of Clifton et al. [4]. All the parties are arranged in a logical ring. All the parties share a symmetric key in advance. A party is chosen as protocol initiator. It chooses a secret random number, and adds to its private data. Now, the party encrypts this sum with the shared key, and sends the encrypted value to the immediate neighbour party in the ring. The receiving party first decrypts the ciphertext to recover the partial sum, and adds its private data to the recovered sum. Now, it encrypts and sends in similar fashion as the previous party did. This process of receive, decrypt, add, encrypt, and send is repeated until the protocol initiator recovers the sum of all the data plus random number. It subtracts the random number to get the sum of actual data and distributes to all the parties. The improvement over the secure sum protocol of Clifton et al. [4] is that our protocol is suitable for insecure networks as the data flow in encrypted form, no intruder can learn the data or the partial sum. We analyse the computation and communication complexity of the protocol.

Secure sum with asymmetric key for insecure networks

The secure sum with symmetric key protocol provides an improvement over secure sum protocol. But there are two limitations with it. One, securely sharing the symmetric key. The key distribution on insecure network is itself is a serious problem. Another, as all parties share a common key and when the key is compromised, private of all the data will be violated. In order to overcome this drawback, we propose using public key cryptography with secure sum architecture. In the new protocol all the parties use the same ring topology used by previous protocols and communicate in only one direction. All the parties are required to generate a public-private key pair using certain algorithm like RSA. Each party must share its public key with the previous party in the ring. Thus, the protocol initiator has already got the public key of the next party. The private keys of all the parties are their secret. The protocol initiator party chooses a secret random number, and adds to its private data. Now, it is

encrypted by the public key of the next party, and the ciphertext is sent to the immediate neighbour party. The party receiving the ciphertext first decrypts it with its private key. The partial sum is recovered as it was encrypted with the public key of the receiving party. Now, this party adds its data to this partial sum and encrypts with the public key of the next party, and sends this ciphertext to the next party. This process of receive, decrypt, add, encrypt, and send is repeated until the protocol initiator recovers the sum of all the data plus random number. It subtracts the random number to get the sum of actual data and distributes to all the parties. The improvement is that as there are different encryption and decryption keys at each point, the whole network cannot be compromised with compromising a single or multiple keys. Another, the distribution of public key can be done on insecure network, and public key is no longer a secret. But the cost of generation of public-private key pair at each party is a drawback.

Secure sum against colluding neighbours for insecure networks

Both the protocols secure sum with symmetric key and secure sum with asymmetric key can see colluding neighbours attack. To solve this problem we propose a protocol secure sum against colluding neighbours in which all parties except the initiator party encrypt with the public key of the initiator. Thus, the encrypted value at the initiator contains data with multiple encryptions with the public key of the initiator. When it is decrypted multiple times with the private key of the initiator, the secure sum is obtained.

Secure sum using homomorphic cryptosystem

We proposed a novel protocol in which we use the additive homomorphic property of a public key cryptosystem. The property ensures that the sum of ciphertexts is equal to ciphertext of sum of individual data. We use again propose to use the ring topology as used in the previous protocols. The concept of the initiator and random number is again applicable here. The new thing is that here the initiator party generates public-private key pair using some homomorphic public key cryptosystem. The random number and the private key are the secrets of the initiator but it distributes its public key to all cooperating parties who wish to compute privacy-preserving sum. The initiator begins by adding random number to its secret data, and encrypting by its public key. It sends the encrypted value to the immediate party in the ring. The receiving party simply adds its public-key encrypted value to the received sum, and sends new sum to the next party. This process is repeated until the initiator receives sum of all the ciphertext encrypted with its public key. As per the homomorphic property it is equal to the ciphertext of the sum of data plus random number. When it is decrypted with the

private key of the initiator, secure sum plus random number is obtained. The initiator subtracts the random number to get the secure sum. The sum is distributed to all the parties.

Secure sum using threshold encryption

This approach of getting secure sum uses Shamir's threshold encryption scheme [6] which ensures that if a value is broken into shares; the same value can be reconstructed by adding certain minimum number of shares. We also use additive homomorphic cryptosystem based on public key cryptography. In this protocol we used a hybrid architecture of real as well as ideal SMC model. We coin the term Semi-ideal SMC model for this architecture. A Trusted Third Party (TTP) assists run the protocol. But there is a limited role of the TTP. It generates public-private key pair using additive homomorphic public key cryptosystem, and shares of the private key using Shamir's threshold encryption scheme. The TTP distributes public key and a share of the private key to all the parties. Now, remaining part of the protocol is run by the parties in real model. The parties first get the sum of ciphertexts of their private data with public key encryption. Now, parties recover the private key by getting sum of their private key shares. The protocol initiator decrypts the ciphertext sum with the recovered private key to get the actual secure sum. The sum is now announced to the parties.

Randomization approach for secure sum in ideal model

The protocol for secure sum proposed by Clifton et al. [4] and many other protocols in the literature employ real SMC model where no TTP exists. We proposed a simple secure sum protocol using randomization approach for ideal model. It uses a TTP to assist in secure sum computation. All the parties simply choose and agree on a common random number which is hidden from the TTP. The participating parties who seek to calculate the sum of their sensitive data multiply their data with the secret random number and provide their product to the TTP. The TTP computes sum of all such products and return back to all the parties. All the parties divide the received sum to get the actual sum of data. The privacy of data among parties is preserved as parties never communicate data with one another. The privacy of the data is preserved against TTP as it does not know the random number. The network lines can be insecure as the intruder cannot learn actual data as it is protected by the random number.

Novel Multiparty Equality Check protocols

Real life scenarios exist where multiple parties are interested to check the equality of their data. But the problem becomes crucial when these parties are worried about the privacy of their individual data. This problem can be solved with protocols of SMC. It motivated us to devise protocols for multiparty equality check. We proposed protocols using homomorphic

cryptosystem, and using hash function. We propose protocols for ideal as well as real SMC model.

Secure Multiparty Equality check using homomorphic cryptosystem

This approach of comparing data of multiple parties is an extension to two-party equality check problem. Two parties can check the equality of their data using homomorphic encryption preserving the privacy of their individual data. We call the protocol as EqualityCheck. For multiparty case the parties are arranged in a logical ring. This is a real model as no TTP is present. The protocol begins when all the parties generate their public-private key pair using homomorphic cryptosystem. A protocol initiator party check equality with the next party using EqualityCheck. If equality holds it checks the equality with next party. Thus, a pair-by-pair equality check continues until the protocol initiator is reached. If the equality holds at the end of the ring, the result is declared as “All are Equal” and if the equality breaks at any point in the ring, the result is declared as “All not Equal”. The parties encrypt using public key and protect their data using a random number.

- To maintain privacy of individual data input while computation: The protocols discussed in this thesis allow parties to compute the sum of their inputs while maintaining their data privacy. We use additive homomorphic cryptosystem and threshold encryption to protect the privacy during computation. Our multiparty equality check protocols use encryption and hashes to protect the privacy during equality check. No party ever passes actual data to other party. Thus the privacy of the individual data is maintained.
- To get the correct result of the computation: Our proposed protocols in this thesis perform addition operation over the encrypted data of the parties. The decryption of this ciphertext is the sum of these actual data. Thus, we are able to provide the correct result to the parties while maintaining their privacy. Similarly, we get the result of equality check while preserving privacy.
- To reduce probability of data leakage: In this thesis protocols are designed keeping an eye over the colluding neighbors who want to learn the secret data of the middle party. We successfully reduced the probability of data leakage in case two neighbors collude to learn the data of a middle party in our protocols is almost nil due to use of cryptographic techniques and the hashes

- To minimize communication cost: We aim to achieve the security with minimum communication cost. For the same reason we analysed our work for minimum communication. In real model we build unidirectional ring between the parties. The links are assumed to be insecure. In ideal model parties communicate with TP only reducing the communication
- To develop Semi-ideal model of SMC: We aim to develop semi-ideal model of SMC in which we use properties of both real model and ideal model to achieve privacy and correctness of the result. The term semi-ideal SMC model is coined by us in [9]. Such a model is highly economical and less vulnerable to security threats.
- To develop secure sum algorithms: Many secure sum algorithms are developed and their performance is analyzed for semi honest parties

Scope of the work

The relevancy of the SMC solutions is enhancing day-by-day because a large volume of cooperative computations are being taking place over data of many parties. Many organizations jointly work on a single project and they frequently share their sensitive information. But each of the organizations is also worried about the privacy of its data. In such scenario SMC solutions play significant role. Secure sum computation is an important example as well as the component of the toolkit for the SMC solution. Equality of data with many joint parties could be of real life use. In this thesis we proposed protocols and the corresponding algorithms for the secure sum computation and secure equality check.

- ❖ Banking Sector: Nowadays many banks work cooperatively and frequently share important information about their customers. SMC solutions allow each of the banks to perform cooperative computations while keeping their individual information a secret. For example the bank may wish to get details of common defaulters from their individual databases. But due to their business interest they do not want to disclose actual database to other banks. A privacy-preserving SMC solution called privacy-preserving data mining is useful in this scenario.
- ❖ Query Over a Group of Databases: When many organizations want to cooperatively work over a group of their databases, traditionally each must know the database of the

other. SMC techniques allow executing queries over the group of these databases without disclosing the individual database to each other. For example police organizations of different Indian states want to find certain information over their collective data records; they can use SMC techniques to explore criminal activities, terrorist activities and many other valuable patterns without disclosing their actual database to other police organization. The same technique can be used by intelligence organizations of different countries.

- ❖ Joint Audit: Two or more organizations working as partners can allow audit over their joint accounts without disclosing the individual account details to one another
- ❖ Privacy-Preserving Social Network Analysis: The police or the intelligence organizations perform social network analysis to find the relation between their social behavior and the criminal activity but the laws prevent them to do so because of the privacy concern of the citizens. SMC solutions can be used here by the government allowing to perform analysis over its database without the organization knowing exact details of the individual records.
- ❖ Privacy-Preserving Auction Management: Auction can be managed without disclosing the individual details to the participating parties.
- ❖ Privacy-Preserving Electronic Voting: Online voting can be performed using the technique of SMC which allows individual voting pattern to be a secret.
- ❖ Privacy-Preserving Mobile Phone Services: The mobile phone service can be implemented such the location information of the subscriber cannot be known to the service provider or to any other party.
- ❖ Privacy-Preserving Medical Diagnostic System: It allows the patients to know their disease without disclosing their identity to others or even the result of the diagnostic.
- ❖ Privacy-Preserving Monitoring in Wireless Sensor Networks: The data sensed by the nodes of a wireless sensor network may be made confidential by using SMC techniques because the sensor network is also deployed at sensitive places like military battlefield where there is a need to make the retrieved data secure while allowing computation over the data by the set of nodes.
- ❖ Privacy-Preserving Signature and Face Detection: Sometimes it becomes necessary to keep the signature or the face of a person secret while matching it with the stored database. In this situation SMC techniques are useful. Our Equality check protocols can be used here.

SECURE SUM PROTOCOLS FOR INSECURE NETWORKS

Secure sum protocol permits calculating the sum of sensitive data multiple cooperating but distrustful parties without revealing their data to one another. Many Protocols are proposed by the researchers but there may be a presumption of secure network lines while devising the protocol. In this chapter secure sum algorithms applicable to insecure networks have been devised. We dropped the assumption of the secure network. To preserve the privacy of the data we used cryptographic techniques so that the actual value of the data is not visible to the user. The data provided by one party to another may be encrypted using symmetric or asymmetric encryption methods. The secure sum algorithms employing a physical ring structure faces a threat of colluding neighbours to capture the data of a middle party in the ring. To overcome this threat we devised a protocol which prevents colluding neighbours hacking the data of a middle victim. We consider developing algorithms for secure sum because all other arithmetic operations can be implemented using addition operation. Many practical situations arise when privacy of data becomes a concern. On the other hand knowing the result of common computation is in the mutual interest of the joint parties. Consider following scenario:

1. Four brothers living independently want to know the total wealth of family but no brother wants to disclose his individual wealth to others.
2. All the students in a class want to know the average marks obtained by students in a test but no student wish to show his marks to other students.
3. Group of a mobile service provider companies wants to know the total number of customers of all the companies in an area but no company wants to disclose its number to other companies in the group.

A protocol was proposed by Cliften et al. in 2002 [4] for computing the secure sum where the parties are set in a ring structure. One of the parties begins the protocol by adding random number to its data. The sum of all the private data plus random number reaches the originator party. The party subtract the random number and sends sum to all the parties. In this secure sum computation scheme the data flow in unencrypted form not preserving the privacy. Another threat is that to parties to a middle victim can share their data to capture middle party's data.

In this chapter we propose protocols which are also suitable for insecure networks. We used symmetric and asymmetric encryption on the data flowing on the network lines. It gives confidentiality to the sensitive data. There are three protocols proposed by us in this chapter.

A secure sum with symmetric key protocol uses a shared secret key by all the parties jointly performing the computation. An initiator party sends its data plus random number encrypted with the shared secret key to the next party in the structure. The receiving party decrypts the sum with the same key to get the data of sending party plus the random number. Now, it adds its data, encrypts with the key and sends encrypted sum to the next party. The next party does the same steps; receive, decrypt, add, encrypt and send. At the end the originator receives the sum of all the data plus random number encrypted with the key. The initiator party decrypts and subtracts the random number from the sum to get the actual sum of the data. In the second protocol secure sum with asymmetric key each party generates its public- private key pair with certain algorithm like RSA. Each succeeding party shares its public key with the previous party. The initiator party adds a random number to its data and encrypts with the public key of the next party to which it sends the encrypted sum. The receiving party decrypts with its private key to get the partial sum. The steps; receive, decrypt, add own data, encrypt with the public key of the next party, and send to the next party continues till the originator is reached. The originator gets sum plus random number at the end and therefore it subtracts the random number to get the actual sum. But the threat of the middle party being victim by its two colluding neighbours prevails in this method. The conspiring neighbours may share the partial sum and take difference of the partial sum to get the actual data of the middle party. To deal with such a threat we propose another method Secure sum against colluding neighbours the protocol initiator party generates publicprivate key pair and distributes its public key to all the joint parties. The initiator encrypts its data with its public key and sends to the next party. No random numbers are chosen. The next party simply adds its data to the received encrypted number, encrypts with the public key of the initiator and sends to the next party. This process is repeated until the protocol initiator is reached. At the initiator multiple decryption operation is to be performed to get the actual sum of the individual data. Since the private key is with the initiator only, no two parties can collude to get the data of the middle party

Secure sum with symmetric key

This method uses a shared secret key which is with all the parties. It also has a unanimously elected initiator party similar to Clifton et al. [4]. The encryption of the sensitive data with symmetric key provides privacy. Thus, the protocol is suitable for insecure networks.

Informal description of secure sum with symmetric key

The proposed protocol uses shared secret key which is assumed to be acquired by all the cooperating parties in advance. One of the parties is to initiate the protocol and it is assumed that there is a consent among parties that who initiates the protocol. The protocol initiator party chooses a random number (which is its secret), and adds to the private data with it. Then it encrypts the partial sum so obtained and sends to the adjacent party in the topology. The party receiving this partial sum decrypts it with the shared key to recover the partial sum sent by the initiator. Now, it adds its sensitive data to compute a new partial sum, encrypts it with the shared key, and sends to the next party. The next party will do the same steps. This process is repeated until the protocol initiator is reached. At this stage the initiator will recover the sum of all the private data and the random number. By subtracting the random number actual sum of all the data is obtained. This sum is broadcasted to all the parties. The encryption by the parties maintains the privacy of partial sum at the network lines. The random number maintains privacy of the data among parties.

Secure sum with asymmetric key

In the secure sum with symmetric key the security of the shared secret key is a great concern. When the key is compromised at any point in the network, the whole objective of the protocol becomes unachievable. The method in this section uses public key cryptography where each party generates a pair of public and private keys. Each of the parties shares its public key with the immediate preceding party in the ring topology. The parties send the private data by encrypting with the public key of the receiver. The receiver can decrypt the received encrypted data with its private key to recover the partial sum.

Informal Description of Secure sum with asymmetric key

All the parties wishing for joint secure sum computation has to generate publicprivate key pair. The succeeding party shares its public key to the previous party. There is consent among the cooperating parties that some party will act as a protocol initiator. The initiator chooses a random number, and adds to its private data. It now encrypts this partial sum with the public key of the next party, and sends the encrypted partial sum to the next party in the topology as shown in the Fig.3.5. The receiving party recovers the partial sum by decrypting with its private key. Now, it adds its private data to the recovered partial sum, and encrypts with the public key of the next party in the ring. It sends this encrypted partial to the next party in the ring. Now the receiving party performs the same steps, and sends the public-key encrypted partial sum to the next party. This process continues until the initiator party is reached. The

initiator party recovers the sum of all the data plus random number. It removes the random number and then broadcasts the result to all the parties.

Secure sum against colluding neighbours

The protocols described in this chapter so far face the problem of colluding neighbours who can share their partial sum to capture the private data of a middle party. Another limitation of the previous protocols is that all the parties perform encryption and decryption of the partial sum. To reduce this computation overhead and solve the problem of colluding neighbours we propose a new protocol called secure sum against colluding neighbours in which the parties except the initiator only perform encryption and not decryption. All the parties use the public key of the initiator party for encryption. The initiator decrypts multiple times to recover the secure sum.

The information security problem where multiple cooperating but distrustful parties wish to evaluate some function of their private data such that during joint computation their data is not revealed to other cooperating parties is called Secure Multiparty Computation (SMC). Secure sum computation is an instance of SMC where multiple parties compute sum of the sensitive data of all such that the result is correctly received and the privacy of the individual data is preserved. Many algorithms for secure sum computation is available in the literature. Many of them are applicable to secure network lines only. In this chapter we devised three novel secure sum protocols which can be used on insecure network lines. To maintain the privacy on the network lines we used symmetric and asymmetric encryption techniques. In first protocol secure sum with symmetric key we used symmetric encryption technique where the data is encrypted with a common secret key. But the problem of key distribution is critical here. When the key is compromised the whole network data will be hacked. In another protocol secure sum with asymmetric key all the parties generate their public – private key pair. The party sending data to other party encrypt the secret data with the public key of the recipient. Now the recipient decrypts it to recover the partial sum. In both the above techniques a party called initiator uses a random number that is added to the initiator data. The initiator recovers the sum and distributes to all. The third method uses single key pair of the initiator.

SECURE SUM PROTOCOLS USING HOMOMORPHIC AND THRESHOLD ENCRYPTION

Computing sum securely has been discussed in chapter 3. In secure sum multiple parties working on a joint task compute the sum of their individual data such that the result is known

to all the cooperating parties and the individual data privacy is preserved. The secure sum computation can work as a component of privacy-preserving data mining tool kit because addition is a basic operation in computer system. All other operations can be implemented using add operation. The protocols assuming secure networks are impractical as almost all network lines are insecure. In chapter 3 of this thesis symmetric and asymmetric encryption is used to protect privacy of the data flowing on the network lines. These protocols need huge number of encryptions and decryptions. In this chapter we use additive homomorphic properties of a cryptosystem to optimize number of encryptions and decryptions. The property ensures that the sum of the encrypted values is same as the encrypted sum of data. First protocol of this chapter named as Secure homomorphic sum uses additive homomorphic property of a cryptosystem as devised by Paillier in 1999 [5]. In another protocol for secure sum named as Secure threshold sum we apply Shamir's threshold scheme together with the additive homomorphic encryption. The threshold scheme shows that a piece of information can be reconstructed from minimum k number of pieces out of total n pieces of information. Less than k pieces are not sufficient to reconstruct the original information. In both the cases our protocol is suitable for semi-honest adversary who follow the steps honestly as listed in the protocol but may perform some extra steps to learn private values with parties. We use ideal model of SMC in Secure homomorphic sum where there is no third party involved during computation. We coin the term semi-ideal SMC model in case of Secure threshold sum as it is the hybrid of the ideal and real model of SMC.

Informal Description of Secure asymmetric sum protocol

The protocol secure asymmetric sum works in similar way as the protocol secure random sum because both use similar architecture of ideal SMC model. The difference is that the secure random sum use random numbers for privacy-preservation and the protocol secure asymmetric sum uses cryptographic approach using additive homomorphic cryptosystem. A protocol initiator party, normally P_0 generate public private key pair (PU, PR) and distribute to all the participating parties. It is assumed in our protocol that this distribution is achieved by some secure means in advance. All the parties encrypt their private data with the private key and send the ciphertext to the TP. The TP takes sum of all such ciphertexts and sends back the computed sum to the participating parties. Owing to the additive homomorphic properties the sum of ciphertext is equal to the ciphertext of the sum of the private data. The participating parties decrypt the sum to get the sum of the private data. The private key is with the participating parties only. Thus the TP doesn't know anything about the actual data.

The parties do not communicate any data with one another. Therefore the privacy of the data is preserved from one another. In this chapter we devise two protocols to get the sum of the private data of certain number of joint but distrustful parties. The problem is called secure sum problem in the literature. It is an instance of more general problem of information security called Secure Multiparty Computation or SMC. In SMC multiple parties compute some function of their private data without disclosing the individual data to one another. The first protocol in this chapter Secure Homomorphic sum allows cooperating parties to compute sum of their private data using additive homomorphic properties. It uses a ring topology in which parties are supposed to be arranged. One of the parties is unanimously chosen as a protocol initiator. The party generates its public/private key pair, distribute the public key to all, and helps in computing the sum of encrypted values. This sum is equivalent to the ciphertext of the sum of the data. By decrypting with the private key, the initiator gets the secure sum. In the second protocol, Secure threshold sum of this chapter additive homomorphic property along with the threshold encryption scheme is used to get the secure sum. A novel SMC model which removes the dependence on the trustworthiness of the TTP is proposed. The TTP generates public-private key pair using additive homomorphic property and shares of the private key using threshold encryption scheme. The protocol is improved in many ways. It is applicable to insecure networks and semihonest adversary. Another secure sum protocol, secure random sum in ideal model is proposed where a TP helps the participating parties computing the sum of private data. A common random number is used to protect the privacy of the data from other parties as well as from the TP. In the last protocol of this chapter, secure asymmetric sum protocol, the architecture of secure random sum is use. But instead of using the random number, a public/private key pair using homomorphic cryptosystem is used for privacy preservation. The sensitive data is encrypted by the public key and then sent to the TP the TP computes the sum of these ciphertexts and sends sum to the parties. The parties recover the sum by decrypting it with the private key. The parties do not communicate directly. Therefore the privacy of the data is preserved among them. The TP receives data encrypted with the public key. As TP doesn't possess the private key, it cannot recover the data. The protocol is suitable for semi-honest adversary. Future work can be extended to devise the secure sum protocols for malicious adversaries. Zero Knowledge proof concepts will be helpful in this case. Our protocols are privacy preserving protocols. Protocols can be devised which protect other security properties like integrity and nonrepudiation.

SECURE MULTIPARTY EQUALITY CHECK PROTOCOLS

Much has been discussed in the previous chapters about the information security problem SMC. Many times the parties simply desire to check their individual data for equality but they are worried about the privacy of their data. During privacy-preserving data mining operation the sensitive data may be present at multiple sites who are involved in computation. Parties need results but they are afraid of the privacy of their data. In this chapter we have proposed a protocol which allows multiple joint parties to check their data for equality without revealing individual data to one another. We have used randomisation technique as well as the secure additive homomorphic property of the encryption to get the correct result. The protocol is suitable for semi-honest adversary.

Secure MultiEqualityCheck protocol

In our proposed protocol we have extended the two-party EqualityCheck protocol to protocol which can perform equality check of the data of more than two parties. We have proposed architecture for multiparty equality check. We also provide informal and formal description of the proposed MultiEqualityCheck algorithm. We also analyse it for its performance.

Informal description of Secure MultiEqualityCheck protocol

Referring to the Fig.5.3, multiple parties P_0 to P_{k-1} run MultiEqualityCheck by choosing unanimously one party as initiator. Normally, P_0 is chosen as an initiator. All the parties generate their public-private key pair using additive homomorphic cryptosystem. Initially, the party P_0 and P_1 run EqualityCheck algorithm to check equality of their data. If the equality holds, then the party P_1 and P_2 run EqualityCheck algorithm. In this way equality is checked pair-by-pair. If equality continues to hold till the initiator is reached, the result that all the data are equal is declared. If equality breaks at any point in the ring the result that the data are not equal is declared.

The privacy preserving function evaluation is the need of today's world as the sensitive data from multiple sources may be shared by multiple sites for joint computation. Secure equality check is one of the cases where a set of parties need to compute the data for equality. For example two police stations situated far apart want to check some biometric data of a criminal such that both the stations do not want to disclose actual data to one another. In this case two-party equality check would be useful. In our work we have extended the two-party equality check to multiparty case. Our protocol secure MultiEqualityCheck uses additive homomorphic cryptosystem where the sum of ciphertext is equal to the ciphertext of the sum. Our work is suitable for semi-honest adversaries. The protocol MultiEqualityCheck uses

public-private key generation using homomorphic encryption, encryption, and decryption. We proposed a novel equality check scheme using hash function. The protocol EqualityHashCheck ensures checking equality among multiple parties with the help of comparison of hash functions. It has an advantage of less computation as compared to the MultiEqualityCheck. We have proposed EqualityHashCheck for both real and ideal model of SMC. In real model the parties compare their hashes pair-by-pair until all the parties are traversed. In the ideal model the hashes are provided to a TP. It is the TP which checks whether all the hashes are equal and provides result to all the parties. The privacy is preserved from TP as it receives hash of data not the actual data. The hash function is an irreversible function. That mean given the hash, it is infeasible to compute data from the hash. Similarly, privacy is preserved among parties as they do not communicate with each other. But, the cost of maintaining the TP is an overhead.

LITERATURE SURVEY

when multiple parties with their private data want to compute some common function of their data jointly such that the privacy of their data is preserved from one another, the problem is called Secure Multiparty Computation (SMC) [11]. In this scenario all the participating parties do not have trust on one another. But the common function evaluation is in their common interest. Mainly, there are two goals of SMC problem, privacy and correctness of the result. Owing to the heavy use of the Internet which is associated with the distributed system, the scenario of SMC has become highly relevant. The parties geographically distributed on distributed sites may wish to evaluate a function providing their data. But at the same time they are worried about privacy of their data. During online transaction many sites cooperatively work to process the transactions. But each site does not necessarily want to show the exact value of the data it shares with other party. Therefore in this age of large number of online transactions the SMC is has become highly relevant. Parties frequently need to perform joint computation on their sensitive data while keeping confidentiality of the data. Consider a situation where two or banks wish to know cooperatively the details of a customer from their individual databases. Knowing the details about the customer is in mutual interest of all the participating banks. But no want wants to share their database to other banks as they are competitor in the same sector.

The development of networking, the Internet, and distributed system has provided huge opportunities for joint computations where multiple parties perform SMC on their secret data. Due to the concern of the privacy and worldwide regulations made by made by different

countries, there exist many scenarios where SMC becomes applicable. Consider following scenarios which will let you understand the practical applicability of the SMC solutions:

1. A person having his DNA pattern wants to about the genetic diseases associated with his DNA pattern. He wants to do some query from a server storing different DNA patters and the diseases associated with those DNA patterns. But at the same time the person does not want to disclose his exact DNA patter to the server. This privacy-preserving database query can be implemented using SMC solutions.
2. A group of mobile service providers wish to compute together to prepare list of total subscribers active in an area within some specified time interval. This could be due to some police investigation and asked by the intelligence of a country. The companies do not want to disclose their customer details. This is a case of privacy-preserving union of sensitive databases and can be solved using SMC solutions.
3. A group of students wish to know their average marks obtained in an examination but know student wish to disclose his actual marks to other students of the group. This is a case of privacy preserving sum computation which is called as secure sum in the literature. The sum in this example is divided by number of students to get the secure average.
4. Five real brothers who live separately wish to compute their total wealth but know brother wish reveal his wealth to other brothers. This problem can also be solved using secure sum solutions.
5. A bank wants to some loan details of a suspicious customer from other bank but the bank don't want to disclose actual customer details to other bank. This privacy-preserving query can be solved using SMC solutions.
6. A group of police stations in a country wish to search details of a criminal from their databases but no police station wish to show its database to other police station. This is a case of privacy preserving data mining.
7. A client computer in a payment system wish to learn that a QR code is matching or not without showing the actual QR code to the server. This is a case of privacy-preserving matching or equality check.

Formally, consider multiple parties P_0 to P_{k-1} with private data d_0 to d_{k-1} respectively. These k parties want to evaluate function $f(d_0, \dots, d_{k-1})$ without revealing their private data to each other. Based on the type of f many specific SMC problems are devised by researchers and thereafter many real life applications emerged. We pointed out these works in our publications . This chapter will explore SMC and will put forward the research gap.

SMC Models

The SMC problem needs paradigms for performance analysis and mathematical modelling. The participating parties with private data are the essential stake of the SMC paradigm. But sometimes these parties seek assistance of a Trusted Third Party (TTP) for function evaluation. Based on the presence and absence of the TTP, two SMC paradigms evolved in the literature:

1. Ideal SMC Model

2. Real SMC Model

Ideal SMC Model

In the Ideal SMC Model the cooperating parties take the services of a TTP for evaluating the function of their data. In ideal case the TTP should be fully honest as it does not share any sensitive information of one party with any other party. But researchers in the literature analysed every kind of behaviour of the TTP. Due to presence of the TTP it becomes easier for the participating parties to preserve privacy of their data from one another as they will share least information with one another. If the TTP behaves maliciously the whole objective of the SMC becomes unachievable. The architecture of the Ideal SMC model is depicted in the Fig. 2.1. The Ideal Model is easier to implement. But there are two drawbacks of this model. One, the cost of maintaining the TTP, and second its trustworthiness. But it is the model is more popular in practice as the TTP is a government agency or an organisation approved by the government. In case of any dispute among the participating parties, the TTP plays an important role to reach at the decision.

SMC SOLUTION TECHNIQUES

When the researchers paid attention to the requirement of privacy-preservation during computation, they started providing solutions using combinational logic circuits. The inputs to these circuits were the sensitive data and the results were obtained at the outputs. As an example, consider a case where two parties wish to check the comparative magnitudes of their data without sharing actual magnitudes with one another. A combinational logic circuit named digital comparator can be useful. It accepts two binary numbers at the input, and provides the result at the output. The comparator can be treated as the TTP of the Ideal SMC Model. The scheme is shown in the Fig.2.3 [14, 15]. The result is obtained over three output lines; E, G, and L. If $E = 1$, both the numbers are equal. If $G = 1$, first number is greater than the second number. If $L = 1$, the first number is less than the second number. Assuming the input lines to be secure, the result is obtained at the output preserving the privacy of two

parties' data from one another. But, for even two-party case the circuit is complex and expensive. For multiparty the complexity and the cost reaches so high as not practically affordable. Soon the researchers realised to use some other economical method which can be practically used. In this connection the work of Yao in 1982 [1] used cryptographic technique to provide solution to millionaires problem. The problem allows two millionaires to know who is richer without sharing actual wealth to one another in Real SMC Model. Yao used symmetric encryption techniques to achieve the result. Yao's two-party SMC problem was extended to multiparty SMC problem by some researchers. For example Goldreich et al. [16] proposed an algorithm for multiple parties to compute some function and leaking incomplete information. But this is suitable when majority of parties were honest. Generalised solutions for multiple SMC problems are proposed by the researchers but these solutions are inefficient. Goldreich et al. [18] showed that specific solutions for a particular SMC problem are more efficient than these general solutions. Presently, there are three techniques which are used to provide solutions to SMC problems:

1. Randomization techniques
2. Cryptographic techniques
3. Anonymization techniques

Randomization SMC techniques

In randomization approach participating parties use random numbers to preserve privacy during computation. The protocols are designed such that the presence of the random number does not influence the result of the computation. But the random number is used to hide the actual value of the data from other parties involved in the computation. Many of the programming languages available today have in-built random function that can be used in these protocols. To understand the randomization technique we explain the secure sum algorithm presented by Clifton et al. [4] where they propose an algorithm for computing sum of individual data from multiple parties such that their privacy is preserved. This is the most easily understood SMC solution using random number. They proposed the protocol by arranging the cooperating parties in a logical ring structure. This is a real SMC model as there is no TTP. All the party unanimously agree on a party which initiates the protocol. The protocol initiator party first chooses a secret random number. It adds the random number to its private data, and sends the partial sum to the next party in the ring. The network lines are assumed to be secure, and therefore it is assume that no intruder can intercept the network lines to capture the data. There is always a unidirectional communication in the logical ring.

The party receiving the partial sum adds its private data to the partial sum and gets new partial sum. This partial sum is sent to the next party in the logical ring structure. This process is repeated until the protocol initiator party is reached. At this point the protocol initiator party gets sum of all the data plus random number. The initiator subtracts random number to get the sum of actual data of all participating parties. The sum is then broadcasted to all the participating parties.

Cryptographic SMC techniques

The cryptographic approach for providing solutions to SMC problems uses symmetric and asymmetric encryption techniques. In symmetric encryption method a common key is used for encryption as well as decryption. But an asymmetric encryption method uses different keys for encryption and decryption. The second method is also known as public key cryptography. Different researchers used their approach and the literature of SMC has many building blocks which can be used to solve SMC problems [19]. These building blocks are listed below:

1. Yao's Millionaires Problem.
2. Homomorphic Cryptosystem.
3. Oblivious Transfer.
4. Private Matching.

MOTIVATION OF THIS WORK

Although the secure sum protocol proposed by Clifton et al. [4] is a milestone in secure arithmetic computation. It has many limitations when we apply it to for the real life applications. One, as mentioned earlier is the threat of two colluding neighbours getting a middle party hacked. This drawback can be removed by many ways. The colluding neighbours simply take difference of what they send and what they receive. There is an immense need to provide protocols which deal with this kind of threat. This is because the secure sum is an integral part of privacy preserving distributed data mining. In our earlier work, we proposed segmentation approach for secure sum computation where the sensitive data of the participating parties is broken into segments. The parties then get the sum by adding segments in different rounds. It also eliminates random number and the threat of getting the middle party victim. The protocol we named as k-secure sum protocol [56]. It has been analysed that the protocol provided better security against data leak. In another version of a protocol based on segmentation approach we suggested distributed the segments among parties before the sum computation. The protocol was called as distributed k-secure sum

protocol [57]. In another protocol we proposed an architecture where each party must change its position after one each round of segment computation, so the two neighbours could only hack segments not full data. The protocol was named as changing neighbours k-secure sum protocol [58]. In thesis we used cryptographic approach with random numbers to deal with the problem of colluding neighbours. Another problem with the secure sum protocol was that it assumed the network to be secure. We dropped this assumption and devised protocols for insecure networks. Our protocols are based on symmetric cryptography, asymmetric cryptography, homomorphic cryptosystem, and threshold encryption. We also devised protocols for equality check among multiple parties.

SMC is an interesting paradigm of computing where multiple parties cooperative evaluate some function of their common interest without revealing actual data to one another. The milestone for this problem is a millionaires' problem proposed by Yao in 1982 [1] where two millionaires can cooperatively work to determine who is richer without disclosing exact wealth to one another. It has been predicted that due to privacy concerns soon SMC may become an integral part of our computing environment. The solutions to SMC problems earlier based on the circuit evaluation. Modern SMC solutions use randomization method, cryptographic methods, and anonymization techniques. Many specific SMC problems and their solutions are suggested by the researchers. Based on these solutions many real life applications are developed. The primary goal of the SMC is to get the correct result while preserving the privacy of the data involved in the computation. Researchers proposed SMC problems and solutions with an objective to reduce communication and computation complexities. In this chapter we explored the subject of SMC and various techniques used by researchers, and specific as well as real life applications provided by them. We also pointed out our earlier research contribution to the subject. We are able to find the research gap between earlier work and the work needed. Earlier protocols assumed the network lines to be secure. Therefore we have developed protocols for insecure networks. Our secure sum protocols work appropriate on insecure lines. Similarly, we have extended two-party equality check protocols to multiparty equality check protocols. These are also suitable for insecure networks and semihonest adversaries.

CONCLUSION AND FUTURE SCOPE

A Voyage to Secure Multiparty Computation

Secure Multiparty Computation is a problem of information security where multiple joint parties cooperatively compute some function of their private data without revealing the actual

data to each other. In modern distributed computing environment the parties may be geographically distributed at different sites. These sites will have to provide data for computation purpose. But all these sites are concerned about the privacy of their data. SMC solutions provide this privacy. The millionaires' problem proposed by Yao [1] is considered to be the milestone to modern SMC era. The problem provides solution as how two millionaires can determine together who is richer without disclosing actual wealth to one another. He provided cryptographic solution to the problem. This two-party problem was extended to multiparty case by many researchers. As of today, there are three techniques which are used to provide SMC solutions; Randomization, cryptographic, and anonymization methods. The randomization technique uses random numbers to hide the actual data during computation. The random number is removed while declaring the result. Secure sum protocol of Clifton et al. [4] is an easier example of randomization approach. In cryptographic techniques, the data is encrypted during computation so that its actual value is not revealed to other parties. Decryption operation is performed while declaring the result. This technique may exploit homomorphic property of the cryptosystem which ensures that sum of the ciphertext is equal to the ciphertext of sum of data. Another important component of the cryptographic technique could be an oblivious transfer where a sender has many inputs and a receiver needs one of them. The receiver sends the index for the input selection and the corresponding input is received by the receiver. But the sender is not aware about the index value and the receiver is not aware about the exact position of the received input. In anonymization method the identity of the participating parties is hidden by anonymizers. Therefore the privacy of data will be maintained during the computation. The architecture of the SMC solution has two flavours in the literature; Ideal SMC model and Real SMC model. In an Ideal SMC model there exists a third party among the participating parties which assists the secure function evaluation. This model is easier to implement. But the trustworthiness of the third party is to be taken into consideration while devising the protocols. In Real SMC model there is no third party and the participating parties only run the protocol to compute the function securely. This model is difficult to implement but is less expensive due to absence of third party. In this thesis we proposed a third model which is a hybrid of above two models.

Secure Sum Protocols and their performance

Secure sum computation is an instance of the generalised SMC problem where the cooperating parties compute sum keeping their privacy. The work of Clifton et al. [4] is considered to be milestone where they used a random number to compute the sum of private

data of all the parties securely. We also focussed on secure sum as all other arithmetic operations can be implemented using addition operation. There are two limitations of secure sum protocol; first it assumes network lines to be secure and second the problem of colluding neighbours in the architecture. But in actual practice the network lines are always insecure. Therefore in this thesis we devised secure sum protocols for insecure networks. They used real SMC model. We proposed secure sum protocols for both the SMC models; real as well as ideal SMC model. We also devised protocol for our novel proposed semi-ideal model. The secure sum protocols proposed in this thesis use symmetric and asymmetric cryptography for making them secure on the insecure network lines. Our secure sum with symmetric key protocol [7] uses a common shared secret key by all the participating parties. The parties also use random numbers similar to Clifton's secure sum protocol for preserving privacy among them. The parties first encrypt the partial sum before sending it to the next party. Therefore the confidentiality is maintained on the insecure network lines. Parties decrypt the received encrypted partial sum and add their private data, encrypt again, and send to next party. Although, the privacy is achieved among parties themselves and on the insecure network lines but the computation involved in the process of decrypt-add-encrypt is very large. Another problem with this is sharing the symmetric key over insecure network lines is also an issue. When the shared key is compromised the whole scheme fails. In another protocol proposed by us in thesis named secure sum with asymmetric key protocol [7] we suggest an architecture where each of the parties generates its public private key pair. Each party shares its public key with the previous party in the ring. The party adds and encrypts the partial sum with the public key of the recipient. The recipient decrypts the encrypted sum with its secret private key. Rest process is similar with secure sum with symmetric key protocol. This method solves problem of key distribution among parties. Both of these protocols suffer from the problem of colluding neighbours. To remove this problem we proposed another protocol named as Secure sum against colluding neighbours protocol[7]. It uses multiple encryptions and decryptions for privacy. The protocols discussed so far employ huge computations for encryption and decryption. In order to reduce amount of computation we used additive homomorphic property of a cryptosystem [5]. This property ensures that the sum of encrypted data is equal to encrypted value of the sum. Therefore irrespective of the number of participants the decryption needs to be done only once. Based on this idea we proposed Secure homomorphic sum protocol [8]. We suggest a ring architecture in which the parties communicate in one direction. An initiator party generates its public private key pair using

additive homomorphic cryptosystem [5]. The party shares its public key with all the parties. All parties encrypt the private data with the public key and send the partial sum to the next party. This process is repeated until the initiator party is reached. The initiator decrypts with its secret private key to recover the secure sum which is distributed to all the parties. Although the amount of computation is much reduced due to the single decryption operation, the protocol faces non-uniform privilege of initiator party generating, distributing public key, and holding private key a secret. In our other protocol services of a third party are used for generating public-private key pair, distributing private key to all the parties. The private key is not held by any single party. Instead we used Shamir's secret sharing concept of threshold encryption [6] where it has been shown that a value can be reconstructed out of certain minimum number of shares of that value. Therefore we suggest the third party generate number of shares of the private key and then distribute one share to each of the parties. The parties can recover the private key only if all provide their share honestly. The protocol is named as Secure threshold sum protocol. The secure sum protocols suggested in [4, 56, 57, 58] and many other researchers are applicable to real SMC model. As pointed out real models are difficult to implement, we have proposed protocols for ideal model where the services of a third party are used. The protocol is named as Secure random sum and is described in chapter 4. The participating parties agree on a common random number which is multiplied to the private data. The products so computed are sent to the third party. The third party computes sum of all such products and sends the sum to all the parties. These parties divide the sum by their secret random number to get the secure sum. The protocol is simple with minimum computation complexity. But sharing of the random number on insecure network lines is an issue. In order to solve the problem of random number distribution we have proposed another secure sum protocol for ideal model named as secure asymmetric sum protocol in ideal model. It is a slightly modified version of Secure random sum protocol. Instead of using random numbers, all the parties generate public private key pair. The data is encrypted with the public key and sent to the third party.

Secure Equality Check Protocols and their performance

The problem of checking equality between two parties has been proposed in [59] which used homomorphic encryption to check equality of two parties in real model. We have extended the problem to multiparty equality check. Our protocol MultiEqualityCheck [10] allows multiple parties to check equality of their data without disclosing their private data. It is based on additive homomorphic public key cryptosystem. It uses real SMC model and suitable for

semi-honest adversaries. It is also applicable to insecure networks. The MultiEqualityCheck protocol use key generation, encryption, and decryption at each party. We also proposed protocol for equality check which simply uses hash function. Each of the parties takes hash of their data and compare these hashes cooperatively. Based on these we developed novel protocol EqualityHashCheck for both real and ideal SMC model. There are obvious benefits of the EqualityHashCheck over MultiEqualityCheck protocol. It does not need any key and random number. Its amount of computation needed is very less. Communication complexity is also very low.

Future Scope

we have proposed novel secure sum protocols and secure equality check protocols. The protocols we propose in this thesis are suitable semi-honest adversaries. These adversaries should obey all the steps described in the protocol but they may also try to learn some private data of other parties. The work can be extended for devising the protocols for malicious adversaries. These adversaries neither follow the steps of the protocol nor remain inactive when it comes to hacking the data of cooperating parties. The protocols when designed for the malicious or corrupt parties must be able to check the inputs for correctness. This is a big deal. But techniques are available in the literature for designing such protocols. There are zero knowledge proof techniques which help no whether an adversary has knowledge of the data and whether the data is correct. These protocols are complex and expensive. But they are robust and need of real life applications. Thus, there is a trade off between its efficiency and cost which needs to be considered while devising such protocols. In this thesis we have proposed some protocols for real model, some for ideal model, and some protocols for both the models. Suitable models can be used to our protocols based on the need. There are many aspects of the network security like privacy, authentication, integrity, and nonrepudiation. Our protocols are designed to deal with the first one, privacy. All these protocols can be designed to take care of other aspects of the network security. Finally, we deigned protocols for scalar values. Protocols can be designed which are applicable for matrix data.

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Studies on Edible Aroid Genera' Botanical and Genetic Diversity

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Abstract

One of the largest monocot families, the Araceae Juss. (Aroideae) contains roughly 105 genera and 3300 species (Mayo et al. 1997). With a few in the temperate zone, they are primarily found in tropical and subtropical regions. Members of the Araceae family are typically scrambling shrubs or climbers with aerial roots. Some grow as herbs and produce corms or tubers. Rarely are they real epiphytes or freely swaying aquatic plants. The Araceae are frequently referred to as the aroid or arum family. Five genera are represented by the cultivated edible aroid tuber crops. They are *Amorphophallus* Blume et Decene, *Cyrtosperma* Griffith, *Colocasia* Schott, *Alocasia* (Schott) G. Don, and *Xanthosoma* Schott. Colocasioideae, Lasioideae, and Aroideae are the three subfamilies to which they belong (Nicolson 1982; Bogner and Nicolson 1991; Mabberly 2005). The genus *Cyrtosperma* belongs to the Lasioideae, Colocasioideae, *Alocasia*, *Colocasia*, and *Xanthosoma* subfamilies, and Aroideae subfamilies of *Amorphophallus*. There are 70 species in the genus *Alocasia* (Mabberly 2005), 200 species in the genus *Amorphophallus* (Jaleel 2002), 8 species in the genus *Colocasia*, 11 species in the genus *Cyrtosperma*, and 57 species in the genus *Xanthosoma*. *Amorphophallus paeoniifolius* (*A. campanulatus*), *Colocasia esculenta*, *Cyrtosperma chamissionis*, *Xanthosoma nigrum*, and *Xanthosoma sagittifolium* are among the most commonly grown edible species.

INTRODUCTION

The category of cultivated plants known as tropical tuber crops is largely ignored. After grains, tuber crops are the most significant human food crops. They serve as a supplemental food source for sizable populations in the tropics and subtropics. According to the FAO (1988), they serve as the primary staple for 5% of the population and a secondary staple for the majority of the rest. The aroids and yams are the most underutilized tuber crops. In South

and Southeast Asia, as well as Oceania, aroid tuber crops including *Alocasia* and *Amorphophallus* are grown as vegetable crops. They are merely incidental in other tropical regions. Even entire information about their region and manufacture is unavailable. Cocoyams are a combination of *Colocasia* (taro) and *Xanthosoma* (yautia, tannia). They are significant crops in tropical Africa, South and Southeast Asia, and Oceania. America's tropics. In various South Pacific Ocean countries, *Cyrtosperma* is commonly grown. islands. Taro was cultivated in an area of 140163 6 hectares in 1999, producing 9652935 t (FAO 1999). 2009 saw an increase with 11312073 t of production from an 1576340ha in size (FAO 2009). This demonstrates that area and production have grown by 12.46% and 17.19%, respectively, from 1999 and 2009. However, these numbers excludes information for a number of nations, notably India, where it is crucial seasonal vegetable in Sri Lanka, Bangladesh, and the east and northeast of India. Only four and seven countries, 20 in Africa, eight in America, nine in Asia, and In Oceania, ten. With 4459650 t of output, Nigeria is the top producer in 2009. a 629787 hectare area. China is in second place with a score just below Nigeria. output of 92697 hectares yielding 1692551 t. In 1999, tannia or yautia was grown on an area of 33649 ha, producing 301211 t (F AO 1999). This grew in 2009 from an area of 49556 ha with 394045 t (FAO 2009). This demonstrates that from 1999 to 2009, area and production both increased by 47.27% and 30.82%, respectively. There are just 12 countries in the American continents included. Cuba produces the most in the world, with a total output of 199400 t from a 27027 ha area. With a yield of 80,000 t from 7000 ha, Venezuela comes in second. However, this information excludes nations like India, where it is a significant seasonal vegetable.

Aroids are eaten as a primary, secondary, or vegetable staple. By fresh weight, tubers are primarily composed of around 75% water, 20% carbs, and about 2% proteins (Onwueme and Charles 1994). On a dry weight basis, they are as nutritious as, and in some cases better than, rice or wheat (Rao et al. 1989). Up to the conclusion of World War II, aroids were the main source of nutrition for the populations of Oceania, particularly Polynesia, Central America, including the Caribbean, and central and west Africa. According to the International Food Policy Research Institute (IFPRI), from the current levels of 5% of food demands, consumption of tuber crops will increase by 20%.

Washington DC recommended in 1995 that by the year 2020 AD, the output and consumption of tuber crops would increase by 20% from the present levels of 5% of food needs. The genus *Alocasia* represents 10 species, *Amorphophallus* 21 species, *Colocasia* 5

species, *Cyrtosperma* 1 species and *Xanthosoma* 2 species in India (CSIR; Wealth of India 1976; Karthikeyan et.al., 1989; Sivadasan 1986, 1989; Sivadasan et al 1994; Hetterscheid et al., 1994; Hetterscheid and Ittenbach 1996; Jaleel 2002). Among these, *Colocasia esculenta* (taro) is the most widely grown species in India. In India, it can also be found in the wild. Additionally grown in India are *Alocasia macrorrhizos* and *Amorphophallus, paeoniifolius*, as well as their wild counterparts. The *Xanthosoma* species are imported from tropical America for cultivation.

One of the major hindrance for the acceptance of the tubers and leafy parts of edible members of the family Araceae as a regular component of the diet is perhaps the presence of calcium oxalate crystals and acidity. Since the aroids are hardy and can be cultivated with relative ease on a wide variety of terrain it has great potential to become important carbohydrate source in the tropical parts of the globe. To fulfill this potential it is imperative to establish the cause of acidity and to verify whether there is any correlation between the presence of calcium oxalate crystals and acidity. It has also to be ascertained whether hybridization can be used as an effective tool to decrease the adverse traits such as acidity and to increase the nutrient content in the cultivable types of aroids. The rate and pattern of growth and morphogenesis in these plants has also to be understood comprehensively before employing them for agricultural purposes. The study of pollination and the factors involved in the pollinations process such as thermogenesis is also crucial in this respect. The present study was carried out with a view to collect, catalogue and characterizes the taxa that are widely prevalent and obtain a better understanding of their botany and genetics. The study carried out includes botanical (growth and development, flowering, pollination biology, thennogenesis, acidity etc.) and genetical (selfing, inter-specific and inter-generic hybridizations, isozyme studies, morphological characterization etc.). Even basic information on aspects like growth and development, thermogenesis acidity etc is lacking. The information from the present studies is used to characterize the diversity and interrelationships of the above genera especially to the commonly cultivated species.

Growth and development

Growth and pattern of development were studied in the species of the genera *Alocasia*, *Amorphophallus*, *Colocasia* and *Xanthosoma*. Growth and development of leaves, corm, cormels and stolons was studied. Corm cuttings were grown in garden pots and in the field.

Observations on leaf development, opening and its longevity in different species and their accessions were studied. The neighbouring wild populations were also observed.

Morphology

The morphological data on aerial and underground parts were recorded as per IPGRI descriptor for taro (1999). The quantitative and qualitative characters were converted into numbers and tabulated. The data was used to study morphological variations. The cladogram was prepared by using the STATISTICA software.

Twenty characters used to describe the morphology in giant taro, taro and tannia are listed below.

1. Plant height (1- dwarf (< 50 cm); 2-medium (50-100 cm) and 3- tall (> 100 cm).
2. Number of stolons (0- None, 1-1-5; 2- 6-10; 3- 11- 20; 4- > 20).
3. Position of leaf lamina (1- drooping, 2-horizontal 3- cup shaped, 4- erect apex up; 5- erect apex down).
4. Leaf blade margin (1- entire; 2 undulate and 3- sinulate).
5. Leaf blade colour (1- whitish; 2- yellow or yellow green; 3- Green; 4- dark green; 5- pink; 6- red ;7- purple and 8- blackish or violet black).
6. Leaf blade margin colour (1- whitish; 2- yellow; 3- orange; 4- green; 5- pink; 6- red and 7- purple).
7. Lamina length width ratio
8. Petiole junction colour.
9. Petiole/lamina length ratio
10. Petiole colour (0- absent; 1- whitish; 2- yellow; 3- orange; 4-light green; 5- green; 6- red; 7- brown and 8- purple).
11. Petiole stripe (0- absent and 1- present).
12. Sheath length / petiole length ratio.
13. Sheath colour (1- whitish; 2- yellow; 3- light green and 4-red purple).
14. Flower formation (0- absent; 1- rare and 3- flowering).
15. Fruit formation (0-absent; 1 — present, 2-rare).
16. Number of spadix in leaf axils.
17. Corm length in cm.

18. Corm shape (1- conical; 2- round; 3- cylindrical; 4-elliptical; 5 dumb-bell; 6-elongated).

19. Number of cormels.

20. Corm flesh colour (1- white; 2-yellow; 3-orange; 4- pink; 5-red; 6-red-purple; 7-purple).

The elephant yam is different in morphology when compared to giant taro, taro and tannia. Hence the following 11 characters were used to describe the morphology.

Plant height(1- dwarf(< 50 cm); 2-medium(50-100 cm) ad 3- tall(> 100 cm).

2.Petiole colour(1- greenish black; 2 black 3-light geen; 4-brown; 5-geen; 6-purple; 7-whitish).

3. Murication(0-absent; I-present and 3-strongly murcate).

4. Mottles(0-absent; 1- present).

5. Mottle colour(I-white; 2-greenish white; 3- purple).

6. Petiole /lamina length ratio.

7. Flower fonnation(0-absent, 1- present)

8. Fruit fnnation(0-absent, I-present)

9. Conn manifestation(0-absent; I-absent).

10. Number of cormels

11. Corm flesh colour(1- pale yellow; 2- Brown; 3-Puple; 4-white).

Flowering and pollination biology

The regular blossoming of all the accessions bearing flowers was noted. The A macrorrhizos (giant taro) accessions all lacked flowers, with the exception of one. The enormous taro was intended to flower by spraying GA3 from the air (Wilson and Cable 1983). The results show that various treatments were applied to the aroid inflorescence to examine their effects on various aspects of flowering pollination and seed establishment. This largely consisted of bagging experiments and the elimination of different flower parts. Floral morphology

Floral morphology was studied in all the flowered accessions. Observations on different floral parts were done to study variations in floral characteristics. This was done at the time of opening of the flowers.

Pollen frtility

In the present study, pollen stainability has been taken as the index of pollen fertility. For studying pollen stainability, 1% aqueous solution of iodine was used. Small quantities of pollen were transferred to a slide; one drop of iodine solution was added to it and mixed well. A cover slip was placed over it and examined under low power of a compound microscope.

Pollen germination

This study was conducted for studying the nature and period of viability in *Alocasia*, *Amorphophallus*, *Colocasia* and *Xanthosoma*. This was done based on the methodology reported by Jos et al. (1967). The pollen grains were collected at the time of pollen shedding and its viability was tested at regular intervals in vitro. The medium was made with 2% sucrose and 0.02% boric acid. Culturing and germination were carried out at regular intervals.

Stigma receptivity

Alocasia odora and *Xanthosoma sagittifolium* were used in this investigation. By using artificial pollination before and after flowering (inflorescence opening), the duration of stigma receptivity was examined. Emasculated, packaged, and labeled inflorescences. The lower spathe was sliced longitudinally after the upper spathe was removed. Flowers with pistillates were visible. Pollen was used to artificially pollinate the flowers. Using a little camel hair brush, separate from other plants. It was labeled and bagged. beginning of the inflorescence's opening to 24 hours later (i.e., until the moment of pollen spilling). Records of the observations were made. In *A. odora*, pollination took place between 126 hours before and 55 hours after the flower opened. Pollinations in *X sagittifolium* took place after 144 hours.

Flowering and pollination biology

Observations were made on natural flowering in all the flowered accessions. The accessions of *A macrorrhizos* (giant taro) except one were found non-flowering. Attempts were done to induce flowering in giant taro by foliar spraying of GA3 (Wilson and Cable 1983). A series of manipulations were done on the aroid inflorescence as listed in the result to study their effect on different aspects of flowering pollination and seed setting. This involved mainly bagging experiments and removal of different floral parts.

Floral morphology

Floral morphology was studied in all the flowered accessions. Observations on different floral parts were done to study variations in floral characteristics. This was done at the time of opening of the flowers.

Pollen fertility

In the present study, pollen stainability has been taken as the index of pollen fertility. For studying pollen stainability, 1% aqueous solution of iodine was used. Small quantities of pollen were transferred to a slide; one drop of iodine solution was added to it and mixed well. A cover slip was placed over it and examined under low power of a compound microscope.

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Stigma receptivity

This study was studied in *Alocasia odora* and *Xanthosoma sagittifolium*. The period of stigma receptivity was studied by artificial pollination before and after flowering (inflorescence opening). Inflorescences were emasculated, bagged and labelled. The upper spathe was removed and the lower spathe was cut longitudinally so that pistillate flowers were exposed. The flowers were artificially pollinated with pollens from other plants by using a small camel hairbrush. This was bagged and labelled.

In *A. odora* pollination was done 126 hrs before opening of the inflorescence to 55 hrs after opening. In *X sagittifolium*, pollinations were done from 144 hr before the commencement of opening of the inflorescence to 24 hr after (i. e., up to the time of pollen shedding). The observations were recorded.

Seed germination

Seed viability and period of viability was tested in all genera. In each genus, except *Alocasia*, 2-3 accessions were used for the studies. Seeds were collected at maturity and dried under shade. Some seeds were also dried under the sunshine. Their viability was tested at regular intervals to study the longevity of seeds. This was done on moistened filter paper in Petri- plates.

Inter specific and Inter-generic hybridization

This was done in all the genera as listed in results. Inflorescences were emasculated, bagged and artificially cross pollinated with pollens taken from other species/accessions. The pollinated flowers were then bagged and labelled. Seeds were collected at maturity and seedlings were raised.

Thermogenesis

Thermogenesis in the inflorescence was studied in *Alocasia macrorrhizos*, *Amorphophallus paeoniifolius*, *Colocasia esculenta* and *Xanthosoma sagittifolium*. Temperature rise in different portions of the inflorescences was recorded. This was done by using an inserting type of digital thermometer. Temperature measurements were taken throughout the flowering period at regular intervals. Ambient temperature was also recorded each time.

Acridity and calcium oxalate estimation

Acridity

This is a major problem in the edibility of aroid tuber crops. This is the painful itching sensation caused when the plant parts or juice is contacted with the skin. The level of acridity was tested in all the cultivated and wild accessions. Small pieces of about 1 cm³ cooked and fresh tubers were rubbed on the forearm skin of three volunteers for checking the acridity. Based on the intensity of 'burning sensation,' the tubers were grouped into 4 categories; highly acrid, medium acrid, less acrid and non acrid. This was also repeated after cooking the corm pieces (2cm³) for about 30min.

Calcium oxalate estimation

Calcium oxalate crystals (raphides) present in corms and leaves are regarded as one of the factors, which causes acidity. This was estimated chemically by the methodology developed by AOAC (1975). Fresh leaves and tubers were cut into small pieces and dried in the oven. The dried tissues were then powdered in a mill. One gram of the powdered sample was then extracted in 20 ml of 0.25 N HCl in a water bath at 60°C for one hour. This was then centrifuged at 8000 rpm and the supernatant was collected in a conical flask. The debris is again extracted and centrifuged to ensure maximum extraction of calcium oxalate. The impurities in the supernatant were collected again and neutralized with dilute ammonia solution (1:1 solution). Calcium oxalate was then precipitated by adding 5 ml acetate buffer (pH 4.5). This was then centrifuged at 8000 rpm to collect the precipitated calcium oxalate. The precipitate was then washed with 6 ml washing liquid (5% acetic acid solution saturated with calcium oxalate). After this, centrifugation was done at 8000 rpm to remove the washing liquid. The washing process is again repeated. The precipitate was then dissolved in 15 ml of 2N H₂SO₄ and transferred into a conical flask. This was titrated with 0.01 N K₂Cr₂O₇ solution which was heated to 60°C. The percentage of calcium oxalate was calculated by using the following formula: % of calcium oxalate = $\frac{\text{Volume of } K_2Cr_2O_7 \text{ solution} \times \text{Normality of } K_2Cr_2O_7 \text{ solution} \times 100}{\text{Weight of sample}}$ The distribution of calcium oxalate crystals was also microscopically tested by taking thin sections of corm and examining in the low power of the microscope (100x).

Somatic chromosome count

Chromosome numbers were counted only in the parents and F₁ progenies inter specific and inter generic hybrids. Young and fresh root tips were treated with 0.002 M solution of 8-Hydroxyquinoline between 10.30 am and 12.00 a at 4°C for 3-4 hours and then fixed in 3:1 acetic acid ethyl alcohol mixture. Fixed root tips were squashed in 2% acetocarmine solution and observed under 1000x of a microscope for counting the metaphase chromosomes.

Isozyme analysis

Isoenzyme analysis was carried out in almost all the accessions by using a single enzyme, esterase. This was done to characterize the phylogenetic variations of individuals in populations and differentiate between accessions. Youngest fully opened leaves were collected in ice filled buckets for analysis. The leaves were washed in distilled water and cut into small pieces of 200 mg. This was then extracted in 2 ml of 0.1 M Tris extraction buffer of pH 7.4 using a mortar and pestle. The extraction was also carried out at 0°C. The extract

was centrifuged at 4°C at 10,000 rpm and the supernatant was collected. To this, 100 mg sucrose was added to increase the density of sample and a drop of bromophenol blue was added as an indicator. This was then electrophoresed in poly acrilamide gel having 7.5% resolving gel. The electrophoresis was also done in cold condition to preserve the stability of the enzyme. Slab gel was prepared between two perfectly clean glass plates held together with two spacers between them along the sides, placed vertically and held together by clamps. The running gel was poured rapidly after mixing, between the glass plates up to 2.5cm below the top. Gel was overlaid with butanol saturated with water to achieve a flat surface. After polymerisation of the gel, stacking gel solution was added on top of it and samples were applied on each slot using a micro syringe. The

wells were filled to the top using sample buffer. The top chamber was filled on the top of the glass plates and it was placed inside the apparatus chamber. The top and bottom chambers of the apparatus are then filled with electrode buffer and the electrodes were connected to power supply. The electrophoresis was done at low temperatures for keeping the stability of the enzyme. It was done at a current of 15mA till the sample enters the resolving gel and 30mA thereafter till the marker dye reached the bottom of the gel. At the end of the run, glass plates were removed from the chamber and after removing the spacers; the gel was taken out gently, washed and stained. The gel was then kept for 30-60 min until the blue-black bands appeared. The staining procedure was the same as used by Tanksley and Orten (1984). After incubation the gel was de-stained with 7% acetic acid solution. Storage was also done in 7% acetic acid solution. Electro morphs were used as isozymic descriptors. Each electro morph was considered as a character, with the presence and absence as the only two possible modalities. Samples were phylogenetically grouped by using UPGMA (Un-weighted Pair Group Analysis with Arithmetic Mean) of Bio-rad's gel documentation programme.

DISTRIBUTION AND ECOLOGY

This aspect was studied in four genera. They are *Alocasia*, *Amorphophallus*, *Colocasia* and *Xanthosoma*. The details collected during germplasm collection and other travels are included in this section.

Alocasia

Three species were collected during the study period. This includes *A. macrorrhizos* (14 accessions), *A. odora* (1 accession) and *A. sanderiana* (2 accessions). *A. sanderiana* is an ornamental plant. *A. odora* was collected from Jharnapani, Nagaland. This occurs naturally along the road side in clumps in partially open areas. *A. macrorrhizos*, giant taro is a cultivated minor tuber crop in Kerala, Karnataka and UP. In Dakshina Kannada and Udupi districts of Karnataka, the tubers are available in the vegetable market. This species occurs naturally in dry and partially shaded areas in clumps. This also occurs as 'escapes' from cultivation (ESS 48). Giant taro is widely distributed in Andamans, Karnataka, Kerala, and Meghalaya and UP.

Amorphophallus

The germplasm collection includes *A. bulbifer* (1 accession), *A. hohenackeri* (4 accessions), *A. paeoniifolius* (24 accessions) and *A. oncophyllus* (1 accession). *A. bulbifer* (ESS 42) is present in Wayanadu district of Kerala. This is also present in Dakshina Kannada and Uduppi districts of Karnataka. This species occurs in the forest under the shade, among bushes. *A. hohenackeri* was collected from Kottayam, Malappuram and Thrissur districts of Kerala. This species also occurs in dry and partially shaded areas among bushes. *A. paeoniifolius* var. *campanulatus* is a cultivated tuber crop, in many states in India. This is being cultivated in the open areas or in partially shaded areas among other crops in soil heaps or in raised beds in order to prevent water logging. *A. paeoniifolius* var. *paeoniifolius* is the wild type which is a widely occurring one. This variety is present under the shade of Coffeeplantation in Wayanadu district, Kerala (ESS 43). It is also usually present under the shade of big forest trees along road side (ESS 13 and ESS 20).

Colocasia

The collected species includes *C. affinis* (4 accessions) and *C. esculenta* (60 accessions). *C. affinis* was collected from Iddukki, Kottayam and Kozhikode districts in Kerala. This is also seen in Dakshina Kannada District of Karnataka. This occurs naturally in wet and partially shaded areas. *C. esculenta* (taro) occurs in both wild and cultivated condition and is widely seen in India. Wild taro occurs throughout the wet lands, in water channels, boundaries of lakes. The species is usually seen in the open areas or rarely in

partially shaded areas. Wild taro can be seen in dry lands in clumps only in the wet season. Taro can develop large populations in wet areas (ESS 1, about 3 Ha).

Wild taro can develop floating population in lakes. This type of floating population was noticed in Aakulam Lake, Thiruvananthapuram. When the first collection (ESS 1) was made there was no floating population of taro. At that time the upper end of the lake was completely filled with *Eichhornia crassipes*. When a revisit was made in the same area after four years, the major species which covered the water body was wild taro. Taro developed floating population over *Eichhornia crassipes* and covered about 90% of the water surface. The fresh and floating debris of *Eichhornia crassipes* gave support to the developing taro population. The average depth at this place of taro population was about 4 meter.

Xanthosoma

Two species were collected, *X nigrum* (1 accession) and *X. sagittifolium* (25 accessions). *Xanthosoma* is a cultivated tuber crop in many parts of India. This is an introduced plant from Central America. Tannia develops natural populations along road side, water channels, and coconut gardens in both wet and partially shaded areas. The centre of origin of this crop is Central America (Harris 1972). Hence the naturalized population observed during the study may be escapes from cultivation. The natural population of tannia can be seen through out Kerala. The present study on the above four genera indicate that wild taro is more ecologically adapted than the other three. Cultivated taro can be seen in a wide variety of soils from tropical laterite soils in coastal Kerala to Maharastra, sandy soils in West coast, alluvial soils along river banks, forest alluvium in the North East, and South India (Velayudhan 2008). The occurrence of wild and cultivated taro in a wide range of environmental conditions points out the adaptive significance of this plant.

GROWTH AND DEVELOPMENT

The growth and development aspects were studied in four genera: *Alocasia*, *Amorphophallus*, *Colocasia* and *Xanthosoma*.

Alocasia

One species, *A. macrorrhizos* (giant taro) with three accessions (ESS 47, 48 and 82) was used for the studies (Plate I a). Giant taro is the common cultivated form of *Alocasia* in India. It also occurs in the wild. Each accession was planted with three corm cuttings. Observations were recorded periodically. The data obtained are presented in (Tables 4.2.1, 4.2.2 and 4.2.3).

The cuttings began to genninate 2-5 weeks after planting (W AP). During sprouting, leaves and roots were initiated simultaneously. Successive leaves were found to emerge from the leaf sheaths of preceding leaves. New leaves emerged after 15-17 days. Lamina expansion was completed in 3-5 days after opening. Petiole growth was observed to be completed in 4-7 days after opening. Longevity of leaves was about 120 days. The plants produced about 85-95 leaves in a period of four years.

The plants produced larger leaves during rainy season. In rainy season, the plants retained 7-9 leaves. In summer, 5-7 leaves were retained. Vegetative growth was vigorous during the first three years (cf., Tables 4.2.1, 4.2.2 and 4.2.3) but was found to decrease gradually during the 4th year. The plants became stunted with 1-2 leaves (Plate Ie).

The growth of corm began with the emergence of the first leaf. Formation of successive leaves was found to add to the height of the corm by 1.0 cm each. The height of the conn was about 1.0 m after 4 years (Plate Id). About nine-tenth of the corm developed was above the soil. Conn is not branched above the soil. However, in few cases, up to eight side shoots are found to develop. These side shoots were developed from the underground portion of the corm. Cormels are stoloniferous. Vegetative reproduction occurs through the development of corm and cormels. However, in some accessions, cormels were not observed.

Period	Months after planting (MAP)	No. leaves	Length	Breadth	Petiole length
I Year	3 (June)	2.0± 0.0	26.3 ± 2.08	24.2 ±2.47	53.7 ± 2.51
	6 (September)	5.7 ± 0.57	44.0 ± 4.00	40.7 ±4.72	86.6 ± 13.57
	9 (December)	7.7 ± 0.57	57.7 ± 2.51	56.3 ±1.53	118.0 ± 2.51
	12 (March)	6.0 ± 0.0	48.7 ± 4.04	47.3 ±3.51	98.6 ± 10.5
II Year	15 (June)	7.0± 0.0	51.8 ± 2.78	51.4±2.78	106.8 ± 5.07
	18 (September)	8.3 ± 0.57	65.5 ± 2.34	64.2 ±1.94	130.8 ± 4.70
	21 (December)	8.3 ± 0.57	82.8 ± 5.45	81.6 ±6.10	168.0 ± 18.0
	24 (March)	6.7 ± 0.57	72.5 ± 3.39	71.3 ±3.44	144.0 ± 7.66
III Year	27 (June)	7.0± 0.0	80.6 ± 5.35	79.2 ±3.56	161.0± 11.86
	30 (September)	8.0± 0.0	94.2 ± 4.16	93.0 ±3.82	188.2 ± 7.46

IV Year	33 (December)	8.3 ± 0.57	87.9 ± 2.94	86.8 ± 2.79	175.7 ± 4.66
	36 (March)	5.7 ± 0.0	69.2 ± 3.86	68.0 ± 3.84	135.8 ± 6.68
	39 (June)	6.3 ± 0.57	56.0 ± 4.69	55.0 ± 4.84	112.7 ± 6.86
	42 (September)	4.0 ± 0.0	41.8 ± 3.91	40.1 ± 3.86	84.7 ± 4.22
	45 (December)	3.3 ± 0.57	33.2 ± 3.43	31.7 ± 3.14	64.5 ± 7.84
	48 (March)	1.3 ± 0.57	27.0 ± 2.82	25.7 ± 2.87	51.2 ± 4.21

Period	Months after planting (MAP)	No. leaves	Length	Breadth	Petiole length
I Year	3 (June)	2.0 ± 0.0	24.2 ± 2.64	22.0 ± 2.64	45.0 ± 5.62
	6 (September)	5.0 ± 0.00	34.5 ± 2.94	32.7 ± 2.49	62.2 ± 7.03
	9 (December)	7.0 ± 1.0	45.0 ± 2.19	43.2 ± 2.48	76.3 ± 2.16
	12 (March)	6.0 ± 0.0	39.5 ± 3.50	38.2 ± 3.81	75.2 ± 5.15
II Year	15 (June)	7.7 ± 0.58	50.7 ± 3.77	47.5 ± 2.79	94.5 ± 7.77
	18 (September)	8.0 ± 0.0	34.5 ± 2.59	52.8 ± 2.56	101.3 ± 2.50
	21 (December)	8.7 ± 0.58	63.5 ± 4.33	61.7 ± 2.56	119.7 ± 7.20
	24 (March)	6.3 ± 6.58	56.5 ± 2.25	55.2 ± 4.17	101.3 ± 6.86
III Year	27 (June)	7.0 ± 0.0	69.0 ± 2.34	67.3 ± 2.25	119.7 ± 2.16
	30 (September)	8.0 ± 0.0	78.7 ± 2.33	77.7 ± 2.06	140.0 ± 5.10
	33 (December)	8.0 ± 0.0	86.3 ± 3.72	85.1 ± 4.35	155.5 ± 3.62
	36 (March)	6.3 ± 0.57	70.7 ± 1.75	69.5 ± 1.76	128.2 ± 3.92
IV Year	39 (June)	6.0 ± 0.0	60.7 ± 2.73	57.3 ± 2.62	107.5 ± 5.24
	42 (September)	4.7 ± 0.58	49.7 ± 2.73	48.0 ± 3.03	85.7 ± 5.20
	45 (December)	3.0 ± 0.0	38.7 ± 2.65	37.5 ± 2.06	70.5 ± 5.08
	48 (March)	1.6 ± 0.57	31.7 ± 3.38	30.2 ± 2.78	57.8 ± 5.40

The corm began to produce stolons 5 months after planting (MAP). Production of stolons continued throughout the growth period. Stolon length reached up to 25 cm. Diameter of the stolon was about 3mm. The growth of the stolons was completed in about 30-45 days of its emergence. Stolons remain below the soil. Their tips become bulged to form a connel. Connels were observed to develop in 12 out of the total 13 accessions. In one accession (ESS 47, Kozhikode, Kerala) up to 602-856 cormels were found to occur. Longevity of stolons was about four months. Cormels get separated from the mother conn by the death and decay of stolons. They remain dormant in the soil. When the mother conn was removed several of the cormels were observed to sprout.

Amorphophallus

The species studied in this genus includes *A. bulbifer*, *A. hohenackeri*, *A. paeoniifolius* and *A. oncophyllus*.

The details of the growth and development in leaves are given in table 4.2.4.

Identity of	Leaf emergence from cataphyll	Opening	No. days required for	No. leaves produced	Leaf longevity in
<i>A. paeoniifolius</i>					
ESS 13	18	26	43	1	198
ESS 17	20	25	40	1	198
ESS 20	18	28	46	1	196
ESS 28	15	25	45	2	193
ESS 29	25	34	52	1	208
ESS 30	15	23	40	1	195
ESS 40	17	24	43	4	201
ESS 41	17	26	46	1	223
ESS 43	16	21	41	1	212
ESS 54	18	27	44	1	201
ESS 56	15	23	41	1	193
ESS 57	22	28	47	2	197
ESS 62	23	30	49	1	194
ESS 66	18	25	40	1	205
ESS 67	17	22	38	2	208
ESS 68	17	23	40	1	217
ESS 69	15	23	41	3	211
ESS 80	20	30	46	1	208
ESS 114	20	29	49	1	201
ESS 115	21	27	45	4	214
<i>A. bulbifer</i>					
ESS 42	21	26	46	1	162
<i>A. hohenackeri</i>					
ESS 44	17	24	39	1	177
ESS 59	20	29	46	1	172
ESS 60	18	28	44	1	170
ESS 65	21	31	47	1	172
<i>A. oncophyllus</i>					

ESS 81	16	21	40	1	132
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The corm began to germinate in 20 days after planting (DAP). The leaf and the roots developed simultaneously during sprouting. Several roots were emerged from the base of the growing bud situated in the center of the corm. Leaves are not observed to be produced on other parts of the mother conn. When the central bud was removed, several lateral buds developed into leaves. Cataphylls protects the growing leaf. The leaf was found emerging from the cataphylls after 15 to 25 days of growth. It opened in 5-10 days after emergence. Expansion of lamina and petiole growth is completed in 15-20 days after opening (Palte 5h-j). Usually, only one leaf was produced in a season. Rarely 2-4 leaves were seen (Table 4.2.4). One accession, where the conn was highly branched, (ESS 40; originated from Tamil Nadu) 33 leaves were found to be produced.

The leaf development or flowering occurs in the months of April - June only after the summer rains. If the growing point is damaged, the corm can produce up to 7 vegetative sprouts. In all the cases, during vegetative growth the mother corm was completely used. As a result, it became shriveled and decayed in the first 3 months after planting.

Longevity of leaves varies with accession. In *A. bulbifer*, *A. hehenackeri* and *A. oncophyllus* longevity varied from 132 - 177 days. In *A. paeoniifolius*, this was 193-223 days (cf., Table 4.2.4).

Accessions of *A. bonaccordensis* and *A. hohenackeri* produced 2-4 stolons. The stolons are unbranched, starchy and remained below the soil. Their tips are bulged. Stolons are 30 to 50 cm long and 3-5 mm in diameter. New plants develop from the tip of the stolons. When the tip is damaged, new plants emerge from the nodal bud.

The vegetative and reproductive phases are separate in *Amorphophallus*. During the vegetative phase, the plant stores food in the corm. This phase lasts for 6-10 months. In the following year, in most cases, if the growing point is not damaged, the corm will produce an inflorescence, irrespective of whether it has remained in the soil, or harvested and stored. If it is grown from seeds, the plant takes about 3-4 years to flower. The reproductive phase lasts for about 8 months. The destruction of the central bud induces the formation of many lateral buds. Adequate number of lateral buds is needed when the

corm is cut into pieces for planting. Hence the destruction of the main bud induces the formation of adequate lateral buds and prevents flowering. Usually the corm produces a single inflorescence or single leaf if the main bud is not damaged.

Colocasia

Only one species, *C. esculenta* (taro) with 3 accessions (ESS 1, ESS 24 and ESS 124) was used for the studies. Three plants from each accession were selected for the studies. Each accession was planted with corms of about 30g. Stolon development was studied in 31 accessions. Wild taro plants occurring in wet stabilized habitats were also observed.

The conns began to germinate 2-4 W AP. Leaves and roots emerged simultaneously during sprouting. The first leaf was irregular in shape. However, from second leaf onwards, nonnal size was regained. Successive leaves emerge from the leaf sheaths of the preceding leaves. New leaves emerged in a span of 8-10 days. They opened 4-5 days after emergence. Lamina expansion is completed in 4-6 days after its opening. Longevity of leaves varies with accessions (ESS 1- 38 days; ESS 124-35 days and ESS 23-30 days). The accessions produced about 20-25 leaves in a season. Flowering took place after the formation of 8-In the cultivated accession (ESS 23), cormel production started after the emergence of 6-10 leaves. All cotmels were produced simultaneously. Connel growth continued till the senescence of the leaves.

Vigorous vegetative growth was observed during 13-14 W AP (Table 4.2.5, 4.2.6 and 4.2. 7). Thereafter the rate of leaf production was found to decrease (about 1-2 leaves in a month) and gradually stop. This happened in 38-47 WAP. Wild taro plants, which occur in wet stabilized habitats, produced leaves continuously. Their growth was indefinite.

Wild accessions in *C. affinis* (Plate 81) and *C. esculenta* produce stolons. Stolons are branches of conn. They are elongated, branched or unbranched structures. They begin to appear after the development of 4-6 leaves. Their production continued throughout the growing season. They develop from the axillary buds present in the leaf scars of the conn. They may also appear from the axils of leaves. The stolon number, their length and growth varied with different accessions (Table 4.2.8a,b). One accession (ESS 1;

Aakulam, Thiruvananthapuram) produced 27 stolons in a season. The length of the stolon in one accession (ESS 9, Pangode, Thiruvananthapuram) was 2.27m

. Leaf growth in taro ESS 1 wild taro, Thiruvananthapuram

Days after planting (DAP)	No. leaves	Length	Breadth	Petiole length
30	1.0 ± 0.0	10.8 ± 1.58	9.2 ± 0.28	24.8 ± 3.30
60	3.0 ± 0.0	17.41 ± 1.50	14.5 ± 1.3	55.0 ± 6.24
90	4.7 ± 0.58	25.7 ± 3.47	20.1 ± 4.61	89.9 ± 14.87
120	5.0 ± 0.0	33.9 ± 4.71	27.6 ± 4.80	138.0 ± 11.26
150	5.0 ± 0.58	37.6 ± 2.68	33.9 ± 2.51	145.8 ± 5.03
180	4.7 ± 0.58	34.8 ± 4.25	29.5 ± 3.04	118.7 ± 6.69
210	4.0 ± 0.0	28.3 ± 3.15	25.3 ± 2.09	107.5 ± 8.04
240	3.0 ± 0.0	22.2 ± 4.78	17.4 ± 2.22	76.8 ± 7.62
270	2.3 ± 0.58	18.6 ± 3.41	14.5 ± 1.32	61.4 ± 3.14
300	1.7 ± 0.58	13.7 ± 1.05	11.2 ± 1.85	41.3 ± 5.08
330	nil	Nil	Nil	nil

SUMMARY

The goal of the current study is to characterize genetic variations and breeding aspects like selfing, inter-varietal, interspecific, and intergeneric hybridizations in some genera of the family Araceae. It also aims to understand growth and development, reproductive biology, ecology, thermogenesis, and the role of calcium oxalate in acidity. The outcomes are listed below in brief. Giant taro, wild taro, and wild type tannia all continue to grow. The growth cycle of cultivated taro and tannia takes a year to complete. If produced from seeds, the elephant yam requires 3–4 years of vegetative growth before it can flower. In accessions of giant taro, flowering is relatively uncommon. Only one of the 14 accessions naturally bloomed. Giant taro does not have natural seed germination. Application of GA3 in an effort to artificially stimulate blooming was unsuccessful. In big taro, taro, tannia, and elephant yam, opening of the spadix (female blooming) and male flowering (pollen shedding) are separated by roughly 24 hours. Except for *Amorphophallus odora*, whose pollen shedding persisted for almost 36 hours, pollen shedding is always finished within 1-2 hours. All of the investigated species don't appear to be subject to wind pollination. Elephant yam was found to be cross-pollinated by taro-pollinating Dipteran flies and dung beetles

(*Onthophagus* sp.). In elephant yam, natural pollination, fertilization, and seed germination were accomplished by keeping pollinators away from the flowers. Elephant yam doesn't appear to self-pollinate. In *C. affinis*, there is no natural seed germination. There were no visible pollinators on this plant. The blooms couldn't be artificially pollinated, unfortunately. In contrast to cultivated taro, wild taro frequently flowers.

fertilization, and ultimately, seed germination. It appears that taro does not naturally self-pollinate. In this plant, the flowers can artificially self-pollinate. Additionally feasible is artificial cross-pollination. Self-pollination is not possible in taro due to protogyny and the unique characteristics of the spathe (constriction and pollen capturing mechanism). When the sterile appendage in a taro inflorescence was removed, the flowering process, including the creation of fragrance, was stopped. The flowers were kept in this state for twenty to twenty-five days. In tannia, stigma receptivity varies from 126 hours prior to spadix opening to 17 hours following spadix opening (stigma is receptive for 5–6 days).

In Tannia, artificial selfing and crossover are both conceivable. There is no seed dormancy in tannia. The seeds can be sown after 75 to 90 days.

All of the genera investigated in this inquiry produce odors. Different species exhibit different smells and produce them at different locations.

64.2% of large taro, 33.5 % of *A. odora*, 39.3 - 91.7% of elephant yam, 89.6 - 94.1% of taro, and 95.6 - 98.6% of tannia produced pollen. Giant taro's pollen viability persisted for 5–6 hours, elephant yam's for 22 hours, taro's for 8 hours, and tannia's for 20 hours. Environmental factors may have an impact on taro tuber formation. When the wild stoloniferous taro plants were planted in garden pots and in the uplands, tuber development was observed. This outcome presents issues with the taxonomy of wild taro as *C. esculenta* var. *aquaticus* and points to the origin of cultivated taro. The results of the current investigation demonstrated that all four genera experience thermogenesis. All thermogenesis was shown to follow a common pattern. Two cycles of temperature rise take place, the first during the female phase and the second during the male phase. The opening process and the volatilization of the odoriferous substances to draw pollinators may be influenced by the increased temperature. This could create an unpleasant environment inside the spadix. The dehiscence of the anther may be aided

by thennogenesis during the second cycle. According to the current study, all four genera experience thermogenesis. It was discovered that all thennogenesis followed a same general pattern. In two cycles, one during the female phase and the other during the male phase, temperatures rise. The process of the flower opening and the volatilization of the odoriferous substances that draw pollinators are both influenced by the increased temperature. This could also create an uncomfortable environment inside the spadix. The second cycle's thennogenesis may be aiding the anther's dehiscence. Studies on acidity and estimates of calcium oxalate show that acidity and calcium oxalate levels alone cannot be connected. Concentrated calcium oxalate crystals of raphides and druses were discovered in the corm's outer layer. It was discovered that leaves have higher calcium oxalate content than corms. Long-term cooking in extremely caustic additions did not get rid of the acidity. The population of wild taro varies in acidity and calcium oxalate. This suggests the potential for choosing this crop and the historical context of taro production. The outcome of the morphological characterization shows that the taro germplasm exhibits the greatest degree of variety. Despite its origin being attributed to either India or Sri Lanka, huge taro is the crop with the least variety. In terms of diversity in India, tannia comes in third and elephant yam comes in second. The analysis of the germplasm using a single isozyme esterase also revealed significant differences between farmed and wild taro. Taro has 33 bands in total that were visible. The cladogram revealed certain wild and domesticated accessions shared genetic traits. The outcomes also showed that there is no relationship between isozyme and morphological variation. Results of isozyme variations in arbitrary samples of *C. elegans* populations in the wild. flora here. Taro populations do, however, include genetic differences. The natural seed-setting process may be to blame. The current study shown that taro, tannia, and *A. odora* are capable of artificial inbreeding. In the species *A. odora* and *A. macrorrhizos*, interspecific hybrids were produced. *Amorphophallus* (*A. paeoniifolius* var. *campanulatus* x *A. paeoniifolius* var. *paeoniifolius*) has also produced intervarietal hybrids. This study's intergeneric hybrids between *Xanthosoma sagittifolium* and *Caladium bicolor* are the first of their kind in the Araceae family.

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Publications

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Review on Role of Mathematics in Coding Theory

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Abstract

Mathematics plays a crucial role in coding theory, a field that deals with error detection and correction in digital communication. Mathematical concepts like linear algebra, finite fields, and combinatorics support the design and analysis of error-correcting codes. These codes are used to transmit data accurately over noisy channels. By employing mathematical principles, coding theory enables the creation of efficient and reliable codes, ensuring data integrity in various applications such as telecommunications, data storage, and cryptography.

Introduction

Coding theory is concerned with transmitting data successfully over a noisy channel and correcting errors in corrupted data. It is critical for a wide range of applications in computer science and engineering. The significance of algebra, combinatorics and geometry in coding theory is a well-known fact, with many deep mathematical results used in the creation of coding theorems. Error-correcting codes are concerned with increasing the reliability of communication over noisy channels and have a wide range of applications, such as the transmission of images from distant locations, space, CD sound quality, phone lines, computer networks, wireless communication, ISBN numbers. In 1948, Claude Shannon published a landmark paper 'A Mathematical Theory of Communication' [94] that started discipline in the concepts of both information theory and coding theory. This paper marked the birth of coding theory, a field of study concerned with the transmission of data across noisy channels and the recovery of corrupted messages. Richard Hamming was among the first to actually build and put error-correction codes in place. The development of the compact disc in 1970 resulted in another application of error correcting codes. The signal on CDs is digitally encoded. Reed Solomon Codes are used to protect against scratches, cracks, and other similar damages.

Importance of Coding Theory

Coding theory is of paramount importance in modern digital communication and technology. Coding theory provides methods for Error Detection and Correction that occur during data transmission. This is crucial for maintaining data integrity and ensuring accurate communication in the presence of noise or interference. By using error-correcting codes, coding theory enables reliable communication over imperfect channels. This is essential for applications like wireless communication, satellite communication, and internet data transfer, where data can get corrupted during transmission. Coding theory helps in optimizing the use of resources such as bandwidth and storage space. Efficient codes can transmit information using minimal resources, leading to improved overall system performance. Coding theory contributes to cryptographic systems, enhancing the security of data transmission and storage. Techniques like error-detection-based authentication and secure key exchange rely on coding theory concepts. Medical imaging technologies such as MRI and CT scans generate vast amounts of data that need to be transmitted accurately for diagnosis. Coding theory ensures the integrity of this crucial medical information. Coding theory has given rise to the concept of network coding, where data packets are mixed at intermediate nodes in a network to optimize transmission efficiency and reliability. In essence, coding theory serves as the foundation for robust and dependable communication in the digital age, impacting a wide range of sectors and technologies that rely on accurate data transmission and storage.

Mathematical scope in Coding Theory

The mathematical scope in coding theory is extensive, as the field heavily relies on various mathematical concepts and structures to design, analyze, and implement error-correcting codes. Some key mathematical areas within coding theory include:

Algebraic Structures: Coding theory draws heavily from algebraic structures such as finite fields (also known as Galois fields), which provide the foundation for constructing linear codes and designing error-correction mechanisms.

Linear Algebra: Linear codes, which form a significant subset of error-correcting codes, are constructed using linear algebra techniques. Concepts like vector spaces, matrices, and linear transformations are essential for understanding and creating these codes.

Polynomial Algebra: Polynomial rings and polynomial arithmetic play a crucial role in defining and analyzing cyclic codes and Reed-Solomon codes, widely used in modern communication systems.

Combinatorial Mathematics: Combinatorial concepts are employed to design codes with specific properties. This includes topics like combinatorial designs, which can be used to create codes with optimal error-correction capabilities.

Graph Theory: Graphs are used to represent and analyze codes, particularly in non-linear codes and iterative decoding algorithms. Tanner graphs and factor graphs are examples of graphical representations used in coding theory.

Number Theory: Certain error-correcting codes, such as BCH (Bose-Chaudhuri-Hocquenghem) codes, are built on number-theoretic principles, making number theory a relevant mathematical area.

Abstract Algebra: Advanced algebraic concepts, such as group theory and ring theory, are used to understand the properties and structure of various types of codes.

Probability and Statistics: Probability theory is utilized to model noise and errors in communication channels. Statistical techniques are applied to analyze the performance of codes under different conditions.

Coding Bounds and Parameters: Mathematical bounds, such as the Singleton bound, Gilbert-Varshamov bound, and Hamming bound, provide theoretical limits and benchmarks for the performance of codes.

Coding Algorithms: The design of efficient encoding and decoding algorithms involves algorithmic techniques, complexity analysis, and optimization, all of which have strong mathematical foundations.

Finite Automata and Formal Languages: These concepts are used in designing and analyzing sequential decoding algorithms for codes.

Information Theory: While not strictly within coding theory, information theory, particularly Shannon's theory of channel capacity, provides a theoretical basis for

understanding the limits of data transmission and the trade-off between data rate and reliability.

Limitations of Coding Theory

While coding theory has many advantages, it's important to acknowledge some of its drawbacks and limitations. Implementing and decoding sophisticated error-correcting codes can be computationally intensive, requiring significant processing power and memory, especially in real-time applications. Many error-correcting codes introduce redundancy to detect and correct errors. This overhead reduces the effective data rate, which can be a concern in applications with limited bandwidth or storage. Also coding schemes are designed for specific block lengths, which might not be suitable for all scenarios and finite fields which requires arithmetic operations. The performance of these codes can degrade when used with data that doesn't match their block length. The decoding process can be complex, especially for advanced codes. Iterative decoding algorithms used for more complex codes might require multiple iterations to achieve desired error correction, leading to increased latency. Static error-correcting codes might not perform optimally in dynamic or rapidly changing channel conditions, requiring the use of adaptive coding schemes. While error-correcting codes can mitigate errors up to a certain extent, there's a limit to the number of errors they can effectively correct. Beyond this limit, the codes might fail to provide accurate correction. More powerful error-correcting codes often come with increased complexity. Balancing complexity against performance requirements can be challenging. Many error-correcting codes assume specific noise models. If the actual noise doesn't match the assumed model, the performance of the code might be suboptimal. Due to factors like hardware limitations, channel variations, and other practical considerations the theoretical performance of error-correcting codes might not always translate exactly to real-world scenarios.

Despite these drawbacks, coding theory remains a critical field that continues to evolve and adapt to new challenges. Researchers and engineers work to address these limitations and develop innovative solutions that strike a balance between performance, complexity, and practicality in various applications.

Conclusion

As technology evolves, coding theory continues to play a role in emerging fields like quantum communication and DNA data storage, ensuring that information can be transmitted accurately and securely in new and challenging contexts. The mathematical scope in coding theory continues to evolve as researchers explore new ways to optimize codes for various applications and challenges. Understanding these mathematical concepts is essential for both theoretical advancements and practical implementations in fields such as communication, data storage, and beyond.

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Publications

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Ethics in Teaching Pedagogy

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Abstract

This paper proposes the idea of pedagogy in the education sector. It is an exercise to reflect the ethics of teaching. It is questioned by scholars why ethics are important for teaching. A light is thrown by several authors from the perspective of educational philosophy so that the teachers highlights the challenges and explore teaching practices. Teachers can be capable to built pedagogical relationship that will be helpful in the learning processes of students. Teacher works as a friend, philosopher and guide. Teacher is capable to establish a sensitive bond with students and enhance the ethics of teaching.

Content

A teacher plays a key role in education and also in student's life. Teaching job is a great responsibility than a mere job. It has an impact on the growth and well-being of the nation. A teacher has far-reaching influence on the society; no other personality can have an influence more profound than that of teacher. Because they have a daily influence on the lives of children, teachers are often held to high standards. Ethic is a fundamental requirement of any profession. As because Teaching creates all other professions, therefore, it must consist a high code of professional ethics so that the same values can be developed among students. Therefore a teacher must serve himself as strong role models and demonstrate ethical behaviors as the interaction with students , colleagues , parents and others. Professional code of ethics help to make sure teachers act in a professional and ethical manner at all times. They are role models within themselves for their students who, consciously or subconsciously, emulate their behavior. They instill values and morals that will influence the lives of their students. Teachers have the responsibility to understand the needs of their student's even outside the school walls and they have the responsibility to nurture and to cherish their students. Rabindra Nath Tagore Says, "Teacher can never truly teach unless he is still learning himself. A lamp never lights another lamp unless it continues

to burn it own flame. The teacher who has come to an end of his subject, who has no living traffic with his knowledge merely repeats his lesson to his students can only load to their mind."¹

All of us are aware that a teacher's main responsibility is to the student and to the development of his/her full potential. In this respect, we need to pursue the truth, devote all our efforts towards excellence, the acquisition of knowledge and the observance of democratic principles. Mahatma Gandhi, (Young India, 24 January 1925) emphasized that " The teacher himself must possess the virtues that he wants to inculcate in the students. This means that the teacher must practice these virtues himself, otherwise his words will have no effect."² He further highlighted (Young India, April 1929)that " the teacher should be able to establish a heart to heart contact with the students....."³ The teacher and the students should be in constant rather than their brains. About the ethical duties of the teacher, Sri Aurobindo says, " The teacher is not an instructor or task-master ; he is a helper and guide. His business is to suggest and not to impose..... he does not impart knowledge to him ; he shows him how to acquire knowledge for himself. he does not call forth the knowledge that is within, he only shows him where it lies and how it can be habituated to raise to the surface."⁴

The relationship between teachers and students is a very important and sensitive one. It must be built on strong foundations, stemming from mutual respect and trust as well as highest of ethical standards. But now the scenario is large no. of teachers are fails to establish this sensitive bond because of lack of ethical values in them. Felicity Haynes (1998) draws from a broad range of ethical theory to argue for an ethics of consequences, consistency, and care- "The greater part of our learning in the school has been a waste because most of our teachers think their subjects are like dead specimens of once living things with which they have a learned acquaintance but no communication of life and love."⁵ The same may be said of the relationship between teachers and other members of the school community that includes, among others, the administration , other professionals and parents. It's teacher's duty to maintain a high level of competence and in order to do with we need to engage in continuous professional development throughout our career in the profession.

But now a day's only few no. of teachers got success in developing proper understanding, mutual, respect and trust with every member of teaching community. This ratio is very low due to the lack of professional ethics in the teaching community and Educational Institutions. first, let us understand what professional ethics are and how we nurture it.

Meaning of Professional Ethics

Every profession, in order to regulate its terms, conditions, norms and quality of service rendered, has its own professional ethics, which is different from general ethics. In the term professional ethics, the word "ethics" adds to the professional obligation that a profession abides by. Professional ethics is a combination of two words, Professional and Ethics. Here, Professional means an expert, specialized, qualified, proficient, skilled, trained, practiced, certified, proficient, licensed, mature etc. So, Professional is a term denoting a level of knowledge and skills possessed by an individual and required of an individual to perform an assignment, that is attained through extensive education and training.

Secondly, Ethics means principles, morals, beliefs, moral principles, moral values, moral code etc. Ethics is a system of moral principles governing the appropriate conduct of a person or a group. Ethics refers to human conduct as to make judgments between what is right and what is wrong. It could be that there are several factors that may encourage one to adopt unethical behavior, but the right person is he who, despite facing ethical dilemmas, assesses the situations and makes differentiation between what is morally good and bad in order to follow the rules and code of professional conduct. Good ethics causes to gain confidence of superiors while promoting integrity, which means doing right things even when we are not watched. Albert Schweitzer says, "Ethics is the activity of man directed to secure the inner perfection of his own personality."⁶

Why ethics for a noble profession-teaching:

- Teachers themselves represent as a role model of students.
- Promote a positive image of the teaching profession.
- Teaching is now a profession rather than a passion.
- Teacher work as a Friend, Philosopher and Guide.
- Establish a sensitive bond with students.
- Great impact in the molding of the next generation.
- To maintain healthy and supportive ambience at workplace.
- Paradigm shift in the perception of teachers.
- It's no longer a service but an occupation with unclear roles, vision and mission.
- Improve Regulation and control of teacher misconduct.
- Erosion in the values, responsibilities, commitment in the profession.

- A Teacher no longer enjoys the same respect and status in the society.
- Perplexed with new development and cultural heritage.
- Commercialization and profit making is the buzz word.
- Availability of varied and financially rewarding opportunities in others sectors has affected teaching profoundly.

Professional Ethics

A member of the teaching profession shall maintain mutual trust and respect with students, act with honesty, integrity and fairness; take responsibility for maintaining the quality of their professional practice; create learning experiences which engage, motivate and challenge students in an inclusive setting with a lifelong learning perspective. A teacher should respect the uniqueness and diversity of students; demonstrate respect for diversity and promote equality irrespective of gender, race, religion, age, language or different abilities; maintain an up to date knowledge and understanding; maintain professional relationship with students and duty bound.

Teaching Professionals shall dwell with other members of the profession in the same manner as the teacher wishes to be treated; respect, support and collaborate with colleagues both in matters concerning the education of students as well as in maintaining relations with colleagues in the highest standard of professional courtesy; be prepared to help junior colleagues; not censure or criticize any fellow in the presence of students or in public. A teacher establishes trust with parents or guardians in the interest of all round development of students; develop good relationships between home and school, respecting the role that parents have in students' education. A teacher can develop respect for the composite culture of country among student and avoid taking part in such activities that spread feelings of hatred among different communities, religions or linguistic groups.

To conclude this paper we can say that teaching is the noblest of all professions. In order to maintain their professional status, teachers have to own certain responsibilities both as individual and as members of a respectable profession. If the teachers bears good professional ethics in relation to his profession, the ethics are automatically transformed to the coming generations. The deteriorating status of the profession will gain back its potential status. Professional ethics will help in the spread of peace and international understanding across the global platform.

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Internet of Things (IoT) and Smart Cities: Innovations for Urban Sustainability

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Abstract

This research paper investigates the role of the Internet of Things (IoT) in shaping smart cities and promoting urban sustainability. As urban areas grapple with challenges such as overpopulation, resource depletion, and environmental degradation, IoT technologies offer innovative solutions that enhance urban living and operational efficiency. Through a comprehensive literature review and analysis of case studies from cities like Barcelona, Singapore, and Amsterdam, the study demonstrates how IoT applications in areas such as energy management, traffic optimization, and waste reduction contribute to sustainable urban development. Findings reveal that IoT facilitates real-time data collection and analysis, enabling informed decision-making and efficient resource allocation. However, challenges such as data privacy concerns, cybersecurity risks, and the digital divide must be addressed to fully realize the potential of IoT in smart cities. This paper emphasizes the importance of collaborative efforts among stakeholders to foster the integration of IoT technologies, ultimately advocating for their critical role in achieving urban sustainability and resilience in the face of rapid urbanization.

1. Introduction

The rapid growth of urban populations has led to an unprecedented demand for innovative solutions to address the challenges faced by cities worldwide. Urban areas are increasingly becoming hotspots for economic activity, but this growth also brings with it a plethora of issues, including traffic congestion, inefficient resource management, environmental degradation, and inadequate infrastructure. According to the United Nations, over 55% of the global population currently resides in urban areas, a figure projected to rise to 68% by 2050. This urbanization trend necessitates a paradigm shift in how cities operate and manage their resources.

The concept of smart cities has emerged as a potential solution to these urban challenges, leveraging advanced technologies to improve the quality of life for residents and promote sustainable practices. Central to this vision is the Internet of Things (IoT), which encompasses a network of interconnected devices that collect, exchange, and analyze data to facilitate intelligent decision-making. By integrating IoT technologies into urban infrastructure, cities can optimize their operations, enhance service delivery, and reduce environmental impacts.

IoT applications in smart cities are diverse, spanning various domains such as transportation, energy, water management, and waste disposal. For instance, smart traffic management systems can reduce congestion by optimizing traffic flow based on real-time data, while smart energy grids enable efficient energy consumption through monitoring and automation. These innovations not only improve urban efficiency but also contribute to sustainability goals by minimizing resource waste and reducing carbon footprints.

Despite the promising potential of IoT in transforming urban environments, the integration of these technologies poses significant challenges. Concerns about data privacy and security, the digital divide that may exclude underserved populations from technological benefits, and the need for substantial investments in infrastructure are critical issues that must be addressed. Moreover, effective governance and collaboration among various stakeholders—government agencies, private sector entities, and citizens—are essential for the successful implementation of smart city initiatives.

This paper aims to explore the transformative impact of IoT on the development of smart cities and their role in fostering urban sustainability. It will review existing literature on IoT applications in urban contexts, analyze case studies of cities that have successfully implemented IoT solutions, and discuss the challenges and barriers to integration. Ultimately, the study seeks to highlight the importance of adopting IoT technologies as part of a comprehensive strategy to create resilient, sustainable urban environments that can meet the needs of present and future generations.

2. Literature Review

2.1 Defining Smart Cities and Urban Sustainability

The term "smart city" refers to urban areas that leverage digital technologies and data-driven solutions to enhance the quality of life for residents, improve city operations, and promote sustainability (Angelidou, 2015). Smart cities integrate information and communication technologies (ICT) with IoT devices to create an interconnected ecosystem that enables real-

time monitoring and management of urban resources. Urban sustainability, on the other hand, is defined as the ability of cities to meet the needs of the present without compromising the ability of future generations to meet their own needs (Wang et al., 2017). This dual focus on technological advancement and sustainable practices positions smart cities as a potential solution to the pressing challenges faced by rapidly urbanizing populations.

2.2 IoT Technologies and Applications in Smart Cities

IoT technologies play a pivotal role in enabling smart city initiatives by providing the infrastructure necessary for data collection, analysis, and action. IoT devices, such as sensors, cameras, and smart meters, gather data on various urban parameters, including air quality, traffic patterns, energy consumption, and waste levels (Khan et al., 2019). This data is then analyzed to inform decision-making and optimize city operations.

One significant application of IoT in smart cities is in transportation management. Intelligent transportation systems utilize IoT technologies to monitor traffic flow, manage public transportation, and reduce congestion (Xie et al., 2020). For instance, smart traffic lights can adjust their timings based on real-time traffic conditions, thereby improving traffic flow and reducing wait times.

In the realm of energy management, IoT technologies facilitate the development of smart grids that optimize energy distribution and consumption (Guan et al., 2020). Smart meters enable consumers to monitor their energy usage in real time, encouraging more efficient consumption patterns and supporting demand response initiatives. Additionally, IoT applications in waste management, such as smart bins that monitor waste levels and optimize collection routes, contribute to more efficient urban services (Vázquez et al., 2021).

2.3 Challenges and Barriers to IoT Implementation

Despite the myriad benefits of IoT technologies in smart cities, their implementation is not without challenges. Data privacy and security concerns are paramount, as the extensive collection and analysis of personal and environmental data raise issues related to surveillance and data misuse (Zhang et al., 2018). Ensuring robust cybersecurity measures is crucial to protect against potential threats that could compromise sensitive urban data and disrupt city operations.

Furthermore, the digital divide presents a significant barrier to the equitable implementation of IoT solutions. Access to technology is often uneven, with underserved populations facing limitations in connectivity and digital literacy (Hollands, 2008). As cities adopt smart

technologies, it is essential to consider inclusivity and ensure that all residents benefit from the advancements.

Investment in infrastructure is another critical challenge. The integration of IoT into existing urban systems often requires substantial financial resources and technical expertise, which may not be readily available in all municipalities (Nam & Pardo, 2011). Therefore, strategic planning and collaboration among government agencies, private sector partners, and communities are necessary to overcome these barriers and facilitate the successful adoption of IoT in smart cities.

2.4 Theoretical Frameworks Supporting IoT Integration

Several theoretical frameworks can guide the integration of IoT technologies into smart city initiatives. The Smart City Framework by Deakin and Al Waer (2011) emphasizes the importance of governance, technology, and social engagement in creating sustainable urban environments. Similarly, the Triple Helix Model of innovation highlights the collaboration between academia, industry, and government as essential for fostering innovation and addressing urban challenges (Etzkowitz & Leydesdorff, 2000).

These frameworks underscore the need for a holistic approach to IoT integration, where technological advancements are complemented by effective governance and community involvement. By leveraging these theoretical foundations, cities can better navigate the complexities of implementing IoT solutions and achieve their sustainability goals.

3. Methodology

This research employs a mixed-methods approach, combining qualitative and quantitative analyses to comprehensively investigate the role of the Internet of Things (IoT) in enhancing urban sustainability within smart cities. The methodology consists of a systematic literature review, case study analysis, and thematic analysis to identify key trends, challenges, and best practices related to IoT applications in urban contexts.

3.1 Systematic Literature Review

The systematic literature review aims to gather existing research on the integration of IoT technologies in smart cities and their implications for urban sustainability. The review process includes the following steps:

1. **Research Question Formulation:** The primary research question guiding this study is: "How do IoT technologies contribute to urban sustainability in smart cities?" This question is further broken down into sub-questions addressing specific applications of IoT in areas such as energy management, transportation, and waste management.

2. **Database Selection:** Academic databases, including Scopus, IEEE Xplore, and Google Scholar, were chosen for the literature search. These databases provide access to peer-reviewed articles, conference papers, and industry reports relevant to the topic.
3. **Inclusion and Exclusion Criteria:** The review focuses on publications from the last ten years (2013-2023) to ensure relevance and currency. Studies included in the review must specifically address IoT applications in smart cities and their impact on sustainability. Non-English publications, opinion pieces, and articles not centered on urban sustainability were excluded.
4. **Data Extraction and Analysis:** Key information from selected studies, including the main findings, methodologies, and identified challenges, was extracted and organized. This data was then analyzed to identify patterns, trends, and gaps in the existing literature.

3.2 Case Study Analysis

To complement the literature review, case studies of selected cities implementing IoT solutions were analyzed. The case studies chosen for this research are Barcelona, Singapore, and Amsterdam, representing diverse geographical and cultural contexts. The analysis of these cities includes:

1. **Selection Criteria:** The cities were selected based on their recognized initiatives in implementing IoT technologies and their commitment to urban sustainability. Each city has demonstrated innovative approaches to leveraging IoT for enhancing urban services.
2. **Data Collection:** Data for the case studies was collected from a combination of primary and secondary sources. Primary data includes interviews with city officials and stakeholders involved in smart city projects, while secondary data encompasses city reports, project documentation, and academic publications.
3. **Data Analysis:** The collected data was analyzed through thematic analysis, focusing on how each city's IoT initiatives contribute to sustainability goals. Key themes explored include the impact of specific IoT applications on resource management, citizen engagement, and environmental outcomes.

3.3 Thematic Analysis

Thematic analysis was utilized to synthesize findings from both the literature review and case study analyses. This approach involves:

1. **Identifying Themes:** Key themes emerging from the data include the benefits of IoT for urban sustainability, challenges related to data privacy and security, and the importance of stakeholder collaboration.
2. **Interpreting Findings:** Each theme was examined in relation to the research questions, highlighting how IoT technologies can address urban sustainability challenges and the barriers that cities may face in their implementation.

3.4 Limitations of the Study

While this methodology provides a comprehensive understanding of the role of IoT in smart cities, several limitations should be acknowledged:

1. **Scope of Literature:** The systematic literature review is limited to publications within a specific time frame, which may overlook earlier foundational studies that are still relevant.
2. **Case Study Selection:** The case studies focus on a select group of cities, which may not be representative of all smart city initiatives globally. Different cities may face unique challenges and contexts that affect their IoT implementations.
3. **Subjectivity in Analysis:** The thematic analysis is subject to interpretation, and different researchers may draw varying conclusions based on the same data.

4. Case Studies

This section presents a detailed analysis of three prominent case studies—Barcelona, Singapore, and Amsterdam—highlighting how these cities have successfully implemented Internet of Things (IoT) technologies to enhance urban sustainability. Each case study focuses on specific IoT applications, the challenges encountered, and the overall impact on urban living and resource management.

4.1 Barcelona, Spain

Barcelona is widely recognized as a pioneer in implementing smart city initiatives through IoT technologies. The city's approach focuses on improving urban services, enhancing citizen engagement, and promoting sustainability.

IoT Applications: Barcelona has developed a comprehensive smart city framework that integrates various IoT applications, including:

- **Smart Lighting:** The city has installed smart streetlights equipped with sensors that adjust brightness based on pedestrian and vehicular traffic. This not only reduces energy consumption but also enhances safety.

- **Smart Waste Management:** The introduction of smart bins equipped with sensors allows for real-time monitoring of waste levels. The system optimizes waste collection routes, reducing operational costs and emissions.
- **Smart Water Management:** IoT sensors monitor water quality and consumption, enabling the city to detect leaks and manage resources more efficiently.

Challenges: Despite its successes, Barcelona faces challenges related to data privacy and the need for comprehensive digital literacy programs to ensure all residents benefit from smart initiatives. The integration of multiple IoT systems also requires significant investment and coordinated governance.

Impact on Sustainability: The implementation of these IoT solutions has resulted in reduced energy consumption, lower carbon emissions, and improved waste management efficiency. Additionally, citizen engagement initiatives, such as mobile applications that allow residents to report issues, have fostered a sense of community involvement in sustainability efforts.

4.2 Singapore

Singapore has emerged as a global leader in smart city development, driven by its commitment to using technology to enhance urban living and sustainability.

IoT Applications: The Smart Nation initiative in Singapore encompasses a wide array of IoT applications:

- **Smart Traffic Management:** The city employs an extensive network of sensors and cameras to monitor traffic patterns and optimize traffic flow. This includes real-time data sharing with commuters through mobile applications.
- **Smart Energy Grids:** Singapore's energy management system uses IoT technologies to monitor and optimize energy consumption across residential and commercial buildings, promoting energy efficiency.
- **Environmental Monitoring:** The city utilizes IoT sensors to monitor air quality and weather conditions, providing real-time data to residents and policymakers.

Challenges: Singapore's challenges include ensuring data privacy and cybersecurity amidst extensive data collection efforts. Additionally, the need for continual investment in infrastructure and technology updates poses financial challenges.

Impact on Sustainability: The integration of IoT in Singapore has led to improved traffic flow, reduced energy consumption, and enhanced public health through better air quality monitoring. The city's commitment to sustainability is evident in its ambitious targets for reducing greenhouse gas emissions and promoting green building initiatives.

4.3 Amsterdam, Netherlands

Amsterdam has embraced IoT technologies to enhance urban living and sustainability through its Smart City program, which emphasizes citizen participation and innovation.

IoT Applications: Key IoT initiatives in Amsterdam include:

- **Smart Energy Solutions:** The city has implemented smart grids and energy storage systems that allow residents to monitor and manage their energy usage effectively. Initiatives encourage the use of renewable energy sources.
- **Smart Mobility:** Amsterdam promotes smart mobility solutions through connected public transportation systems and bike-sharing programs. IoT technology enables real-time tracking of public transport, enhancing user experience.
- **Waste Management:** Similar to Barcelona, Amsterdam employs smart bins that provide data on waste levels, optimizing collection routes and reducing operational costs.

Challenges: Amsterdam faces challenges related to maintaining citizen trust in data security and privacy. Additionally, integrating new technologies with existing urban infrastructure requires careful planning and investment.

Impact on Sustainability: The city's smart initiatives have led to a significant reduction in energy consumption, improved public transport efficiency, and enhanced waste management. Citizen participation in decision-making processes fosters a sense of ownership and accountability for sustainable practices.

5. Discussion

The findings from the literature review and case studies highlight the transformative potential of Internet of Things (IoT) technologies in shaping smart cities and promoting urban sustainability. This section discusses the implications of these findings, the challenges that cities face in adopting IoT solutions, and the future direction for urban sustainability initiatives leveraging IoT technologies.

5.1 Implications of IoT on Urban Sustainability

The integration of IoT technologies in urban environments presents significant opportunities for enhancing sustainability. The case studies of Barcelona, Singapore, and Amsterdam illustrate how IoT applications can lead to improved resource management, enhanced citizen engagement, and increased operational efficiency. For example, smart waste management systems not only reduce collection costs but also promote recycling and waste reduction

initiatives by providing real-time data on waste levels. Similarly, smart traffic management systems can alleviate congestion, reducing emissions and improving air quality.

The effectiveness of these technologies underscores the importance of data-driven decision-making in urban planning and management. Cities that leverage IoT data to inform policy decisions can respond more effectively to emerging challenges, such as climate change and population growth. The ability to monitor environmental conditions and resource consumption in real-time allows for proactive interventions that support sustainability goals.

5.2 Challenges to IoT Integration

Despite the benefits, the integration of IoT technologies in smart cities is fraught with challenges. Data privacy and security concerns are paramount, as extensive data collection can lead to potential misuse of personal information. As cities increasingly rely on interconnected systems, the risk of cyberattacks also escalates, highlighting the need for robust cybersecurity measures.

Moreover, the digital divide presents a significant barrier to the equitable implementation of IoT solutions. Vulnerable populations may lack access to technology or the skills necessary to engage with smart city initiatives, potentially exacerbating existing inequalities. It is crucial for city planners and policymakers to prioritize inclusivity in their smart city strategies, ensuring that all residents can benefit from technological advancements.

5.3 The Role of Collaboration and Governance

The successful implementation of IoT in smart cities requires effective collaboration among stakeholders, including government agencies, private sector partners, and citizens. The case studies demonstrate that cities with strong governance frameworks and active citizen engagement are more likely to achieve their sustainability goals. Collaborative initiatives that involve community input not only enhance the legitimacy of smart city projects but also foster a sense of ownership among residents.

Furthermore, establishing clear regulatory frameworks is essential for addressing data privacy and security concerns. Policymakers must create guidelines that promote transparency in data collection and use while safeguarding citizens' rights. Engaging with citizens in the development of these frameworks can help build trust and acceptance of smart city initiatives.

5.4 Future Directions for Research and Practice

Looking ahead, several future directions emerge for the research and practice of IoT in smart cities:

1. **Exploring New Technologies:** As IoT continues to evolve, it will be essential to explore emerging technologies, such as artificial intelligence (AI) and blockchain, that can enhance the capabilities of smart city systems. Integrating these technologies could further improve data analysis, decision-making, and security.
2. **Longitudinal Studies:** More longitudinal studies are needed to assess the long-term impacts of IoT implementations on urban sustainability. Understanding how these technologies evolve over time and their sustained effects on cities will provide valuable insights for future initiatives.
3. **Cross-Cultural Comparisons:** Research that examines IoT applications in diverse cultural and economic contexts can identify best practices and common challenges, contributing to a more comprehensive understanding of smart city dynamics globally.
4. **Addressing the Digital Divide:** Future research should focus on strategies to bridge the digital divide, ensuring that smart city initiatives are inclusive and accessible to all residents, particularly marginalized communities.

6. Conclusion

This research paper has explored the transformative role of Internet of Things (IoT) technologies in advancing urban sustainability within smart cities. Through a systematic literature review and detailed case studies of Barcelona, Singapore, and Amsterdam, it is evident that IoT applications are critical in enhancing resource management, improving urban services, and fostering citizen engagement. These technologies provide cities with the tools needed to address complex urban challenges such as resource depletion, traffic congestion, and environmental degradation.

The findings indicate that successful IoT implementation in smart cities not only leads to operational efficiencies but also contributes significantly to sustainability goals. For instance, smart waste management systems optimize collection processes and encourage recycling, while smart traffic solutions reduce congestion and enhance air quality. Moreover, the case studies demonstrate that the involvement of citizens in smart city initiatives is vital for building trust and ensuring that these technologies serve the needs of all community members.

However, the integration of IoT technologies is not without challenges. Issues of data privacy and security, as well as the potential for exacerbating existing inequalities, must be addressed to ensure equitable access to the benefits of smart city innovations. Policymakers must

prioritize inclusive strategies that engage all stakeholders and develop robust regulatory frameworks to safeguard citizen rights.

In conclusion, while the journey towards fully realizing the potential of IoT in smart cities is ongoing, the evidence presented in this paper underscores the significant promise that these technologies hold for achieving urban sustainability. Future research should continue to investigate innovative IoT applications, explore cross-cultural comparisons, and focus on strategies to bridge the digital divide. By doing so, cities can harness the power of IoT to create more sustainable, resilient, and livable urban environments for generations to come.

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AI-Driven Fraud Detection in Banking: A Machine Learning Approach

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Abstract

This white paper delves into the use of machine learning and AI algorithms for detecting fraudulent activities in banking, including real-time transaction monitoring, anomaly detection, and pattern recognition.

Highlight the role of AI in enhancing cybersecurity in the banking sector, including AI-based fraud detection, threat monitoring, and predictive analytics to prevent data breaches and cyberattacks.

Investigate the growing role of AI in optimizing trading strategies, improving market predictions, and managing portfolios. It can also explore the challenges of using AI in high-frequency trading and algorithmic strategies.

Challenges in Traditional Fraud Detection: Limitations of rule-based systems, reliance on historical data, and human oversight.

Keywords:

Machine Learning Algorithms, Artificial Intelligence (AI), Fraudulent Activities Detection, Real-time Transaction Monitoring, Anomaly Detection

1. Introduction

The rapid evolution of digital banking and online financial transactions has led to a significant increase in fraudulent activities. With millions of transactions occurring daily, cybercriminals continuously develop more sophisticated techniques to exploit vulnerabilities in financial systems. “The sheer volume of digital transactions, coupled with the complexity of emerging threats, has rendered traditional fraud detection methods insufficient in providing robust protection.

While rule-based fraud detection systems have been effective in the past, they often struggle to keep up with modern threats due to their reliance on static rules and historical data. These

systems typically detect fraud by flagging transactions that violate predefined conditions, such as large fund transfers, multiple transactions within a short timeframe, or transactions from unusual locations. However, such approaches are prone to high false positive rates, which lead to unnecessary transaction declines and customer frustration. Furthermore, fraudsters are constantly evolving their tactics, finding ways to bypass traditional security measures by mimicking legitimate transactions.

This evolving landscape has necessitated the adoption of Artificial Intelligence (AI) and Machine Learning (ML) technologies, which bring advanced fraud detection capabilities through real-time monitoring, predictive analytics, and adaptive learning. AI-powered fraud detection systems continuously evolve, adapting to new fraud patterns and detecting anomalies more efficiently than conventional methods.

Objectives of the Chapter:

1. Examine the role of AI and ML in detecting banking fraud.
2. Explore real-time monitoring and anomaly detection techniques.
3. Highlight cybersecurity improvements through AI-driven solutions.
4. Analyze AI's impact on trading strategies and financial market predictions.
5. Discuss the challenges and limitations of AI-based fraud detection systems.

2. Traditional Fraud Detection Methods: Limitations and Challenges

Traditional fraud detection methods have long served as the foundation for financial security in digital banking and online transactions. These methods primarily rely on rule-based systems, static thresholds, and historical transaction analysis to detect fraudulent activities. While these approaches have been effective to some extent, they struggle to keep up with the increasing sophistication of cybercriminal tactics, rapidly evolving fraud patterns, and the vast volume of transactions processed daily.

1. Dependence on Rule-Based Systems

Traditional fraud detection relies on predefined rule-based systems that flag suspicious transactions based on specific criteria. These rules may include:

- Transactions exceeding a predefined amount.
- Multiple transactions from different geographic locations within a short period.
- Repeated login attempts from an unrecognized device.

Limitations:

- **Lack of Adaptability:** Rule-based systems struggle to keep pace with new fraud tactics, requiring constant manual updates to remain effective.
- **High False Positives:** Legitimate transactions often get flagged as fraud due to rigid rules, leading to customer inconvenience.
- **Inefficiency Against Emerging Threats:** Fraudsters can easily bypass static rules by slightly modifying their attack strategies, rendering rule-based systems ineffective.

2. Reliance on Historical Data

Many traditional fraud detection methods rely heavily on historical transaction data to identify fraudulent behavior. These systems compare incoming transactions against past patterns to detect anomalies.

Limitations:

- **Inability to Detect New Fraud Techniques:** Since historical data only reflects past fraud cases, these systems fail to recognize novel fraud schemes that have not been encountered before.
- **Delayed Detection:** Fraudulent activities often go undetected until a suspicious pattern emerges in transaction history, allowing fraudsters time to execute multiple fraudulent transactions before being caught.

3. High False Positive and False Negative Rates

Traditional fraud detection systems often struggle with accuracy, resulting in high false positives (legitimate transactions flagged as fraud) and false negatives (fraudulent transactions going undetected).

False Positives (Legitimate Transactions Blocked)

- Causes frustration for customers, leading to a poor user experience.
- Increases customer service workload as users dispute blocked transactions.
- Can result in financial loss for businesses due to declined legitimate transactions.

False Negatives (Fraudulent Transactions Approved)

- Allows cybercriminals to bypass detection and commit fraud.
- Results in financial losses for both banks and customers.

- Damages the reputation of financial institutions if fraud goes unchecked.

4. Ineffectiveness Against Sophisticated Fraud Techniques

Modern fraudsters use advanced tactics such as social engineering, identity theft, account takeovers, and AI-driven attacks to bypass traditional fraud detection systems.

Limitations:

- **Account Takeover (ATO) Fraud:** Traditional systems may fail to detect when a fraudster **gains access to a legitimate user's account** and conducts transactions that appear normal.
- **Phishing and Social Engineering Attacks:** These attacks trick users into willingly providing their credentials, making it difficult for rule-based systems to detect fraud.
- **Synthetic Identity Fraud:** Fraudsters create **fake identities** using a mix of real and fictitious data to open fraudulent accounts that bypass traditional security checks.

5. Slow Response Times and Manual Investigations

Traditional fraud detection relies heavily on manual reviews by fraud analysts, which slows down the process of identifying and responding to threats.

Limitations:

- **Time-Consuming Investigations:** Manual fraud investigations take time, allowing fraudsters to complete multiple transactions before being detected.
- **Resource-Intensive:** Banks must employ large fraud investigation teams to handle flagged transactions.
- **Scalability Issues:** As digital banking grows, traditional systems cannot scale efficiently to handle millions of transactions per second.

6. Lack of Real-Time Fraud Prevention

Many traditional fraud detection systems operate in batch processing mode, meaning transactions are analyzed after they are completed. This delay allows fraudsters to complete fraudulent transactions before detection occurs.

Limitations:

- **Inability to Stop Fraud in Real-Time:** Without real-time monitoring, fraudulent transactions may only be flagged after the funds have been transferred.

- **Slow Adaptation to New Fraud Schemes:** Fraud patterns evolve quickly, making after-the-fact detection less effective.

7. Difficulty in Handling Large-Scale Digital Transactions

With the rise of e-commerce, mobile banking, and cryptocurrency transactions, financial institutions process billions of transactions daily. Traditional fraud detection methods struggle with scalability due to their reliance on manual processes and static rules.

Limitations:

- **Limited Processing Power:** Rule-based systems struggle to analyze vast amounts of transaction data in real-time.
- **Inability to Detect Cross-Channel Fraud:** Many fraudsters operate across multiple platforms (e.g., bank accounts, credit cards, mobile wallets), making it harder for traditional systems to detect coordinated attacks.

The Need for AI-Driven Fraud Detection

Given the limitations of traditional fraud detection methods, financial institutions are increasingly adopting Artificial Intelligence (AI) and Machine Learning (ML) to enhance fraud prevention. AI-powered fraud detection systems:

- Continuously learn and adapt to new fraud patterns.
- Monitor transactions in real-time to detect and prevent fraud before completion.
- Reduce false positives and false negatives, improving accuracy.
- Identify anomalous behavior beyond rule-based thresholds.
- Scale efficiently to process billions of digital transactions.

As cyber threats evolve, AI-driven fraud detection will be critical in securing online transactions and protecting financial institutions from increasingly sophisticated fraud schemes.

3. AI and Machine Learning in Fraud Detection

The rapid growth of digital transactions, online banking, and e-commerce has led to an increased risk of fraudulent activities. Traditional fraud detection methods, which rely on predefined rules and static data analysis, often struggle to keep up with evolving fraud tactics and the scale of financial transactions. As fraudsters develop more sophisticated techniques,

financial institutions must adopt Artificial Intelligence (AI) and Machine Learning (ML) to strengthen fraud detection systems.

AI and ML provide advanced capabilities such as real-time monitoring, adaptive learning, and predictive analytics, allowing financial institutions to detect fraudulent activities with greater accuracy”. These technologies reduce false positives, identify emerging fraud patterns, and enhance security in financial transactions.

3.1 How AI and Machine Learning Enhance Fraud Detection

AI and ML models analyze vast amounts of transaction data, detecting anomalous behavior, suspicious patterns, and high-risk activities in real time. The fraud detection process typically involves the following steps:

Step 1: Data Collection and Preprocessing

AI systems gather transaction data from various sources, including:

- **Bank transactions** (credit card purchases, wire transfers, deposits, withdrawals).
- **User behavior** (login times, device information, IP addresses, geolocation).
- **Historical fraud records** (previous fraudulent transactions).

ML algorithms preprocess and clean this data, removing noise and inconsistencies before feeding it into fraud detection models.

Step 2: Anomaly Detection and Risk Scoring

ML models establish a baseline of normal user behavior and detect deviations. If a transaction deviates significantly from established patterns, it is flagged as potential fraud.

Example: If a customer who normally spends \$500 per month suddenly makes a \$10,000 purchase from an unfamiliar location, the system flags it for further review.

Each transaction is assigned a risk score, helping financial institutions determine whether to approve, decline, or request additional authentication.

Step 3: Real-Time Monitoring and Decision Making

AI-powered fraud detection operates in real time, ensuring that high-risk transactions are blocked before they are completed. The system:

- Flags transactions exceeding predefined risk thresholds.
- Automatically requests multi-factor authentication (MFA) if necessary.
- Blocks transactions from compromised accounts or devices.

Step 4: Continuous Learning and Model Updates

AI models continuously learn from new fraud attempts, improving detection over time. This ensures that fraudsters cannot exploit the same loopholes repeatedly.

Example: A fraudster may use stolen card details to make small, unnoticeable transactions. If AI detects this pattern, it adjusts its fraud detection model to identify similar fraud attempts in the future.

3.2 Advantages of AI and Machine Learning in Fraud Detection

1. Real-Time Fraud Prevention

- Traditional fraud detection methods identify fraud after it has occurred. AI prevents fraud before a transaction is completed.
- Real-time transaction monitoring blocks suspicious activities instantly, preventing financial losses.

2. Higher Accuracy and Reduced False Positives

- Traditional systems often flag legitimate transactions as fraud (false positives), causing inconvenience to customers.
- AI analyzes multiple factors (transaction history, user behavior, device fingerprinting) to accurately assess fraud risk.

3. Adaptive Learning and Evolving Fraud Detection

- AI continuously learns and adapts to new fraud techniques, unlike static rule-based systems.
- ML models evolve with new fraud trends, phishing attacks, and account takeovers, ensuring up-to-date fraud protection.

4. Multi-Layered Fraud Detection Using Big Data Analytics

- AI analyzes huge volumes of financial data, detecting complex fraud patterns that human analysts might miss.
- Fraud detection models use deep learning, behavioral analytics, and anomaly detection for multi-layered security.

5. Cost-Effective Fraud Detection

- AI automates fraud detection, reducing the need for manual reviews by fraud analysts.

- Reduces operational costs while improving fraud prevention efficiency.

6. Enhanced User Authentication with Behavioral Biometrics

- AI integrates biometric authentication (fingerprints, facial recognition, keystroke dynamics) to verify users securely.
- Detects anomalies in user behavior, preventing account takeovers.

3.3 AI and Machine Learning Techniques Used in Fraud Detection

Machine Learning Model	How It Works	Use Case in Fraud Detection
Supervised Learning	Trained on labeled fraud and non-fraud data	Detects known fraud patterns
Unsupervised Learning	Identifies anomalies without labeled fraud cases	Detects new and emerging fraud tactics
Neural Networks (Deep Learning)	Learns complex patterns in large datasets	Identifies synthetic identity fraud
Random Forest & Decision Trees	Uses multiple decision pathways to classify transactions	Flags high-risk transactions
Natural Language Processing (NLP)	Analyzes text data for phishing attempts	Detects fraudulent emails and scam messages

3.4 Real-World Applications of AI in Fraud Detection

1. AI in Banking and Financial Fraud Detection

- Banks like JPMorgan Chase, Citibank, and HSBC use AI-driven fraud detection.
- AI models analyze millions of transactions daily to identify fraud patterns.

2. AI for Credit Card Fraud Prevention

- AI detects card-not-present (CNP) fraud in online transactions.
- Visa and Mastercard use AI to analyze cardholder spending behavior and flag unusual transactions.

3. AI in E-Commerce and Payment Fraud Prevention

- Online retailers like Amazon and eBay use AI to detect fake reviews, fraudulent returns, and chargeback fraud.
- AI models flag suspicious refund requests and fake accounts created for scams.

4. AI in Cryptocurrency and Blockchain Fraud Detection

- AI tracks suspicious crypto transactions, preventing fraud in digital wallets and exchanges.
- Fraud detection tools like Chainalysis and CipherTrace use AI to detect illegal activities on the blockchain.

3.5 Challenges of AI in Fraud Detection

- **Data Privacy Concerns:** AI requires access to sensitive transaction data, raising privacy issues.
- **Bias in AI Models:** AI may incorrectly flag certain demographics if trained on biased datasets.
- **Evolving Fraud Techniques:** Cybercriminals develop AI-resistant fraud strategies, requiring continuous AI model updates.
- **False Negatives:** AI may fail to detect highly sophisticated fraud attempts, allowing fraud to occur.

Financial institutions must ensure AI fraud detection systems are transparent, ethical, and compliant with regulations like GDPR and CCPA.

AI and Machine Learning have transformed fraud detection, offering real-time monitoring, predictive analytics, and adaptive learning to combat cyber threats. By integrating deep learning, behavioral biometrics, and big data analytics, AI enhances fraud prevention while reducing false positives and improving security.

4. Enhancing Cybersecurity Through AI

Cybersecurity threats have become more sophisticated and widespread as organizations and individuals increasingly rely on digital platforms for financial transactions, communications, and data storage. “Traditional cybersecurity measures, such as rule-based firewalls and signature-based antivirus software, struggle to keep pace with evolving cyber threats such as ransomware, phishing, advanced persistent threats (APTs), and zero-day exploits.

Artificial Intelligence (AI) has emerged as a transformative force in cybersecurity, enabling real-time threat detection, predictive analytics, automated responses, and enhanced risk

management. By leveraging Machine Learning (ML), Natural Language Processing (NLP), and behavioral analytics, AI enhances cybersecurity by identifying, analyzing, and mitigating threats more effectively than traditional methods.

4.2 The Role of AI in Cybersecurity

AI strengthens cybersecurity by continuously learning from vast amounts of data, detecting patterns, anomalies, and malicious activities that may go unnoticed by human analysts. Unlike traditional security tools that rely on predefined rules, AI-driven security systems adapt to new threats in real time, providing proactive protection.

1. Real-Time Threat Detection and Response

- AI-powered security systems monitor network traffic in real time, detecting unusual patterns and preventing cyberattacks before they cause harm.
- AI-driven Intrusion Detection Systems (IDS) and Intrusion Prevention Systems (IPS) analyze packet flows, login attempts, and network behaviors to identify cyber threats.

2. Anomaly Detection and Behavior Analysis

- AI identifies deviations from normal user behavior, flagging potential security breaches.
- Behavioral analytics monitor login patterns, keystroke dynamics, and device activity to detect compromised accounts.

3. Predictive Threat Intelligence and Risk Assessment

- AI uses predictive analytics to assess potential cyber threats before they occur by analyzing historical attack data.
- Threat intelligence platforms use AI-powered threat modeling to anticipate and mitigate risks.

4. Automated Incident Response and Threat Mitigation

- AI-powered Security Orchestration, Automation, and Response (SOAR) platforms automate cybersecurity responses, reducing the need for manual intervention.
- AI quarantines infected devices, blocks malicious IP addresses, and isolates threats without human intervention.

5. Phishing and Email Fraud Detection

AI-powered Natural Language Processing (NLP) analyzes email content to detect phishing attempts, impersonation fraud, and scam emails.

AI flags suspicious emails by detecting fake URLs, urgent language, and malicious attachments.

6. Malware and Ransomware Detection

- AI-powered antivirus and endpoint security solutions detect and neutralize malware before it executes.
- AI analyzes malware behavior rather than relying on traditional signature-based detection, making it effective against zero-day attacks.

7. Enhancing Identity and Access Management (IAM)

- AI strengthens user authentication through biometric security, multi-factor authentication (MFA), and behavioral analytics.
- AI-powered adaptive authentication adjusts security levels based on user risk profiles.

4.3 Benefits of AI in Cybersecurity

1. Faster Threat Detection and Response

- AI identifies cyber threats in, reducing the time needed to respond to attacks.
- Traditional security systems may take hours or days to detect a breach, allowing cybercriminals to cause damage.

2. Reduced False Positives and Enhanced Accuracy

- AI reduces false positives in cybersecurity alerts by analyzing multiple threat indicators.
- This ensures security teams focus on genuine threats rather than wasting time on harmless anomalies.

3. Continuous Learning and Adaptive Security

- AI-based security solutions evolve over time, improving their ability to detect new and emerging threats.
- Unlike rule-based security systems, AI adapts to zero-day vulnerabilities and advanced cyber threats.

4. Automation and Cost Efficiency

- AI automates security operations, reducing the need for large cybersecurity teams.
- Automating threat detection and response lowers operational costs while improving security efficiency.

5. Proactive Defense Against Advanced Attacks

- AI detects Advanced Persistent Threats (APTs) that remain hidden within networks for months.
- AI-powered deception techniques (e.g., honeypots, decoy networks) lure attackers and expose cyber threats.

4.4 AI-Powered Cybersecurity Tools and Applications

AI Cybersecurity Solution	Functionality	Use Case
AI-Powered Firewalls	Monitors network traffic and blocks threats in real time	Preventing DDoS attacks, malware intrusions
AI-Driven Endpoint Security	Detects malware and unauthorized device activity	Protecting workstations and mobile devices
AI-Based SIEM (Security Information & Event Management)	Analyzes security logs for attack patterns	Detecting cyber threats across IT infrastructure
AI in Cloud Security	Identifies unauthorized access and insider threats	Protecting cloud data and SaaS applications
AI-Powered Fraud Detection	Analyzes banking transactions for fraud	Preventing financial scams and identity theft

4.5 Real-World Applications of AI in Cybersecurity

1. AI in Banking and Financial Security

- Banks use AI for fraud detection, identity verification, and transaction monitoring.
- AI-powered biometric authentication prevents unauthorized access.

2. AI in Enterprise Cybersecurity

- Corporations deploy AI-driven threat intelligence platforms to secure sensitive data.
- AI identifies insider threats, privilege misuse, and unauthorized access attempts.

3. AI in Government and National Security

- Governments use AI for cyber defense, surveillance, and intelligence gathering.
- AI helps detect nation-state cyberattacks and cyber espionage threats.

4.6 Challenges and Risks of AI in Cybersecurity

- **AI-Powered Cyber Attacks** – Hackers use AI to automate and enhance cyberattacks (e.g., deepfake scams).
- **Data Privacy Concerns** – AI systems require access to sensitive user data, raising privacy issues.
- **Bias in AI Algorithms** – AI may misidentify threats if trained on biased datasets.
- **High Implementation Costs** – AI-driven cybersecurity solutions can be expensive for small organizations.

AI is revolutionizing cybersecurity by enhancing threat detection, automating responses, and strengthening risk management. AI-driven security solutions provide real-time monitoring, predictive analytics, and adaptive learning, offering a proactive defense against cyber threats.

5. AI in Trading and Market Optimization

Financial markets operate in a highly dynamic and data-driven environment where **speed, accuracy, and predictive analysis** are critical for success”. Traditional trading strategies, which rely on human expertise, fundamental analysis, and historical data, often struggle to adapt to rapid market fluctuations, high-frequency trading, and the growing complexity of global markets.

Artificial Intelligence (AI) and Machine Learning (ML) have revolutionized financial trading by enabling data-driven decision-making, automated trading strategies, and real-time market analysis. AI-powered trading systems can process vast amounts of market data, detect patterns, and execute trades faster and more accurately than human traders.

5.2 The Role of AI in Trading

AI enhances trading strategies by analyzing market trends, predicting price movements, and optimizing trade execution. AI-powered systems leverage big data, deep learning algorithms, and quantitative models to identify trading opportunities in real time.

1. Algorithmic Trading (Algo-Trading)

- AI-driven algorithmic trading systems execute buy and sell orders based on pre-defined conditions.
- These systems use high-frequency trading (HFT) to execute trades within milliseconds, maximizing profit opportunities.

2. Predictive Market Analytics

- AI uses machine learning models to analyze historical market data and predict future price movements.
- Predictive analytics help traders make informed decisions based on market sentiment, news, and technical indicators.

3. Sentiment Analysis and News-Based Trading

- AI-powered Natural Language Processing (NLP) scans financial news, earnings reports, and social media to determine market sentiment.
- AI traders adjust portfolios based on breaking news and global events that impact financial markets.

4. Risk Management and Portfolio Optimization

- AI assesses market risks, volatility, and economic indicators to optimize portfolio allocation.
- AI-driven risk models adjust trading strategies to minimize losses during market downturns.

5. Fraud Detection and Market Surveillance

- AI monitors trading activities to detect insider trading, market manipulation, and fraudulent transactions.
- Regulatory bodies and stock exchanges use AI for real-time surveillance of suspicious trading patterns.

5.3 Key Benefits of AI in Trading and Market Optimization

1. Increased Trading Speed and Efficiency

- AI-powered high-frequency trading (HFT) executes thousands of trades per second.
- Automated trading eliminates human delays and emotional biases, ensuring faster decision-making.

2. Enhanced Market Predictions and Accuracy

- AI analyzes millions of market data points to identify profitable trading opportunities.
- AI models continuously improve, adapting to changing market conditions.

3. Reduced Trading Risks and Losses

- AI predicts market downturns and adjusts trading strategies to minimize risks.
- AI-powered risk management algorithms optimize portfolio performance.

4. Better Market Liquidity and Price Stability

- AI-powered market-making ensures continuous liquidity, reducing price volatility.
- AI traders adjust prices in real time, improving market efficiency.

5. Cost Reduction and Automation

- AI automates trade execution, analysis, and reporting, reducing operational costs.
- Fewer human traders are required, lowering transaction costs for financial institutions.

5.4 AI-Driven Trading Strategies

AI Trading Strategy	Functionality	Use Case in Trading
High-Frequency Trading (HFT)	Executes thousands of trades per second	Capitalizing on small price fluctuations
Machine Learning-Based Trading	Predicts price movements using AI models	Optimizing buy/sell decisions
Sentiment Analysis Trading	Analyzes news and social media sentiment	Reacting to market-moving events
Reinforcement Learning Trading	AI adapts trading strategies based on rewards	Improving long-term portfolio performance
Risk-Based Portfolio Optimization	AI manages asset allocation to minimize risk	Enhancing investment diversification

5.5 Real-World Applications of AI in Trading

1. AI-Powered Hedge Funds

- Hedge funds use AI for market prediction, risk assessment, and trade execution.
- AI-driven funds outperform traditional funds by adjusting strategies based on market conditions.

2. AI in Stock Market Trading

- Stock traders use AI algorithms to analyze earnings reports, technical indicators, and trading patterns.
- AI enhances quantitative trading strategies by detecting price movements before human traders react.

3. AI in Cryptocurrency Trading

- AI-powered bots analyze crypto price trends, trading volumes, and blockchain activity.
- Crypto exchanges use AI for market-making, fraud detection, and automated trading.

4. AI in Forex Trading

- AI models analyze macroeconomic indicators, interest rates, and geopolitical events to forecast forex market trends.
- AI optimizes currency trading strategies by identifying arbitrage opportunities.

5.6 Challenges of AI in Trading

- **Market Manipulation Risks** – AI-driven high-frequency trading can create flash crashes and liquidity crises.
- **Overfitting in AI Models** – ML models may perform well on past data but fail in real-world trading conditions.
- **Regulatory and Ethical Concerns** – AI-driven trading raises concerns about fairness, transparency, and compliance.
- **Cybersecurity Threats** – AI trading systems are vulnerable to hacking, algorithmic exploitation, and AI-driven attacks.

Regulators and financial institutions must ensure responsible AI adoption while balancing market efficiency and risk management.

AI and Machine Learning have revolutionized trading and market optimization, enabling faster, smarter, and more efficient trading strategies. AI-powered trading systems enhance market predictions, automate trade execution, and reduce trading risks, making them indispensable for modern financial markets.

As AI continues to evolve, its role in trading will expand, providing greater market insights, improved risk management, and enhanced decision-making capabilities. However, financial

institutions must address regulatory challenges, ethical concerns, and cybersecurity risks to ensure responsible AI-driven trading.

6. Challenges and Ethical Considerations

AI and Machine Learning (ML) have transformed industries such as finance, cybersecurity, healthcare, and trading, enabling faster decision-making and automation. However, their widespread adoption also introduces challenges and ethical concerns related to data privacy, security, transparency, and adversarial threats.

6.1 Data Privacy and Security

6.1.1 Privacy Risks in AI Systems

AI-driven systems rely on massive datasets containing personal, financial, and behavioral information to improve performance. However, improper handling of this data can lead to privacy violations and regulatory breaches.

- **Unauthorized Data Collection** – AI models often collect and store user data without explicit consent, raising concerns under GDPR, CCPA, and other data protection laws.
- **Re-identification Risks** – Even anonymized data can sometimes be reverse-engineered to identify individuals.
- **Surveillance Concerns** – AI-powered facial recognition and behavioral tracking may violate individual privacy rights.

6.1.2 Security Risks in AI Systems

AI models are vulnerable to cyberattacks, data breaches, and unauthorized access, posing risks to financial and cybersecurity applications.

- **Data Poisoning Attacks** – Attackers inject false data into AI training sets to mislead decision-making.
- **Model Inversion Attacks** – Hackers exploit AI models to extract sensitive user data.
- **Weak Encryption** – Poor security practices can expose AI algorithms, financial records, and trade secrets.

6.1.3 Solutions to Privacy and Security Risks

- **Federated Learning:** AI models learn from decentralized data sources without directly accessing raw data, enhancing privacy.
- **Differential Privacy:** AI adds mathematical noise to datasets, preventing individual identification.
- **Zero-Trust Security Models:** AI systems verify every user and device before granting access.
- **Regulatory Compliance:** AI should align with GDPR, CCPA, and financial data protection laws to protect users.

6.2 Model Interpretability ("Black Box" Problem)

6.2.1 Lack of Transparency in AI Decision-Making

Many AI models, especially deep learning and neural networks, function as black boxes, meaning their internal decision-making processes are not easily interpretable.

Key Issues with Model Interpretability:

- **Regulatory Challenges** – In finance and cybersecurity, organizations must explain AI decisions to regulators and users.
- **Trust and Accountability** – If AI incorrectly flags a legitimate banking transaction, users should understand why.
- **Ethical Concerns** – Biased AI decisions (e.g., loan denials or fraud detection errors) can harm individuals unfairly.

6.2.2 Explainable AI (XAI) Solutions

- **Decision Trees and Rule-Based Models:** Offer transparent decision-making, making AI easier to audit.
- **SHAP (Shapley Additive Explanations):** Breaks down AI decisions to show which factors influenced an outcome.
- **LIME (Local Interpretable Model-Agnostic Explanations):** Generates simplified models to explain complex AI behavior.

6.2.3 Balancing Performance and Interpretability

- **Trade-Off Between Accuracy and Explainability:** Some AI models, like deep learning, achieve high accuracy but lack interpretability, while simpler models (e.g., decision trees) are easier to understand but may be less effective.

- **Hybrid AI Models:** Combining machine learning and rule-based systems improves both transparency and performance.

6.3 Adversarial Attacks on AI Systems

6.3.1 What Are Adversarial Attacks?

Adversarial attacks occur when hackers manipulate AI models by introducing small, deceptive inputs that cause incorrect predictions. These attacks pose a significant threat to AI applications in finance, cybersecurity, and trading.

Types of Adversarial Attacks:

- **Evasion Attacks** – Attackers modify input data to trick AI into misclassifying fraud, malware, or transactions.
- **Data Poisoning Attacks** – Hackers inject false training data, corrupting AI decision-making.
- **Model Extraction Attacks** – Attackers reverse-engineer AI models to steal intellectual property.

6.3.2 Impact of Adversarial Attacks in Finance and Cybersecurity

- **Financial Losses:** Attackers trick AI-driven stock trading bots into making incorrect investments.
- **Data Breaches:** AI cybersecurity tools misclassify malware, allowing hackers to bypass security defenses.
- **Identity Theft:** AI-powered facial recognition systems are fooled by adversarial attacks (e.g., deepfake manipulation).

6.3.3 Defending AI Against Adversarial Attacks

- **Adversarial Training:** AI models are trained with manipulated inputs to improve resistance to attacks.
- **AI Security Audits:** Regular vulnerability assessments ensure AI remains robust against cyber threats.
- **Anomaly Detection Systems:** AI monitors its own behavior, detecting suspicious modifications in real time.

7. Conclusion

Artificial Intelligence (AI) and Machine Learning (ML) have revolutionized industries such as finance, cybersecurity, and trading, enabling faster decision-making, enhanced security, and optimized market operations. AI-powered systems have significantly improved fraud detection, real-time risk assessment, automated trading, and cybersecurity defenses, making them indispensable in today's digital world.

However, despite its advantages, AI comes with challenges and ethical considerations that organizations must address. Data privacy concerns, AI bias, model interpretability issues, and adversarial attacks present significant risks to AI-driven systems. Ensuring transparent, secure, and fair AI models is crucial for maintaining user trust and regulatory compliance.

To fully harness the potential of AI, industries must:

- Implement privacy-preserving AI techniques (e.g., federated learning, differential privacy).
- Develop explainable AI (XAI) models to improve transparency and trust.
- Strengthen adversarial defense mechanisms to protect AI from cyber threats.
- Comply with global AI regulations and ethical AI standards to prevent misuse.
- Maintain human oversight in AI-driven decision-making to reduce errors and biases.

As AI continues to evolve, its role in fraud prevention, market optimization, and cybersecurity will expand. However, responsible AI governance, ethical considerations, and security best practices must remain a priority to ensure AI-driven systems are fair, transparent, and resilient against emerging threats.

By adopting AI responsibly, businesses and institutions can unlock greater efficiency, security, and innovation, paving the way for a future where AI enhances decision-making while upholding ethical and security standards

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The Evolution of Kinds of Intelligencer

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Abstract

This research paper explores the evolving understanding of human intelligence, moving beyond traditional notions of general intelligence (IQ) to encompass a broader spectrum of cognitive and cognitive-related abilities. It examines various theoretical frameworks, including multiple intelligences, emotional intelligence, social intelligence, and creative intelligence, and discusses their implications for education, personal development, and societal understanding. The paper analyzes the impact of technological advancements, such as artificial intelligence and neuroscience, on our understanding of intelligence and its evolution.

Keywords:

Intelligence, Multiple Intelligences, Emotional Intelligence, Social Intelligence, Creative Intelligence, General Intelligence (IQ), Cognitive Abilities, Neuroscience, Artificial Intelligence, Education, Personal Development, Human Potential

Introduction

For centuries, the concept of intelligence has been a subject of intense philosophical and scientific inquiry. Traditionally, intelligence was largely equated with general intelligence (g), often measured by IQ tests, which primarily assess logical-mathematical and linguistic abilities. However, this narrow view has been challenged by a growing body of research that suggests a more nuanced and multifaceted understanding of human intelligence. This paper explores the evolution of our understanding of intelligence, moving beyond the limitations of traditional IQ tests to encompass a broader spectrum of cognitive and cognitive-related abilities.

Overview of Literature

The historical perspective on intelligence has evolved significantly. Early theories, influenced by psychometrics, focused on general intelligence as a single, unitary trait. However, this view has been challenged by:

- **Multiple Intelligences Theory:** Proposed by Howard Gardner, this theory suggests that intelligence encompasses a range of distinct abilities, including linguistic, logical-mathematical, spatial, bodily-kinesthetic, musical, interpersonal, intrapersonal, naturalistic, and existential intelligences. This framework emphasizes the importance of recognizing and nurturing diverse talents and abilities.
- **Emotional Intelligence (EI):** Developed by John Mayer and Peter Salovey, EI encompasses the ability to understand, manage, and utilize emotions effectively. It includes self-awareness, self-regulation, empathy, and social skills.
- **Social Intelligence:** This concept refers to an individual's ability to navigate social situations effectively, build and maintain relationships, and understand and respond appropriately to social cues.
- **Creative Intelligence:** This aspect of intelligence emphasizes the ability to generate novel ideas, think outside the box, and find innovative solutions to problems.

Research Methodology

This research paper primarily relies on a comprehensive literature review of existing research articles, books, and reports on intelligence, cognitive psychology, and related fields. Relevant databases, such as PubMed, Google Scholar, JSTOR, and PsycINFO, were searched using appropriate keywords, including "intelligence," "multiple intelligences," "emotional intelligence," "social intelligence," "creative intelligence," "cognitive abilities," "neuroscience," "artificial intelligence," "education," and "human potential." The collected data was analyzed to identify key trends, challenges, and future directions in the study of human intelligence.

Evolution of Intelligence Theories

- **From General Intelligence to Multiple Intelligences:** The concept of general intelligence, often measured by IQ tests, has been challenged by the emergence of the Multiple Intelligences theory. This theory, proposed by Howard Gardner, suggests that intelligence is not a single, unitary trait but rather a multifaceted construct encompassing various distinct abilities. These include:
 - **Linguistic Intelligence:** The ability to use language effectively, including speaking, writing, and understanding.
 - **Logical-Mathematical Intelligence:** The ability to reason logically, solve mathematical problems, and think abstractly.

- **Spatial Intelligence:** The ability to visualize and manipulate objects in space, such as in art, architecture, and navigation.
- **Bodily-Kinesthetic Intelligence:** The ability to use one's body effectively, including skills in sports, dance, and crafts.
- **Musical Intelligence:** The ability to perceive, create, and understand music.
- **Interpersonal Intelligence:** The ability to understand and interact effectively with other people.
- **Intrapersonal Intelligence:** The ability to understand oneself, including one's own emotions, strengths, and weaknesses.
- **Naturalistic Intelligence:** The ability to understand and interact with the natural world.
- **Existential Intelligence:** The ability to ponder philosophical questions about life, death, and the meaning of existence.
- **The Rise of Emotional Intelligence:** Emotional intelligence (EI) has emerged as a significant area of research in recent decades. EI encompasses the ability to understand, manage, and utilize emotions effectively. Key components of EI include:
 - **Self-awareness:** The ability to recognize and understand one's own emotions.
 - **Self-regulation:** The ability to manage one's emotions and impulses effectively.
 - **Social skills:** The ability to build and maintain healthy relationships, communicate effectively, and navigate social situations.
 - **Empathy:** The ability to understand and share the feelings of others.
- **The Importance of Social Intelligence:** Social intelligence refers to an individual's ability to navigate social situations effectively, build and maintain relationships, and understand and respond appropriately to social cues. It encompasses a wide range of skills, including:
 - **Social perception:** The ability to accurately read and interpret social cues, such as facial expressions, body language, and tone of voice.
 - **Social cognition:** The ability to understand and predict the behavior of others.

- **Social skills:** The ability to effectively communicate, build rapport, and resolve conflicts.
- **The Role of Creative Intelligence:** Creative intelligence emphasizes the ability to generate novel ideas, think outside the box, and find innovative solutions to problems. It involves:
 - **Divergent thinking:** The ability to generate multiple solutions to a problem.
 - **Flexibility:** The ability to adapt to new situations and think in new ways.
 - **Originality:** The ability to generate unique and novel ideas.
 - **Elaboration:** The ability to develop and refine ideas.

The Impact of Technology on Intelligence

Technological advancements are significantly impacting our understanding and definition of intelligence.

- **Artificial Intelligence (AI):** The development of AI systems is pushing the boundaries of our understanding of intelligence. AI algorithms are capable of performing complex tasks, such as learning, problem-solving, and decision-making, that were once thought to be uniquely human. This raises questions about the nature of intelligence and the relationship between human and artificial intelligence.
- **Neuroscience and Brain Research:** Advances in neuroscience are providing new insights into the neural mechanisms underlying intelligence. Brain imaging techniques, such as fMRI and EEG, are allowing researchers to study brain activity in real-time and gain a deeper understanding of how the brain processes information, learns, and creates.
- **The Digital Age:** The digital age has created new forms of intelligence, such as digital literacy, computational thinking, and the ability to navigate and utilize digital information effectively. These skills are becoming increasingly important in today's interconnected world.

Implications for Education

The evolving understanding of intelligence has significant implications for education.

- **Personalized Learning:** By recognizing the diversity of intelligences, educators can create more personalized learning experiences that cater to the individual needs and strengths of each student.

- **Developing Multiple Intelligences:** Educational programs can be designed to nurture a wide range of intelligences, not just focusing on traditional academic skills like reading, writing, and mathematics.
- **Promoting Emotional and Social Intelligence:** Schools can play a crucial role in developing students' emotional and social intelligence, including skills such as empathy, communication, and conflict resolution.
- **Fostering Creativity and Innovation:** Educational environments should encourage creativity and innovation by providing students with opportunities for exploration, experimentation, and problem-solving.

Assessing Intelligence and Its Evolution

Many methods are used to assess intelligence and its evolution. These include (1) *behavioral measures*, which may involve naturalistic observation or analyzing responses in laboratory experiments, (2); *artifactual measures*, which involve analysis of tools, art, and so forth,; and (3) *anatomical/neurological measures*, which involve studies of the brain and cranium. Ideally, all three would converge upon a unified picture of how intelligence evolved. However, this is not always the case, and indeed, the assessment of intelligence is fraught with challenges.

An obvious one is that we cannot perform behavioral or neurological studies of our ancestors, so we are forced to rely on bones and artifacts. Moreover, the further back in time one looks, the more fragmentary the archaeological record becomes. To explore the ancestral roots of our intelligence, we therefore also partly rely on studying the intelligence and brains of the great apes, our closest biological relatives. We share a common ancestor with great apes as recently as 4–6 million years ago (mya): No living species are more closely related. Other species such as dolphins and crows share some complex intellectual abilities with great apes and humans, but their abilities probably evolved independently and operate differently. Dolphins' and crows' brains differ strikingly from ours, for instance, whereas great ape brains are exceptionally similar to ours Emery & Clayton, 2004; Hof, Chavis & Marino, 2005; MacLeod, 2004).

What the great apes offer to the study of the evolution of human intelligence is the best living model of the intelligence that existed in our common great ape ancestors before our unique evolutionary lineage, the hominins, diverged. Modern human intelligence evolved from earlier forms of intelligence in response to selective pressures generated by ancestral living conditions. Understanding its evolution therefore entails looking into the past for the

changes that occurred within the hominins – but also for earlier intellectual traits upon which the hominins built and the changes that led to their divergence from ancestral great apes. If we can identify complex behaviors that great apes share with humans but not with other nonhuman primates, then these behaviors and the intellectual qualities they imply may have been shared by our common ancestors.

To use great apes to contribute to understanding the evolution of human intelligence, especially inferring what intellectual capacities evolved uniquely in the hominins, we need to assess their intellectual ceiling, that is, their top adult-level capabilities near the human boundary. The intelligence of great apes is highly malleable and dependent on the developmental and learning history of the individual (Matsuzawa, Tomonaga & Tanaka, 2006; Parker & McKinney 1999; de Waal, 2001), as it is in humans. Conclusions about great ape cognition and comparisons with human cognition must therefore be made with care. In part because this care has not always been taken, the literature on how human intelligence evolved does not present as straightforward a picture as one might hope. Nevertheless, an integrated account is starting to emerge.

What Distinguishes Human From Nonhuman Intelligence?

Many have attempted to specify what marks the intellectual divide between humans and other species. Some follow Aristotle's proposal that it is reason (French, 1994), or symbolic thinking. Symbols are arbitrary signs with conventional meanings that are used to represent (stand for) other things or relationships between them, and that generally have conventionally accepted meanings. Another suggestion is that human intelligence is distinguished by the ability to develop complex, abstract, internally coherent systems of symbol use (Deacon, 1997). Others propose that it is creativity, such as is required to invent tools, or abilities associated with creativity, such as or cognitive fluidity (combining concepts or ideas, or adapting them to new contexts), or the ability to generate and understand analogies (Fauconnier & Turner, 2002; Mithen, 1996). Still other proposals single out key abilities for dealing with the social world, such as demonstration teaching, imitative learning, cooperative problem solving, or communicating about the past and future. A related proposal is that the divide owes to the onset of what Premack and Woodruff (1978) refer to as *theory of mind*—the capacity to reason about mental states of others (Mithen, 1998).

The more we learn about nonhuman intelligence, however, the more we find that abilities previously thought to be uniquely human are not. For example, it was thought until the 1960s that humans alone make tools. But then Jane Goodall (1963) found wild

chimpanzees making them. Later, several other species were found making tools too (Beck, 1980). Thus, ideas about what marks the boundary between human and nonhuman intelligence have undergone repeated revision.

Although a large gulf separates human abilities from those of other species, it is not as easy as we hoped to pinpoint in a word or two what distinguishes humans. That does not mean that a more complex explanation is not forthcoming. For example, it may be that it is not creativity per se that distinguishes human intelligence, but the proclivity to take existing ideas and adapt them to new contexts or to one's own unique circumstances – that is, to put one's own spin on them, such that they become increasingly complex. The question of what separates human intelligence from that of other species is a recurring theme that will be fleshed out in the pages that follow.

Intelligence in Our Closest Relatives: The Great Apes We now possess a rich body of data on great ape intelligence (Byrne, 1995; Gómez, 2004; Matsuzawa *et al.*, 2006; Parker & Gibson, 1990; Povinelli, 2000; Rumbaugh & Washburn, 2003; Russon, Bard & Parker, 1996; Tomasello & Call, 1997; de Waal, 2001). This section summarizes the current picture of great ape intelligence, focusing on qualities once thought to be uniquely human. While some monkeys have shown similar achievements, great apes consistently achieve higher levels (Parker & McKinney, 1999). Great apes have shown many social cognitive abilities thought uniquely human. They show imitative learning and demonstration teaching powerful enough to sustain simple cultures (Boesch, 1991; Byrne & Russon, 1998; Parker 1996; van Schaik *et al.*, 2003; Whiten *et al.*, 1999). Some have solved problems cooperatively (Boesch & Boesch-Achermann, 2000; Hirata & Fuwa, 2007) and show some understanding of others' mental states (e.g., knowledge, competence) (Parker & McKinney, 1999). Captives have acquired basic sign language, including learning and inventing arbitrary conventional signs and simple grammar (Blake, 2004). Some great ape gestures qualify as symbolic by standards used in early language studies, including tree-drumming, holding thumb and finger together and blowing through them to represent a balloon, and making twisting motions toward containers they wanted opened (Blake, 2004). Great apes can understand simple analogies and engage in analogical reasoning (Thompson & Oden, 2000). They are considered to achieve basic symbolic abilities in several problem domains; they can do simple arithmetic and master simple language, for example (Parker & McKinney, 1999; Thompson & Oden, 2000).

A certain degree of creativity may be normal in great apes (and other nonhuman species);

Reader & Laland 2003). Their creativity includes smearing leaf pulp foam on their body (perhaps as an analgesic), inventing new tools (e.g., branch hook tools, termite fishing brush tools), primitive swimming, and fishing (Russon et al., 2009; Sanz & Morgan, 2004). They have invented gestures and signs such as hand shaking and tree drumming (Boesch, 1996; Goodall, 1986). Some have mimed inventively; examples are making hitting actions toward nuts they want cracked, blowing between thumb and forefinger to represent a balloon, and making twisting motions at containers they want opened (Miles et al., 1996; Russon, 2002; Savage-Rumbaugh et al., 1986).

One approach to assessing great ape intelligence is measuring their performances against children's on the same cognitive task. Chimpanzees can use scale models, for instance, which children first master in their third year (Kuhlmeier, Boysen & Mukobi, 1999). Chimpanzees and orangutans have solved *reverse contingency* tasks, which allow a subject to choose one of two sets of items (e.g., different amounts of candies) but then give the subject the set *not* chosen (Boysen et al., 1996; Shumaker et al., 2001). Chimpanzees who understood number symbols solved this task (chose the smaller amount to receive the larger) when amounts were shown by symbols, but failed with real foods. Children first solve this task between three and three and a half years of age and three year olds show limitations like the chimpanzees' (Carlson, Davis, & Leach, 2005). Thus some great apes show certain symbolic logical abilities comparable to those of three and a half year old children. To date, great apes have not shown evidence of the symbol systems that Deacon (1997) proposes to distinguish human intelligence.

Summary and Implications of Great Ape Research for Human Intelligence

There is now fairly strong agreement that great apes share a grade of intelligence of intermediate complexity that goes beyond that of other nonhuman primates and includes abilities previously thought uniquely human (Byrne, 1995; Gómez, 2004; Langer, 1996; Matsuzawa, 2001b; Parker & McKinney, 1999, Russon, 2004). A minority of primatologists view great ape intelligence as not significantly different from that of other nonhuman primates, on the one hand, or as more powerful but not reaching the currently defined human boundary, on the other (e.g., Povinelli, 2000; Suddendorf & Whiten, 2002; Tomasello & Call, 1997). Disagreement is due partly to emphasizing weak performances, interpreting monkey evidence too generously, neglecting great apes' most complex achievements, or incorrectly discounting them as artificially boosted by human enculturation. All-in-all, however, the evidence remains consistent with Premack's (1988) rule of thumb: Under normal

circumstances great apes can reach levels of intelligence of 3.5-year-old children, but not beyond.

In short, within the primates, many of the intellectual enhancements once considered uniquely hominin adaptations probably originated in the older and broader great ape lineage. Paleological evidence is consistent with a great ape grade of intelligence evolving with midMiocene hominids, as part and parcel of a biological package that includes larger brains, larger bodies, longer lives, and the mix of socioecological pressures the hominids faced and created (Russon & Begun, 2004). If so, these intellectual enhancements evolved as hominid adaptations to increasingly difficult life in moist tropical forests – not hominin adaptations to savanna life.

The Intelligence of Early Humans

This section examines the archaeological evidence for the earliest indications of human intelligence and anthropological evidence for concurrent changes in the size and shape of the cranial cavity. It discusses the implications for the evolution of human intelligence.

Homo Habilis

Ancestral humans started diverging from ancestral great apes approximately six million years ago. The first Homo lineage, *Homo habilis*, appeared approximately 2.4 million years ago in the Lower Paleolithic and persisted until 1.5 mya. The earliest known human inventions, referred to as *Oldowan* artifacts (after Olduvai Gorge, Tanzania, where they were first found), are widely attributed to *Homo habilis* (Semaw *et al.*, 1997), although it is possible that they were also used by late australopithecenes (de Baune, 2004). They were simple, mostly single faced stone tools, pointed at one end (Leakey, 1971). These tools were most likely used to split fruits and nuts (de Baune, 2004), although some of the more recently constructed ones have sharp edges, and are found with cut-marked bones, suggesting that they were used to sharpen wood implements and butcher small game (Leakey, 1971; Bunn & Kroll, 1986).

Although these carefully planed and deliberately fashioned early tools are seen as marking a momentous breakthrough for our lineage, they were nevertheless simple and unspecialized; by our standards they were not indicative of a very flexible or creative kind of intelligence. The same tools were put to many uses instead of adapting them to precisely meet the task at hand. Mithen (1996) refers to minds at this time as possessing *generalized intelligence*, reflecting his belief that associative-level domain-general learning mechanisms, such as operant and Pavlovian conditioning, predominated. The minds of these early hominins have been referred to as *pre-representational*, because available artifacts show no

indication that the hominins were capable of forming representations that deviated from their concrete sensory perceptions; their experience is considered to have been *episodic*, or tied to the present moment (Donald, 1993). Donald characterized their intelligence as governed by procedural memory. They could store perceptions of events and recall them in the presence of a reminder or cue, but they had little voluntary access to episodic memories without environmental cues. They were therefore unable to voluntarily shape, modify, or practice skills and actions, and they were unable to invent or refine complex gestures or means of communicating.

The Massive Modularity Hypothesis

Evolutionary psychologists claim that the intelligence of *Homo* arose due to *massive modularity* (Buss, 1999, 2004; Buss *et al.*, 1994; Cosmides & Tooby, 2002; Dunbar *et al.*, 1994; Rozin, 1976; for an extensive critique see Buller, 2005 and Byrne, 2000). Cosmides and Tooby (1992) proposed that human intelligence evolved in the form of hundreds or thousands of functionally encapsulated (that is, not accessible to each other) cognitive modules. Each module was specialized to accomplish a specific task or solve a specific problem encountered by ancestral humans in their *environment of evolutionary adaptedness*, taken to be hunter-gatherer life in the Pleistocene. Modules for language, theory of mind, spatial relations, and tool use are among the modules proposed. These modules are supposedly content rich, pre-fitted with knowledge relevant to hunter-gatherer problems. It is also claimed that these modules exist today in more or less the same form as they existed in the Pleistocene, because too little time has passed for them to have undergone significant modification.

What is the current status of these ideas? Although the mind exhibits an intermediate degree of functional and anatomical modularity, neuroscience has not revealed vast numbers of hardwired, encapsulated, task-specific modules; indeed, the brain has been shown to be more highly subject to environmental influence than we thought (Wexler, 2006). Nevertheless, evolutionary psychology has made a valuable contribution by heightening awareness that the human mind is not an optimally designed machine; its structure and function reflect the pressures it was subjected to in over its long evolutionary history.

Homo erectus

Approximately 1.9 million years ago, *Homo ergaster* and *Homo erectus* appeared, followed by archaic *Homo sapiens* and *Homo neanderthalensis*. The size of the *Homo erectus* brain was approximately 1,000 cc, about 25% larger than that of *Homo habilis*, at least twice as

large as those of living great apes, and 75% the cranial capacity of modern humans (Aiello, 1996; Ruff *et al.*, 1997; Lewin, 1999). *Homo erectus* exhibited many indications of enhanced ability to adapt to the environment to meet the demands of survival, including sophisticated, task-specific stone hand axes, complex stable seasonal home bases, and long-distance hunting strategies involving large game. By 1.6 mya, *Homo erectus* had dispersed as far as Southeast Asia, indicating the ability to adjust its lifestyle to different climates and habitats (Anton & Swisher, 2004; Cachel & Harris, 1995; Swisher, Curtis, Jacob, Getty, & Widiasmoro, 1994; Walker & Leakey, 1993). By 1.4 mya in Africa, West Asia, and Europe, *Homo erectus* had produced the Acheulean handaxe (Asfaw *et al.*, 1992), a do-it-all tool that may have functioned as a social status symbol (Kohn & Mithen, 1999). The most notable characteristic of these tools is their biface (two-sided) symmetry. They probably required several stages of production, bifacial knapping, and considerable skill and spatial ability to achieve their final form.

Though anatomical evidence indicates the presence of Broca's area in the brain, suggesting that the capacity for language was present by this time (Wynn, 1998), verbal communication is thought to have been limited to (at best) pre-syntactical proto-language involving primarily short, nongrammatical utterances of one or two words (Dunbar, 1996). Mental processes during this time period probably strayed little from concrete sensory experience. The capacity for abstract thought, and for thinking about what one is thinking about (that is, metacognition), had not yet appeared.

Social Explanations

There are multiple versions of the hypothesis that the origins of human intellect and onset of the archaeological record reflect a transition in cognitive or social abilities. *Homo erectus* were indeed probably the earliest humans to live in hunter-gatherer societies. One suggestion has been that they owe their achievements to onset of theory of mind (Mithen, 1998). However, as we have seen, there is evidence that other species possess theory of mind (Heyes, 1998), yet do not compare to modern humans in intelligence.

Self-triggered Recall and Rehearsal Loop

Donald (1991) proposed that with the enlarged cranial capacity of *Homo erectus*, the human mind underwent the first of three transitions by which it – and the cultural matrix in which it is profoundly embedded – evolved from the ancestral, pre-hominin condition. Each transition entailed a new way of encoding representations in memory and storing them in collective memory so that they can later be drawn upon and shared with others.

This first transition is characterized by a shift from an *episodic* to a *mimetic mode* of cognitive functioning, made possible by onset of the capacity for voluntary retrieval of stored memories, independent of environmental cues. Donald refers to this as a “self-triggered recall and rehearsal loop.” Self-triggered recall enabled hominins to access memories voluntarily and thereby act out¹ events that occurred in the past or that might occur in the future. Thus not only could the mimetic mind temporarily escape the here and now, but by miming or gesture, it could communicate similar escapes in other minds. The capacity to mime thus ushered forth what is referred to as a *mimetic* form of cognition and brought about a transition to the mimetic stage of human culture. The self-triggered recall and rehearsal loop also enabled hominins to engage in a stream of thought. One thought or idea evokes another, revised version of it, which evokes yet another, and so forth recursively. In this way, attention is directed away from the external world toward one’s internal model of it. Finally, self-triggered recall allowed actors to take control over their own output, including voluntary rehearsal and refinement, and mimetic skills such as pantomime, reenactive play, self-reminding, imitative learning, and proto-teaching. In effect, it allows systematic evaluation and improvement of motor acts and adapting them to new situations, resulting in more refined skills and artifacts, and the capacity to use one’s body as a communication device to act out events.

Donald’s scenario becomes even more plausible in light of the structure and dynamics of associative memory (Gabora, 1998, 2003, 2007; Gabora & Aerts, 2009). Neurons are sensitive to *microfeatures* – primitive stimulus attributes such as a sound of a particular pitch, or a line of a particular orientation. Episodes etched in memory are *distributed* across a bundle or cell assembly of these neurons, and likewise, each neuron participates in the encoding of many episodes. Finally, memory is *content-addressable*, such that similar stimuli activate and get encoded in overlapping distributions of neurons. With larger brains, episodes are encoded in more detail, allowing for a transition from more coarse-grained to more fine-grained memory. Fine-grained memory means more microfeatures of episodes tend to be encoded, so there are more ways for distributions to overlap. Greater overlap meant more routes by which one memory can evoke another, making possible the onset of self-triggered recall and rehearsal, and paving the way for a more integrated internal model of the world, or worldview.

Over a Million Years of Stasis

The handaxe persisted as the almost exclusive tool preserved in the archaeological record for over a million years, spreading by 500,000 years ago into Europe, where it was used until about 200,000 years ago. During this period, there was almost no change in tool design and little other evidence of new forms of intelligent behavior, with the exception of the first solid evidence for controlled use of fire, approximately 800,000 years ago (Goren-Inbar et al., 2004). There is, however, some evidence (such as charred animal bones at *Homo ergaster* sites) that fire may have been used substantially earlier.

A Second Increase in Brain Size

Between 600,000 and 150,000 years ago there was a second spurt in brain enlargement (Aiello, 1996; Ruff *et al.*, 1997), which marks the appearance of anatomically modern humans. It would make our story simple if the increase in brain size coincided with the burst of creativity in the Middle/Upper Paleolithic (Bickerton, 1990; Mithen, 1998), to be discussed shortly. But although *anatomically* modern humans had arrived, *behavioral* modernity had not. Leakey (1984) writes of anatomically modern human populations in the Middle East with little in the way of evidence for the kind of intelligence of modern humans and concludes, “The link between anatomy and behavior therefore seems to break” (p. 95). An exception to the overall lack of evidence for intellectual progress at this time is the advancement of the Levallois flake, which came into prominence approximately 250,000 years ago in the Neanderthal line. This suggests that cognitive processes were primarily first-order—tied to concrete sensory experience—rather than second-order—derivative, or abstract.

Perhaps this second spurt in encephalization exerted an impact on expressions of intelligence that left little trace in the archaeological record, such as ways of coping with increasing social complexity, or manipulating competitors (Baron-Cohen, 1995; Byrne & Whiten, 1988; Dunbar, 1996; Humphrey, 1976; Whiten, 1991; Whiten & Byrne, 1997; Wilson et al., 1996). Another possible reason for the apparent rift between anatomical and behavioral modernity is that while genetic changes necessary for cognitive modernity arose at this time, the fine-tuning of the nervous system to fully capitalize on these genetic changes took longer, or the necessary environmental conditions were not yet in place (Gabora, 2003). It is worth noting that other periods of revolutionary innovation, such as the Holocene transition to agriculture and the modern Industrial Revolution, occurred long after the biological changes that made them cognitively possible.

The Spectacular Intelligence of Modern Humans

The European archaeological record indicates that an unparalleled transition occurred between 60,000 and 30,000 years ago at the onset of the Upper Paleolithic (Bar-Yosef, 1994; Klein, 1989a; Mellars, 1973, 1989a, 1989b; Soffer, 1994; Stringer & Gamble, 1993). Considering it “evidence of the modern human mind at work,” Richard Leakey (1984:93-94) writes: “unlike previous eras, when stasis dominated, ... [with] change being measured in millennia rather than hundreds of millennia.” Similarly, Mithen (1996) refers to the Upper Paleolithic as the ‘big bang’ of human culture, exhibiting more innovation than in the previous six million years of human evolution.

At this time we see the more or less simultaneous appearance of traits considered diagnostic of behavioral modernity. They include the beginning of a more organized, strategic, season-specific style of hunting involving specific animals at specific sites, elaborate burial sites indicative of ritual and religion, evidence of dance, magic, and totemism, the colonization of Australia, and replacement of Levallois tool technology by blade cores in the Near East. In Europe, complex hearths and many forms of art appeared, including naturalistic cave paintings of animals, decorated tools and pottery, bone and antler tools with engraved designs, ivory statues of animals and sea shells, and personal decoration such as beads, pendants, and perforated animal teeth, many of which may have indicated social status (White 1989a,b). White (1982:176) also wrote of a “total restructuring of social relations”. What is perhaps most impressive about this period is not the novelty of any particular artifact but that the overall pattern of change is cumulative; more recent artifacts resemble older ones but have modifications that enhance their appearance or functionality. This cumulative change is referred to as the *ratchet effect* (Tomasello, Kruger & Ratner, 1993), and some suggest it is uniquely human (Donald, 1998).

Whether this period was a genuine revolution culminating in behavioral modernity is hotly debated because claims to this effect are based on the European Palaeolithic record, and largely exclude the African record (McBrearty & Brooks, 2000); Henshilwood & Marean, 2003). Indeed, most of the artifacts associated with a rapid transition to behavioral modernity at 40,000– 50,000 years ago in Europe are found in the African Middle Stone Age tens of thousands of years earlier. These artifacts include blades and microliths, bone tools, specialized hunting, long distance trade, art and decoration (McBrearty & Brooks, 2000), the Berekhat Ram figurine from Israel (d’Errico & Nowell, 2000), and an anthropomorphic figurine of quartzite from the Middle Ascheulian (ca. 400 ka) site of Tan-tan in Morocco

(Bednark, 2003). Moreover, gradualist models of the evolution of cognitive modernity well before the Upper Palaeolithic find some support in archaeological data (Bahn, 1991; Harrold, 1992; Henshilwood & Marean, 2003; White, 1993; White *et al.*, 2003). If modern human behaviors were indeed gradually assembled as early as 250,000–300,000 years ago, as McBrearty and Brooks (2000) argue, the transition falls more closely into alignment with the most recent spurt in human brain enlargement. However, the traditional and currently dominant view is that modern behavior appeared in anatomically modern humans in Africa between 50,000 and 40,000 years ago due to biologically evolved cognitive advantages, and that anatomically modern humans spread replacing existing species, including the Neanderthals in Europe e.g., Ambrose, 1998; Gamble, 1994; Klein, 2003; Stringer & Gamble, 1993). Thus, from this point onward, there was only one hominin species: the modern *Homo sapiens*.

Despite lack of overall increase in cranial capacity, the prefrontal cortex, and particularly the orbitofrontal region, increased disproportionately in size (Deacon, 1997; Dunbar, 1993; Jerison, 1973; Krasnegor, Lyon, & Goldman-Rakic, 1997; Rumbaugh, 1997) and it was likely a time of major neural reorganization (Henshilwood, d'Errico, Vanhaeren, van Niekerk, & Jacobs, 2000; Klein, 1999). These brain changes may have given rise to metacognition, or what Feist (2006) refers to as “meta-representational thought,” that is, the ability to reflect on representations and think about thinking. Whether or not it is considered a “revolution,” it is accepted that the Middle/Upper Paleolithic was a period of unprecedented intellectual activity. How and why did it occur? Let us now review the most popular hypotheses for how and why behavioral modernity and its underlying intellectual capacities arose.

Syntactic Language and Symbolic Reasoning

It has been suggested that at this time humans underwent a transition from a predominantly gestural to a vocal form of communication (Corballis, 2002). Although the ambiguity of the archaeological evidence means we may never know exactly when language began (Bednark, 1992:30; Davidson & Noble, 1989), most scholars agree that earlier *Homo* and even Neanderthals may have been capable of primitive proto-language and the grammatical and syntactic aspects emerged at the start of the Upper Palaeolithic (Aiello & Dunbar, 1993; Bickerton, 1990, 1996; Dunbar, 1993, 1996; Tomasello, 1999). Carstairs-McCarthy (1999) presented a modified version of this proposal, suggesting that although some form of syntax was present in the earliest languages, most of the later elaboration, including recursive

embedding of syntactic structure, emerged in the Upper Paleolithic. Syntax enabled the capacity to state more precisely how elements are related and to embed them in other elements. Thus it enabled language to become general-purpose and applied in a variety of situations.

Deacon (1997) stresses that the onset of complex language reflects onset of the capacity to internally representing complex, abstract, internally coherent systems of meaning using symbols—items, such as words, that arbitrarily stand for other items, such as things in the world. The advent of language made possible what Donald (1991) refers to as the *mythic* or story-telling stage of human culture. It enhanced not just the ability to communicate with others, spread ideas from one individual to the next, and collaborate (thereby speeding up cultural innovation), but also the ability to think things through for oneself and manipulate ideas in a controlled, deliberate fashion (Reboul, 2007).

Cognitive Fluidity, Connected Modules, and Cross-Domain Thinking

Another proposal is that the exceptional abilities exhibited by *Homo* in the Middle/Upper Paleolithic were due to the onset of *cognitive fluidity* (Fauconnier and Turner, 2002). Cognitive fluidity involves the capacity to draw analogies, to combine concepts and adapt ideas to new contexts, and to map across different knowledge systems, potentially employing multiple ‘intelligences’ simultaneously (Gardner, 1983; Langer, 1996; Mithen, 1996). Cognitive fluidity would have facilitated the weaving of experiences into stories, parables, and broader conceptual frameworks, and thereby the integration of knowledge and experience (Gabora & Aerts, 2009).

A related proposal has been put forward by Mithen (1996). Drawing on the evolutionary psychologist’s notion of massive modularity, he suggests that the abilities of the modern human mind arose through the interconnecting of preexisting intellectual modules (that is, encapsulated or functionally isolated *specialized intelligences*, or cognitive domains) devoted to natural history, technology, social processes, and language. This interconnecting, he claims, is what enabled the onset of cognitive fluidity and allowed humans to map, explore, and transform conceptual spaces. Sperber (1994) proposed that the connecting of modules involved a special module, the “module of meta-representation,” which contains “concepts of concepts” and enabled cross-domain thinking, and particularly analogies and metaphors.

Contextual Focus: Shifting Between Explicit and Implicit Modes of Thought

These proposals for what kinds of cognitive change could have led to the Upper Paleolithic transition stress different aspects of cognitive modernity. Acknowledging a possible seed of truth in each, we begin to converge toward a common (if more complex) view. Concept combination is characteristic of *divergent thought*, which tends to be intuitive, diffuse, and associative. Divergent thought is on the opposite end of the spectrum from the *convergent thought* stressed by Deacon, which tends to be logical, controlled, effortful, and reflective and symbolic. Converging evidence suggests that the modern mind engages in both (Arieti, 1976; Ashby & Ell, 2002; Freud, 1949; Guilford, 1950; James, 1890/1950; Johnson-Laird, 1983; Kris, 1952; Neisser, 1963; Piaget, 1926; Rips, 2001; Sloman, 1996; Stanovich & West, 2000; Werner, 1948; Wundt, 1896). This is sometimes referred to as the dual-process theory of human cognition (Chaiken & Trope, 1999; Evans & Frankish, 2009) and it is consistent with some current theories of cognition (Finke, Ward, & Smith, 1992; Gabora, 2000, 2002, 2003, under revision; S.B. Kaufman, this volume. Divergent processes are hypothesized to facilitate insight and idea generation, while convergent processes predominate during the refinement, implementation, and testing of an idea.

It has been proposed that the Paleolithic transition reflects genetic changes involved in the fine-tuning of the biochemical mechanisms underlying the capacity to shift between these modes of thought, depending on the situation, by varying the specificity of the activated cognitive receptive field (Gabora, 2003, 2007; for similar ideas see Howard-Jones & Murray, 2003; Martindale, 1995). This capacity is referred to as *contextual focus*² because it requires the ability to focus or defocus attention in response to the context or situation one is in. Defocused attention, by diffusely activating a broad region of memory, is conducive to divergent thought; it enables obscure (but potentially relevant) aspects of the situation to come into play. Focused attention is conducive to convergent thought; memory activation is constrained enough to hone in and perform logical mental operations on the most clearly relevant aspects. Note that contextual focus enables dynamic “resizing” of the activated brain region in response to the situation (as opposed to rigid compartmentalization).

Once the capacity to shrink or expand the field of attention came about, thereby improving the capacity to tailor one’s mode of thought to the demands of the current situation, tasks requiring convergent thought (e.g., mathematical derivation), divergent thought (e.g., poetry), or both (e.g., technological invention) could be carried out more

effectively. When the individual is fixated or stuck, and progress is not forthcoming, defocusing attention enables the individual to enter a more divergent mode of thought, and peripherally related elements of the situation begin to enter working memory until a potential solution is glimpsed. At this point attention becomes more focused, and thought becomes more convergent, as befits the fine-tuning of the idea and manifestation of it in the world.

Thus, the onset of contextual focus would have enabled hominins to adapt ideas to new contexts or combine them in new ways through divergent thought and fine-tune these unusual new combinations through convergent thought. In this way, the fruits of one mode of thought provide the ingredients for the other, culminating in a more fine-grained internal model of the world. A related proposal is that this period marks the onset of the capacity to move between explicit and implicit modes of thought (Feist, 2007). Explicit thought involves the executive functions concerned with control of cognitive processes such as planning and decision making, while implicit thought encompasses the ability to automatically and nonconsciously detect complex regularities, contingencies, and covariances in our environment (Kaufman, DeYoung, Gray, Jiménez, Brown, & Mackintosh, N., under revision). A contributing factor to the emergence of the ability to shift between them may have been the expansion of the prefrontal cortex. This expansion probably enhanced the executive functions as well as the capacity to maintain and manipulate information in an active state in working memory. Indeed, individual differences in working memory capacity are strongly related to fluid intelligence in modern humans (Conway, Jarrold, Kane, & Miyake, 2007; Engle, Tuholski, Laughlin, & Conway, 1999; Kane, Hambrick, & Conway, 2005; Kaufman, DeYoung, Gray, Brown, & Mackintosh, 2009).

Evolution of intelligence

Intelligence, as a concept and a field of study, has been a cornerstone of human progress. It spans cognitive abilities, emotional insight, problem-solving skills, and decision-making. Recent advancements in artificial intelligence (AI), neuroscience, and cognitive science have redefined the boundaries of what intelligence is and how it can be enhanced. This chapter explores these breakthroughs, examining their implications and challenges.

Artificial Intelligence: Redefining Cognitive Capabilities

Artificial Intelligence (AI) is no longer confined to science fiction. Over the past decade, it has transitioned from basic algorithms to sophisticated systems capable of mimicking human thought processes. Advances in deep learning, natural language processing, and

reinforcement learning have enabled AI to achieve milestones like generating human-like text, recognizing emotions, and even creating art.

For instance, the advent of generative AI models like GPT-4 and beyond demonstrates remarkable capabilities in language understanding and generation. These systems are not just tools but collaborative partners in tasks ranging from medical diagnoses to creative writing. AI's integration into daily life, from virtual assistants to autonomous vehicles, marks a shift in how intelligence is perceived—blurring the lines between human and machine.

Neuroscience: Mapping the Brain

Parallel to AI, neuroscience has made strides in understanding the biological foundations of intelligence. The Human Connectome Project and advancements in neuroimaging have provided detailed maps of neural connections in the brain. These maps reveal how different regions collaborate to produce complex behaviors and thought patterns.

Recent studies have identified specific neural correlates of intelligence, such as the role of the prefrontal cortex in abstract thinking and decision-making. Techniques like transcranial magnetic stimulation (TMS) and deep brain stimulation (DBS) are being explored to enhance cognitive functions and treat neurological disorders.

Moreover, the growing field of brain-computer interfaces (BCIs) aims to bridge the gap between biology and technology. Companies like Neuralink are developing implants that allow direct communication between the brain and external devices, offering possibilities such as restoring mobility in paralyzed individuals and enhancing memory retention.

Emotional Intelligence: A Growing Focus

In contrast to traditional measures of intelligence, such as IQ, emotional intelligence (EQ) emphasizes understanding and managing emotions. Recent developments have highlighted the role of EQ in leadership, team dynamics, and mental well-being.

Technology has started to play a role in this domain as well. AI systems are being trained to recognize and respond to human emotions through facial expressions, voice tones, and text cues. These systems have applications in customer service, therapy, and even conflict resolution, where understanding emotions is critical.

Ethical and Social Implications

While the advancements in intelligence are exciting, they raise important ethical questions. The rapid development of AI has sparked debates about privacy, bias, and accountability. Autonomous systems must be programmed with ethical frameworks to ensure they align with human values.

In neuroscience, concerns about "neuroenhancement"—using technology to artificially boost intelligence—have prompted discussions about equity and consent. If only a select few can afford such enhancements, societal divisions could deepen.

Similarly, as AI begins to outperform humans in certain tasks, questions arise about the future of work and the role of human creativity. Balancing progress with ethical considerations is paramount to ensure these technologies benefit humanity as a whole.

The Road Ahead

The fusion of AI, neuroscience, and emotional intelligence marks the dawn of a new era in understanding and enhancing intelligence. Future research may lead to systems that seamlessly integrate human and machine intelligence, creating hybrid models of cognition.

The journey to unravel the mysteries of intelligence is far from over. As we navigate this uncharted territory, we must remain vigilant, ensuring that innovation serves humanity while preserving the essence of what makes us intelligent beings—our creativity, empathy, and ethical judgment.

Conclusion

The concept of intelligence is complex and multifaceted, and our understanding of it continues to evolve. Moving beyond the traditional focus on general intelligence, contemporary research emphasizes the importance of multiple intelligences, including emotional, social, and creative intelligence. Technological advancements, such as AI and neuroscience, are further expanding our understanding of the human mind and the nature of intelligence. By recognizing the diversity of human intelligence and fostering the development of a wide range of skills and abilities, we can create a more inclusive and equitable society that values and celebrates individual strengths and talents. This chapter not only examines the state-of-the-art developments but also emphasizes the need for responsible innovation, paving the way for a future where intelligence—both natural and artificial—enhances our collective potential.

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Publications

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Innovative Approaches in Additive Manufacturing: Enhancing Efficiency and Sustainability in Mechanical Production

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Abstract

The evolution of manufacturing technologies has significantly transformed industrial production, with Additive Manufacturing (AM) emerging as a pivotal innovation in enhancing efficiency and sustainability. Unlike traditional subtractive manufacturing methods, AM constructs components layer by layer, reducing material waste and energy consumption while enabling intricate and lightweight designs. This paper explores innovative approaches in AM, including topology optimization, multi-material printing, AI-driven process enhancements, and automation. These advancements improve material efficiency, production speed, and defect detection, positioning AM as a critical tool for sustainable mechanical production. Additionally, AM's role in minimizing environmental impact through localized production, energy-efficient processes, and circular economy principles is discussed. Case studies from aerospace, automotive, and biomedical industries highlight successful AM applications that enhance both economic viability and ecological responsibility. By integrating AI, robotics, and sustainable materials, Additive Manufacturing is poised to revolutionize modern manufacturing, addressing global challenges in resource efficiency and environmental sustainability.

Keywords: Additive Manufacturing (AM), Sustainability in Manufacturing, Topology Optimization, Multi-Material Printing, AI-Driven Defect Detection, Circular Economy

1.0. Introduction

The rapid evolution of manufacturing technologies has significantly reshaped modern industrial processes, with additive manufacturing (AM) emerging as a transformative innovation. AM, commonly known as 3D printing, is a process that fabricates components

layer by layer, enabling precise material deposition and complex geometric designs (Gibson, Rosen, & Stucker, 2021). Unlike traditional subtractive manufacturing techniques, which involve material removal through cutting, milling, or drilling, additive manufacturing (AM) optimizes resource utilization by minimizing waste and energy consumption (Baumers, Dickens, Tuck, & Hague, 2016). This shift has positioned AM as a key technology in promoting efficiency and sustainability in mechanical production.

One of the most significant advantages of AM lies in its ability to enhance design flexibility and material efficiency. Engineers can now develop intricate structures that were previously unachievable using conventional manufacturing methods (Leary et al., 2019). For instance, topology optimization—an advanced computational design approach—enables the production of lightweight yet high-strength components, which is particularly beneficial in industries such as aerospace and automotive manufacturing (Bai, Xie, Zhang, & Liu, 2021). Additionally, multi-material printing has revolutionized product development by allowing the integration of diverse materials within a single component, enhancing mechanical properties and functionality (Bartolo & Kruth, 2017).

The sustainability impact of AM is another crucial factor driving its adoption. The traditional manufacturing sector is one of the largest contributors to global carbon emissions due to high energy consumption and material waste (Kellens et al., 2017). AM addresses these challenges by enabling localized production, reducing transportation-related emissions, and facilitating the use of recycled materials (Ford & Despeisse, 2016). Moreover, AM supports circular economy principles by allowing end-of-life recycling of components, further promoting environmental sustainability (Tlegenov, Lu, & Xu, 2021).

Despite these advantages, AM is not without challenges. Issues such as material limitations, production speed, and quality control continue to hinder large-scale adoption (Peng, Chen, & Zhou, 2018). However, recent advancements in artificial intelligence (AI) and machine learning have significantly improved AM processes by optimizing print parameters, detecting defects in real-time, and enhancing overall production efficiency (Zhang, Wang, & Liu, 2021). These technological integrations have propelled AM from a prototyping tool to a viable solution for full-scale manufacturing.

This paper explores innovative approaches in AM that enhance efficiency and sustainability in mechanical production. The study focuses on topology optimization, multi-material printing, bio-inspired designs, and AI-driven process enhancements. Furthermore, the environmental benefits of AM and its role in shaping a more sustainable manufacturing

ecosystem are discussed. By examining recent developments and industrial applications, this research aims to highlight AM's potential in revolutionizing modern manufacturing while addressing sustainability challenges.

2.0. Enhancing Efficiency in Additive Manufacturing

Additive manufacturing (AM) has revolutionized mechanical production by enabling highly efficient, precise, and sustainable fabrication processes. However, optimizing AM efficiency remains a critical focus for researchers and industry professionals. Efficiency in AM can be enhanced through various innovative approaches, including topology optimization, multi-material printing, process parameter optimization, automation and robotics, and artificial intelligence (AI)-driven defect detection and quality control. These advancements contribute to increased production speed, improved material utilization, reduced energy consumption, and superior mechanical properties of printed components.

2.1. Topology Optimization for Material Efficiency in Additive Manufacturing

Topology optimization is a computational design method that strategically distributes material within a given design space to achieve maximum strength with minimal weight (Sigmund & Maute, 2013). This method significantly enhances efficiency in Additive Manufacturing by minimizing material usage without compromising structural integrity. This results in lightweight, high-strength components with minimal material waste. Table 1 elaborates on the key aspects of topology optimization, its advantages, applications, and case studies.

For example, Airbus has implemented topology optimization in manufacturing aircraft brackets, leading to weight reductions of up to 55%, which in turn reduces fuel consumption and operational costs (Leary et al., 2019). Similarly, the automotive industry has leveraged topology optimization to manufacture lightweight but durable vehicle components, enhancing fuel efficiency and sustainability (Schmid et al., 2019).

By integrating topology optimization with AM, engineers can design intricate lattice structures and bio-inspired geometries that would be impossible or highly inefficient using conventional manufacturing techniques (Bai, Xie, Zhang, & Liu, 2021).

Table 1: Topology Optimization for Material Efficiency in Additive Manufacturing

Aspect	Details	References
Definition	Topology optimization is a mathematical approach	Sigmund &

	that optimizes material distribution in a given design space to achieve maximum performance with minimal weight.	Maute (2013)
Key Principle	Uses iterative finite element analysis (FEA) to remove non-essential material while maintaining structural integrity.	Leary et al. (2019)
Material Efficiency	Minimizes raw material consumption , reducing waste by up to 50% compared to traditional methods .	Schmid et al. (2019)
Lightweight Design	Generates high-strength, lightweight structures , reducing overall weight by 30-55%.	Bai et al. (2021)
Enhanced Performance	Produces components with optimized stiffness, strength, and load distribution , leading to superior mechanical performance.	Gibson, Rosen, & Stucker (2021)
Applications in Aerospace	Boeing and Airbus use topology optimization for aircraft brackets and structural components , reducing weight and fuel consumption.	DebRoy et al. (2018)
Applications in Automotive	Automakers use optimized chassis, suspension arms, and engine mounts for enhanced fuel efficiency.	Schmid et al. (2019)
Applications in Biomedical	Used to create patient-specific implants with porous structures that enhance bone integration and reduce weight.	Choi, Kim, & Kim (2021)
Case Study: Boeing	Boeing's topology-optimized aircraft brackets reduced part weight by 55% , improving fuel efficiency.	Leary et al. (2019)
Case Study: Airbus	Airbus used 3D-printed, topology-optimized components in the A350 XWB, cutting material waste by 45% .	Schmid et al. (2019)
Software Tools	Common software includes ANSYS, Altair OptiStruct, Siemens NX, and Autodesk Fusion 360 for design optimization.	Gibson, Rosen, & Stucker (2021)
Challenges	Computationally intensive, requires high processing	Sigmund &

	power ; final designs may need post-processing modifications .	Maute (2013)
Future Trends	AI-integrated topology optimization for real-time adaptive manufacturing and improved design flexibility.	Zhang, Wang, & Liu (2021)

Insights from table 1:

1. Material Reduction & Cost Savings:
 - Traditional machining wastes up to 70-90% of raw material due to subtractive processes (DebRoy et al., 2018).
 - With topology optimization, material usage can be reduced by 50%, leading to significant cost savings and sustainability improvements (Schmid et al., 2019).
2. Improved Performance & Durability:
 - Optimized structures distribute stress more effectively, preventing premature failure.
 - Studies show that optimized aerospace brackets can handle higher loads with 30% less weight (Bai et al., 2021).
3. Wider Adoption Across Industries:
 - Aerospace & Automotive: Lightweight but durable parts improve fuel efficiency.
 - Biomedical: Custom implants with optimized porous structures enhance biocompatibility and bone integration (Choi et al., 2021).
 - Industrial Machinery: Optimized machine components reduce energy consumption and improve operational life.
4. Future Prospects with AI & Machine Learning:
 - AI-powered topology optimization can automatically refine designs based on performance data, reducing computational time (Zhang et al., 2021).
 - Real-time optimization will enable adaptive manufacturing, adjusting designs during the printing process for optimal material efficiency.

2.2. Multi-material Printing for Functional Efficiency in Additive Manufacturing

Traditional AM primarily relied on single-material fabrication, limiting functional performance. However, advancements in multi-material printing allow manufacturers to

combine different materials within a single print, enhancing mechanical properties, thermal resistance, and electrical conductivity (Bartolo & Kruth, 2017).

For instance, in the biomedical sector, multi-material AM is used to fabricate prosthetics with rigid and flexible sections, ensuring comfort and durability (Choi, Kim, & Kim, 2021). In mechanical production, hybrid metal-polymer printing has resulted in improved wear resistance and mechanical strength, extending the lifespan of components (Gibson, Rosen, & Stucker, 2021).

This approach not only enhances efficiency by eliminating post-processing **assembly steps** but also enables the creation of smart components with embedded sensors for real-time monitoring of mechanical performance (Wang et al., 2020).

Multi-material printing is an advanced additive manufacturing (AM) technique that allows for the fabrication of components using two or more materials in a single build process. This approach enhances functional efficiency, enabling the development of stronger, lighter, and more complex structures with tailored properties. Table 2 outlines the major aspects of multi-material printing, its advantages, applications, and case studies.

Table 2: Multi-material Printing for Functional Efficiency

Aspect	Details	References
Definition	Fabrication of single components with multiple materials to achieve enhanced mechanical, thermal, and electrical properties.	Espinosa et al. (2021)
Key Principle	Utilizes simultaneous or sequential deposition of different materials, often in a voxel-based approach.	Choi & Kim (2019)
Material Combinations	Common combinations include metal-metal, metal-polymer, polymer-ceramic, and bio-compatible materials.	Tumbleston et al. (2017)
Functional Efficiency	Enhances mechanical strength, thermal resistance, flexibility, electrical conductivity, and biocompatibility.	Bandyopadhyay et al. (2019)
Lightweight Structures	Multimaterial topology optimization allows for targeted reinforcement in high-stress areas , reducing weight.	Zhang et al. (2021)

Enhanced Durability	Different materials can be combined to create wear-resistant surfaces while maintaining structural flexibility.	Gu et al. (2018)
Applications in Aerospace	Used in aircraft turbine blades, heat shields, and structural components , improving performance under extreme conditions.	Yang et al. (2020)
Applications in Automotive	Enables graded structures for impact absorption and lightweight designs in engine components and chassis parts .	Liu & Wang (2022)
Applications in Electronics	Printed flexible circuits, sensors, and conductive pathways , allowing for smart electronics integration .	Zardini et al. (2021)
Applications in Biomedical	Used for patient-specific implants, prosthetics, and bone scaffolds with varied stiffness and biocompatibility.	Jakus et al. (2017)
Case Study: NASA	NASA used multimaterial 3D printing to develop rocket nozzles with graded thermal properties , increasing efficiency.	Espinosa et al. (2021)
Case Study: MIT	Researchers at MIT developed multimaterial robotic skins , integrating soft and rigid materials for adaptable motion.	Bandyopadhyay et al. (2019)
Challenges	Requires precise control of material deposition, advanced software, and post-processing techniques .	Choi & Kim (2019)
Future Trends	AI-driven real-time material adaptation for optimized performance in aerospace and medical applications.	Zhang et al. (2021)

Insights from table 2:

- Enhanced Functional Efficiency Through Material Integration:**

- Unlike traditional manufacturing, multimaterial printing enables localized material properties, reducing the need for multiple assembly steps (Espinosa et al., 2021).
- Example: A multimaterial jet engine blade can have a heat-resistant ceramic core and a tough metal exterior, balancing thermal and mechanical performance (Yang et al., 2020).

2. Weight Reduction & Structural Optimization:

- Combining lightweight polymers with high-strength metals reduces weight while maintaining mechanical integrity (Liu & Wang, 2022).
- Example: In automotive chassis, impact-resistant zones can be printed with graded stiffness to absorb shocks efficiently (Gu et al., 2018).

3. Biomedical & Electronics Advancements:

- Medical implants with graded porosity improve bone integration and reduce stress shielding (Jakus et al., 2017).
- Flexible electronics with conductive pathways allow for wearable sensors and smart textiles (Zardini et al., 2021).

4. Future Trends & AI Integration:

- AI-based optimization will allow real-time adaptation of materials during printing, improving energy efficiency and material utilization (Zhang et al., 2021).
- Researchers are exploring self-healing multimaterial structures, useful for autonomous repair in aerospace and defense applications (Bandyopadhyay et al., 2019).

2.3. Process Parameter Optimization for Faster Production

Optimizing process parameters, such as laser power, scanning speed, layer thickness, and cooling rates, plays a crucial role in improving AM efficiency. Poor parameter selection can lead to defects such as porosity, warping, or excessive material usage (DebRoy et al., 2018).

Researchers have developed AI-driven algorithms that analyze historical print data and dynamically adjust process parameters in real time. For example, AI-based models have successfully reduced print failure rates by 30% and improved production speeds by 25% in metal AM applications (Zhang, Wang, & Liu, 2021).

Moreover, advanced thermal management techniques, such as selective laser melting (SLM) with optimized heat distribution, have been shown to significantly reduce cooling times and

residual stresses, improving part quality and efficiency (Peng, Chen, & Zhou, 2018). Table 3 illustrates the relationship between these parameters and their impact on production speed and efficiency.

Table 3: Process Parameter Optimization for Faster Production

Process Parameter	Optimization Strategy	Impact on Faster Production	References
Laser Power (P)	Higher power leads to faster melting but requires control to avoid overheating.	Reduces build time but may affect precision.	DebRoy et al. (2018)
Scanning Speed (V)	Increased speed enhances productivity but can cause defects if too high.	Balance is needed to ensure both speed and quality.	Yang et al. (2020)
Layer Thickness (h)	Thicker layers speed up production but may compromise surface finish.	Reducing layer count significantly lowers print time.	Gibson et al. (2021)
Hatch Spacing (S)	Wider spacing reduces overlaps and speeds up material deposition.	Optimized hatch spacing minimizes defects and print time.	Bai et al. (2021)
Cooling Rate	Controlled cooling minimizes warping and enhances layer adhesion.	Faster cooling enables higher throughput without distortion.	Zhang et al. (2021)
Powder Feed Rate	Adjusting feed rate ensures consistent material flow and energy absorption.	Prevents interruptions and enhances continuous printing.	Gu et al. (2018)
Gas Flow Rate	Ensures uniform powder distribution and prevents oxidation.	Improves consistency, reducing post-processing needs.	Leary et al. (2019)

Key Insights on Faster Production as indicated in table 3:

1. Balancing Speed & Quality:
 - Example: Higher laser power (200-300W) combined with optimized scanning speeds (1000-1500 mm/s) can reduce print time by 35% without affecting precision (DebRoy et al., 2018).
2. Reducing Printing Defects for Efficiency:
 - Optimized cooling rates prevent microcracks, reducing the need for post-processing and increasing overall throughput (Yang et al., 2020).
3. Industry Case Study: Aerospace & Automotive Sectors
 - Boeing & Tesla use optimized laser power and hatch spacing to speed up metal AM production while maintaining part strength and aerodynamics (Leary et al., 2019).
4. Future Trends: AI & Machine Learning in Process Optimization
 - AI-driven systems adjust process parameters in real-time, reducing waste and boosting efficiency by up to 50% (Zhang et al., 2021).

2.4. Automation and Robotics in Additive Manufacturing (AM) Workflows

Automation and robotics are transforming additive manufacturing (AM) by enhancing precision, reducing manual intervention, and improving production speed. Integrating robotic arms, AI-driven systems, and automated quality control in AM workflows streamlines operations, reduces human error, and enables mass customization (Gibson et al., 2021). The integration of automation and robotics in AM processes as shown in table 4 has further enhanced efficiency by reducing manual interventions, minimizing human errors, and enabling continuous production cycles (Ford & Despeisse, 2016).

Industrial robots equipped with automated build plate handling and post-processing systems can seamlessly transition from printing to finishing, reducing production lead times. For example, robotic arms in metal AM have been used to remove support structures and smooth surfaces, reducing post-processing times by 40% (Garcia et al., 2020).

Additionally, automated powder handling systems in powder bed fusion (PBF) technologies have improved material recycling efficiency, lowering production costs while maintaining high-quality outputs (Kellens et al., 2017).

Table 4: Key Areas of Automation & Robotics in Additive Manufacturing AM

Automation Aspect	Description	Impact on AM Efficiency	References
Robotic Arms in Printing	Robots assist in material extrusion and direct energy deposition (DED) for large-scale structures.	Enhances precision and repeatability , enabling complex geometries .	Williams et al. (2019)
Automated Material Handling	AI-driven conveyor and robotic loading systems transport materials seamlessly.	Reduces waste, errors, and downtime in multi-material AM.	Leary et al. (2020)
Real-Time Process Monitoring	AI-driven sensors track temperature, laser power, and layer fusion quality during printing.	Detects defects early , reducing waste and rework costs .	Zhang et al. (2022)
Automated Post-Processing	Robots perform support removal, polishing, and surface finishing after printing.	Eliminates manual post-processing bottlenecks .	DebRoy et al. (2018)
AI-Powered Workflow Optimization	Machine learning adjusts print speed, energy input, and cooling rates in real-time.	Increases print speed by up to 50% while ensuring quality.	Liu et al. (2021)
Collaborative Robots (Cobots)	Human-robot collaboration allows safe and efficient AM operation .	Reduces operator fatigue and enhances precision.	Espinosa et al. (2021)

Key Insights on Robotics & Automation in AM as shown in table 4:

1. Enhancing Speed & Precision:
 - Example: GE Aviation integrates robotic laser-based AM to manufacture jet engine components with sub-micron accuracy (Williams et al., 2019).
2. Reduction in Waste & Rework:

- AI-driven defect detection systems analyze real-time data to adjust parameters, reducing material waste by 30-40% (Zhang et al., 2022).
- 3. Industrial Case Study: Aerospace & Automotive
 - Boeing & Tesla use robotic arms for automated large-scale AM, significantly reducing labor costs and improving consistency (Leary et al., 2020).
- 4. Future Trends: Fully Autonomous AM Factories
 - Emerging lights-out manufacturing facilities use robots, AI, and IoT for 24/7 autonomous AM production (Liu et al., 2021).

2.5. AI-Driven Defect Detection and Quality Control in Additive Manufacturing

The integration of Artificial Intelligence (AI) in defect detection and quality control is revolutionizing additive manufacturing (AM) by enhancing accuracy, reducing waste, and improving production efficiency table 5. AI-driven systems use machine learning (ML), deep learning (DL), and computer vision to identify defects in real time, predict failures, and optimize process parameters (Zhang et al., 2022).

One of the major challenges in AM is ensuring consistent print quality across batches. AI and machine learning algorithms are now being employed for real-time defect detection, significantly improving process efficiency (Tlegenov, Lu, & Xu, 2021).

By using in-situ monitoring systems, AI-driven models can analyze layer-by-layer build data and detect anomalies such as cracks, porosity, and dimensional inaccuracies (Zhang et al., 2021). These systems enable adaptive process corrections, where laser intensity or scanning speed can be adjusted mid-print to rectify errors before they compromise the final product.

For instance, deep learning models have improved defect detection rates by up to 98% in metal AM applications, leading to substantial cost savings and increased productivity (Bogue, 2020).

Table 5: AI-Driven Defect Detection and Quality Control in Additive Manufacturing

AI Technique	Function in AM Quality Control	Benefits	References
Machine Learning (ML)	Analyzes past defect data to predict new failures.	Reduces scrap rates and optimizes material usage.	Wang et al. (2021)
Deep Learning (DL)	Uses neural networks to identify micro-scale defects	Increases defect detection accuracy up	Gao et al. (2020)

	in printed layers.	to 95%.	
Computer Vision	Captures high-resolution images to detect layer-wise inconsistencies.	Enables real-time monitoring and automatic corrections.	Leary et al. (2020)
AI-Powered Process Monitoring	Adjusts laser power, scanning speed, and material feed rate in real-time.	Minimizes deviations and improves final product quality.	Yang et al. (2021)
X-Ray & Thermal Imaging Analysis	AI scans internal defects using non-destructive techniques.	Identifies voids, cracks, and porosity without damaging parts.	Zhang et al. (2022)
AI for Predictive Maintenance	Predicts machine wear and tear before failures occur.	Reduces downtime and increases machine lifespan.	Liu et al. (2019)

Key AI Techniques for Defect Detection in Additive Manufacturing as shown in table 5:

1. Machine Learning for Pattern Recognition
 - ML algorithms analyze large datasets of previous defects to predict and prevent recurring errors in AM (Wang et al., 2021).
 - Example: In powder bed fusion (PBF), AI predicts optimal laser parameters to prevent overheating and distortion (Gao et al., 2020).
2. Deep Learning for Image-Based Defect Identification
 - Convolutional Neural Networks (CNNs) detect porosity, cracks, and layer misalignment in real-time (Leary et al., 2020).
 - Example: A deep learning model reduced AM defect rates by 85% in a recent aerospace study (Yang et al., 2021).
3. Computer Vision for Real-Time Monitoring
 - AI-powered cameras capture high-resolution images of each printed layer, allowing for instant corrections (Gibson et al., 2021).
 - Example: Tesla’s AI-driven AM system automatically adjusts material flow based on real-time defect detection (Zhang et al., 2022).
4. AI in Non-Destructive Testing (NDT)

- AI analyzes X-ray and thermal images to detect internal voids and porosity without damaging the part (Liu et al., 2019).
- Example: Boeing uses AI-powered X-ray imaging to verify AM-produced aircraft components before deployment (Leary et al., 2020).

2.5.1. Industrial Applications & Case Studies

- Boeing & Airbus: Use AI-driven quality control in AM to reduce defects in aircraft components by 40% (Zhang et al., 2022).
- NASA: Uses machine learning-based thermal imaging to detect voids in 3D-printed rocket nozzles, increasing structural integrity (Wang et al., 2021).
- Automotive Industry (Tesla & BMW): AI-powered real-time defect detection in metal 3D printing improves precision and consistency (Gibson et al., 2021).

2.5.6. Future Trends in AI-Driven Quality Control

- Self-Healing AI Models: AI will autonomously adjust AM process parameters to prevent defects before they occur (Liu et al., 2019).
- Quantum AI for AM: Advanced quantum computing algorithms will enhance defect detection speed and accuracy (Yang et al., 2021).
- Fully Automated AM Quality Systems: AI-powered "lights-out" AM factories will run 24/7 without human intervention (Zhang et al., 2022).

3.0. Sustainability in Additive Manufacturing: Advancing Eco-Friendly Production

Additive manufacturing (AM), commonly known as 3D printing, is revolutionizing sustainable production by minimizing material waste, reducing energy consumption, and enabling localized manufacturing (Gibson et al., 2021). Compared to traditional subtractive manufacturing, which removes excess material, AM builds products layer by layer, using only the necessary material. This resource-efficient approach aligns with global sustainability goals such as carbon footprint reduction, circular economy practices, and eco-friendly material utilization (Thompson et al., 2022).

Table 6: Sustainable Practices in Additive Manufacturing

Sustainability Aspect	Key Benefits	Industry Examples	References
Material Efficiency	Reduces raw material waste by up to 90%.	Boeing & Airbus use AM to print lightweight aircraft parts , cutting	Yang et al. (2023)

		fuel consumption.	
Energy Consumption Reduction	AM uses 30–50% less energy than conventional machining.	GE’s AM-powered jet engine nozzles consume 35% less energy during production.	Gao et al. (2021)
Recycling & Circular Economy	Enables recycling of metal powders, biopolymers, and composite materials.	HP’s Multi Jet Fusion printers recycle 80% of used powder.	Kumar & Rosen (2020)
Localized Manufacturing & Supply Chain Optimization	Reduces global transportation emissions by 50–60% .	NASA uses AM to 3D print tools in space , eliminating supply chain dependency.	Zhang et al. (2022)
Biodegradable & Bio-Based Materials	Supports PLA (polylactic acid), biocomposites, and algae-based resins.	Adidas uses AM for eco-friendly, biodegradable shoe soles.	Torres et al. (2021)
Lightweight & High-Strength Design	Reduces material use and fuel consumption in transport sectors.	Tesla’s AM-designed battery components improve energy efficiency.	Wang et al. (2021)

3.1. Sustainable Materials in Additive Manufacturing

3.1.1. Material Efficiency and Waste Reduction

AM significantly reduces material waste compared to traditional subtractive manufacturing processes (Baumers et al., 2016). By precisely depositing material where needed, AM eliminates excess machining and scrap production.

Recycled materials, such as reclaimed polymers and metal powders, are increasingly used in AM, further promoting circular economy principles (Tlegenov et al., 2021).

3.1.1.1. Recyclable Metal Powders for Sustainable AM

- Aluminum, titanium, and stainless steel powders can be reused multiple times, reducing material waste (Gao et al., 2021).
- Example: Airbus implemented recycled titanium powder in AM to cut waste by 50% in aircraft component production (Kumar & Rosen, 2020).

3.1.1.2. Biodegradable Polymers for Eco-Friendly Printing

- Polylactic Acid (PLA) is a plant-based, biodegradable polymer widely used in 3D printing (Torres et al., 2021).
- Example: Adidas 3D prints sneakers from PLA and algae-based filaments, promoting sustainable fashion (Zhang et al., 2022).

3.1.1.3. Bio-Based and Composite AM Materials

- Research is advancing cellulose-based, algae-infused, and bio-resin materials to replace synthetic plastics (Yang et al., 2023).
- Example: Researchers at MIT developed wood-based AM filaments, reducing plastic waste by 60% (Gibson et al., 2021).

3.2. Energy Efficiency in Additive Manufacturing

Energy consumption in AM is typically lower than in conventional manufacturing due to the elimination of tooling and machining operations (Peng et al., 2018). Studies have shown that AM can reduce energy consumption by up to 50% in certain applications, particularly in aerospace and automotive industries (Kellens et al., 2017).

Additionally, decentralized AM production reduces transportation-related emissions by enabling localized manufacturing (Garcia et al., 2020).

- Laser-based AM processes, such as Selective Laser Melting (SLM) and Electron Beam Melting (EBM), use less energy compared to traditional casting and milling (Wang et al., 2021).
- Case Study: General Electric (GE) reduced energy consumption by 35% in jet engine production using AI-optimized AM processes (Kumar & Rosen, 2020).
- Sustainable AM Factories: Companies are now integrating solar power and AI-driven energy management for carbon-neutral AM production (Thompson et al., 2022).

3.3. Circular Economy and Waste Reduction Strategies

- Powder Bed Fusion (PBF) recycling: Unused metal powders from AM processes can be collected, sieved, and reused (Gao et al., 2021).
- Closed-loop AM systems: Companies like HP and Stratasys recycle used polymer powders to minimize landfill waste (Torres et al., 2021).
- Case Study: Boeing developed a recycling system for excess titanium powder, reducing raw material costs by 50% (Zhang et al., 2022).

3.4. Sustainable Supply Chains

AM facilitates localized production, reducing dependency on global supply chains and minimizing transportation emissions (Ford & Despeisse, 2016). This decentralization enhances supply chain resilience and supports sustainable industrial practices.

For instance, during the COVID-19 pandemic, AM was crucial in rapidly producing medical components, such as ventilator parts, to meet urgent demands (Bogue, 2020).

4.0. Case Studies on Sustainability in Additive Manufacturing

4.1. Case Study 1: Airbus – Sustainable Lightweight Aircraft Components

Airbus, a global aerospace leader, has integrated additive manufacturing (AM) to reduce material waste, fuel consumption, and carbon footprint. Table 7 shows how the company employs 3D printing for lightweight aircraft components, ensuring both structural efficiency and sustainability (Kumar & Rosen, 2020).

Table 7: Sustainability Innovations in Airbus' Additive Manufacturing Strategy

Innovation	Sustainability Impact	Outcome	Reference
Lightweight AM-printed brackets	Used topology optimization to reduce part weight by 30%	Lower aircraft weight results in fuel savings and reduced CO ₂ emissions	Zhang et al. (2022)
Recycled titanium powder in AM	Reduced raw material waste by 50%	Less dependence on mined titanium, lowering environmental impact	Wang et al. (2021)
3D-printed spare parts for aircraft maintenance	Eliminated global shipping for spare parts	40% reduction in logistics-related carbon emissions	Torres et al. (2021)
Bionic-designed AM components	Inspired by nature-inspired structures to maximize material efficiency	Reduced material usage and energy consumption in part production	Gao et al. (2021)

Sustainability Achievements include the following;

- Airbus' A320neo aircraft, equipped with AM-optimized fuel nozzles and brackets, reduces CO₂ emissions by 5,000 tons per aircraft per year (Zhang et al., 2022).
- The company has eliminated 60% of machining waste by transitioning from subtractive to additive manufacturing (Kumar & Rosen, 2020).

4.2. Case Study 2: General Electric (GE) – AM for Sustainable Jet Engine Production

General Electric (GE) Aviation has pioneered AM for jet engine parts, aiming for higher fuel efficiency and reduced material waste. Table 8 indicates the company’s most notable success is the LEAP fuel nozzle, a 3D-printed jet engine component that replaced traditional 20-part assemblies with a single AM unit (Thompson et al., 2022).

Table 8: Sustainable Additive Manufacturing Initiatives by General Electric GE

AM Innovation	Sustainability Benefit	Impact on Aviation Industry	Reference
AM-printed LEAP fuel nozzle	Uses 25% less material , making it lighter and more fuel-efficient	Reduces aircraft fuel consumption by 15%	Kumar & Rosen (2020)
Reduction in jet engine component count	Consolidated 20 parts into 1 using AM	Lowers material waste, reducing CO ₂ emissions during production	Gao et al. (2021)
Energy-efficient AM process for titanium parts	Reduces energy use by 30% compared to machining	Cuts carbon footprint and lowers production costs	Torres et al. (2021)
On-site AM for maintenance and repairs	Eliminates need for shipping replacement parts	Reduces aviation logistics-related CO ₂ emissions	Wang et al. (2021)

Sustainability Achievements are;

- GE’s AM-printed LEAP fuel nozzle saves 100,000 metric tons of CO₂ annually by improving engine fuel efficiency (Thompson et al., 2022).
- The company’s on-site AM repair system has cut jet engine maintenance costs by 40% while minimizing material waste (Zhang et al., 2022).

4.3. Case Study 3: Adidas – 3D-Printed Sustainable Footwear

Adidas has leveraged additive manufacturing to create sustainable, lightweight, and customizable shoe soles, reducing material waste and promoting a circular economy (Gibson et al., 2021). The Futurecraft 4D shoe line is a prime example of AM-driven sustainability in fashion table 9.

Table 9: Sustainable Additive Manufacturing Practices in Adidas Footwear Production

Additive Manufacturing-Based Sustainability Innovation	Eco-Friendly Impact	Industry Contribution	Reference
3D-printed midsoles using Carbon's Digital Light Synthesis (DLS)	Reduces material waste by 60%	Enhances shoe durability and longevity	Zhang et al. (2022)
Use of recycled and bio-based materials	Lowers dependence on petroleum-based plastics	Introduces sustainable polymers like biodegradable TPU	Kumar & Rosen (2020)
Localized Additive Manufacturing production for on-demand manufacturing	Eliminates over production and excess inventory	Lowers logistics-related emissions	Torres et al. (2021)
AM-optimized lattice structures for lightweight design	Reduces material usage while enhancing comfort	Creates strong yet lightweight footwear	Wang et al. (2021)

Sustainability Achievements include;

- Adidas' Futurecraft 4D shoes use 50% less material compared to conventional footwear, reducing landfill waste (Gibson et al., 2021).
- The company has reduced global shipping emissions by 35% by adopting localized 3D printing hubs (Thompson et al., 2022).

4.4. Case Study4: National Aeronautics and Space Administration (NASA) – Sustainable Additive Manufacturing for Space Missions

National Aeronautics and Space Administration (NASA) has embraced Additive Manufacturing for in-space manufacturing, reducing reliance on Earth-bound supply chains and minimizing material waste. The agency uses 3D printing to produce tools, spare parts, and even habitats for future Mars missions (Gibson et al., 2021).

Table 10: National Aeronautics and Space Administration NASA's Sustainable Additive Manufacturing Innovations

Innovation	Sustainability Benefit	Impact on Space Missions	Reference
In-space 3D printing of tools and spare parts	Reduces reliance on Earth-based supply chains	Lowers launch mass and cuts resupply mission costs	Gao et al. (2021)
Recyclable AM materials for space applications	Uses recycled polymers from used packaging	Promotes circular economy in space missions	Torres et al. (2021)
3D-printed Mars habitat structures using in-situ materials	Eliminates need for transporting construction materials	Reduces environmental impact and launch costs by 80%	Wang et al. (2021)
AM-printed metal components for spacecraft	Uses lighter materials with high strength	Reduces fuel consumption for space travel	Thompson et al. (2022)

Sustainability Achievements are;

- NASA’s in-space Additive Manufacturing systems have cut logistical costs by \$500 million annually by reducing dependence on Earth-based resupply missions (Zhang et al., 2022).
- Recycled 3D printing materials have enabled sustainable waste repurposing in space (Gibson et al., 2021).

4.5. Case Study 5: Adidas’ 3D-Printed Footwear – Reducing Material Waste

Adidas has leveraged AM to develop sustainable footwear using Carbon's Digital Light Synthesis (DLS) technology. The Futurecraft 4D shoe line is an example of how AM enables precise material use and reduced waste.

Key Material Efficiency Strategies

- Lattice-structured midsoles: Uses 30% less material while maintaining durability.
- Recycled TPU material: Reduces reliance on virgin polymers.
- Localized AM production: Cuts logistics emissions by 35%.

By shifting to 3D-printed midsoles, Adidas has achieved a 50% reduction in raw material waste, making its AM footwear one of the most sustainable innovations in the fashion industry (Thompson et al., 2022).

4.6. Case Study: Boeing’s Additive Manufacturing (AM) Efficiency Innovations

Boeing has been a pioneer in integrating additive manufacturing (AM) into aerospace production, significantly improving efficiency, sustainability, and cost-effectiveness. The company utilizes 3D printing for aircraft parts, reducing material waste, production time, and fuel consumption. Below are key areas where Boeing has enhanced efficiency using AM:

4.6.1. Material Waste Reduction (50% Reduction in Waste)

- Traditional machining involves subtractive processes (cutting, milling, drilling), leading to up to 90% material waste (DebRoy et al., 2018).
- With AM, Boeing precisely deposits material layer by layer, leading to a 50% reduction in waste compared to traditional methods (Schmid et al., 2019).
- This translates to lower raw material costs and reduced environmental impact.

4.6.2. 3D-Printed Components (60,000+ Parts in Aircraft)

- Boeing has incorporated over 60,000 3D-printed components in commercial and military aircraft (Leary et al., 2019).
- Examples include:
 - Cabin brackets and ducting (lightweight but strong).
 - Titanium structural components (used in the Boeing 787 Dreamliner).
- These parts contribute to lower fuel consumption and operational efficiency.

4.6.3. Faster Production Cycles (30% Reduction in Time)

- Additive Manufacturing eliminates several time-intensive steps (tooling, assembly, and machining).
- Boeing's AI-integrated Additive Manufacturing systems have reduced production cycles by 30%, enabling faster part production and shorter aircraft assembly times (Bogue, 2020).
- Example:
 - Traditional aerospace part production takes weeks or months.
 - With 3D printing, it can be reduced to days or hours.

4.6.4. Lightweight Structures (Fuel-Efficient Designs)

- Boeing uses topology optimization and biomimetic designs to produce lightweight yet strong aircraft parts.
- Example:
 - The Boeing 787 Dreamliner incorporates 3D-printed titanium parts that reduce aircraft weight, improving fuel efficiency by 20% (Peng et al., 2018).

- These lightweight structures reduce fuel consumption, operational costs, and carbon footprint.

4.6.5. AI & Robotics Integration (Automated Quality Control)

- Boeing integrates robotic systems and AI-driven defect detection to enhance AM efficiency.
- Automated real-time monitoring ensures:
 - Early defect detection, reducing scrap rates.
 - Precise material deposition, improving component integrity.
- AI-based quality control has improved defect detection rates by 98%, ensuring high reliability (Zhang et al., 2021).

5.0. Challenges & Future Prospects for Sustainable Additive Manufacturing

5.1. Current Challenges

- High Energy Consumption for Metal AM: Metal AM processes, such as SLM and EBM, still require significant power (Thompson et al., 2022).
- Recycling Limitations: Some polymer and composite materials degrade after multiple recycling cycles (Gibson et al., 2021).
- Cost of Sustainable Materials: Biodegradable AM materials are still expensive compared to traditional plastics (Torres et al., 2021).

5.2. Future Developments

- AI-Optimized Energy Use: AI-powered AM processes will reduce energy consumption by up to 50% in the next decade (Yang et al., 2023).
- Advancements in Biodegradable & Recyclable Materials: New research into bio-based and algae-infused AM materials will enhance eco-friendliness (Zhang et al., 2022).
- Green AM Factories: Integration of renewable energy sources in AM production will make factories carbon-neutral by 2035 (Wang et al., 2021).

Conclusion

Additive Manufacturing (AM) is revolutionizing mechanical production by enhancing efficiency, material utilization, and sustainability. Through advancements such as topology optimization, multi-material printing, and AI-driven quality control, AM enables the creation of lightweight, high-strength, and functionally optimized components. Its ability to minimize waste, reduce energy consumption, and facilitate localized manufacturing aligns with the principles of sustainable production. Case

studies from leading industries, including aerospace and automotive sectors, demonstrate AM's potential to significantly lower carbon footprints and production costs. Despite challenges such as material limitations and production speed constraints, the integration of AI and automation continues to drive improvements, making AM a viable solution for large-scale manufacturing. Future developments in AI-powered self-healing systems, quantum computing for defect detection, and sustainable material innovations will further solidify AM's role in shaping the future of industrial production. As industries continue to adopt and refine AM technologies, its impact on efficiency, customization, and environmental sustainability will become increasingly profound, cementing its status as a cornerstone of modern manufacturing.

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Publications

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Enhancing Concrete Performance: The Impact of Polymer Additives

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Abstract

Concrete remains a fundamental material in modern construction, yet traditional concrete exhibits limitations such as low tensile strength, high permeability, and susceptibility to cracking. The integration of polymer additives into concrete has emerged as an effective strategy to enhance its mechanical properties, durability, and sustainability. Polymer-modified concrete (PMC) and polymer cement concrete (PCC) utilize polymers in various forms, including latexes, liquid resins, and water-soluble compounds, to improve workability, tensile strength, flexural strength, impact resistance, and durability. Research indicates that polymer addition densifies the interfacial transition zone (ITZ), reduces porosity, and enhances mechanical performance. Additionally, self-healing properties imparted by superabsorbent polymers (SAPs) further contribute to extended service life. Despite certain challenges, such as potential reductions in compressive strength with some polymer types, the benefits of polymer incorporation outweigh the drawbacks. This paper comprehensively examines the quantitative and qualitative effects of polymer additives on concrete's performance, including enhanced durability, permeability reduction, and self-healing capabilities. The findings underscore the potential of polymer-modified concrete as a viable solution for sustainable and high-performance construction materials.

Key words: *Polymer Additives, Concrete Performance, Tensile Strength, Self-Healing, Durability, Impact Resistance*

1.0 Introduction

Concrete, a composite material composed of fine and coarse aggregates bonded together with a fluid cement, has been a cornerstone in construction due to its versatility and compressive

strength. However, traditional concrete exhibits limitations such as low tensile strength, susceptibility to cracking, and permeability issues, which can compromise durability and longevity. To address these challenges, the incorporation of polymers into concrete mixtures has been explored, leading to the development of Polymer Modified Concrete (PMC) and Polymer Cement Concrete (PCC). These innovations aim to enhance the mechanical properties, durability, and overall performance of concrete structures (Krause et al., 2019).

The concept of integrating polymers into concrete dates back to the mid-20th century, with significant advancements occurring in the subsequent decades. Polymers can be introduced in various forms, including latexes, liquid resins, and water-soluble polymers, each contributing uniquely to the properties of the resulting composite material. The selection of polymer type and dosage is critical, as it influences factors such as workability, setting time, and the hardened properties of the concrete (Nawy, 2012).

Research has demonstrated that polymer addition can significantly improve the tensile and flexural strengths of concrete. For instance, the inclusion of polymer emulsions has been shown to densify the interfacial transition zone (ITZ) and fill pores within the cement matrix, leading to enhanced mechanical performance (Kujawa et al., 2021). Moreover, polymers contribute to reduced permeability, thereby limiting the ingress of harmful substances like chlorides and sulfates, which are known to cause reinforcement corrosion and concrete degradation. This reduction in permeability enhances the durability of concrete structures, especially in aggressive environmental conditions (Kujawa et al., 2021; Rashad, 2015).

In addition to mechanical enhancements, certain polymers impart self-healing capabilities to concrete. Superabsorbent polymers (SAPs), for example, can absorb water and swell, facilitating the closure of microcracks and promoting further hydration of unreacted cement particles. This autogenous healing process can restore mechanical properties and extend the service life of concrete structures (Snoeck et al., 2012; Guo et al., 2021).

However, the incorporation of polymers into concrete is not without challenges. The type and amount of polymer used can affect the fresh and hardened properties of concrete. For example, the addition of shredded recycled polymer fibers has been observed to reduce compressive strength slightly but can significantly mitigate shrinkage cracking (Kujawa et al., 2021). Therefore, optimizing the polymer content is essential to achieve a balance between workability, strength, and durability.

The addition of polymers to concrete represents a significant advancement in construction materials, offering improvements in strength, durability, and functionality. Ongoing research

and development continue to refine these materials, making them more adaptable to various construction demands and environmental conditions.

2.0 Enhancement of Mechanical Properties through Polymer Addition in Concrete

The incorporation of polymers into concrete has been extensively studied to improve its mechanical properties, including compressive strength, tensile strength, flexural strength, and impact resistance. These enhancements are primarily attributed to the modification of the cement matrix and the interfacial transition zone (ITZ) between the cement paste and aggregates.

2.1 Compressive Strength

The effect of polymer addition on the compressive strength of concrete varies depending on the type and amount of polymer used. Some studies have reported a slight reduction in compressive strength with the incorporation of certain polymers, such as shredded recycled polymer fibers, which may reduce compressive strength by approximately 0.5% to 8.5% (Kujawa et al., 2021). However, other research indicates that specific polymer additives can enhance compressive strength by refining the pore structure and increasing the density of the cement matrix (Daneti & Meddah, 2013). Additionally, the use of polymers in cement mortar can improve workability with a lower water-cement ratio, contributing to strength development (Kumar & Prakash, 2015).

2.2 Tensile and Flexural Strength

Polymers significantly enhance the tensile and flexural strengths of concrete by improving the bond between the cement paste and aggregates. This improvement is due to the formation of a polymer film within the cement matrix, which increases cohesion and reduces the propagation of microcracks under load (Ohama, 1995). Studies have shown that polymer-modified concrete exhibits higher flexural strength compared to conventional concrete, making it suitable for applications where improved tensile performance is required (Kujawa et al., 2021). Furthermore, the incorporation of polymers has been demonstrated to enhance flexural strength and modulus of elasticity, contributing to better resistance against mechanical stresses (Kakade et al., 2023).

2.3 Impact Resistance

The addition of polymers contributes to improved impact resistance in concrete. The enhanced toughness is attributed to the energy-absorbing capacity of the polymer phase, which helps in dissipating impact energy and preventing crack propagation (Daneti &

Meddah, 2013). This property is particularly beneficial in structures subjected to dynamic or impact loads.

2.4 Modulus of Elasticity

The incorporation of polymers can influence the modulus of elasticity of concrete. Generally, polymer-modified concrete exhibits a lower modulus of elasticity compared to conventional concrete, which can be advantageous in reducing brittleness and enhancing deformation capacity under load (Ohama, 1995).

The addition of polymers to concrete significantly enhances its mechanical properties, such as compressive strength, tensile strength, flexural strength, impact resistance, and durability. The table 1 below summarizes these enhancements and their effects.

Table 1: Enhancement of Mechanical Properties of Polymer-Modified Concrete

Property	Enhancement Due to Polymer Addition	References
Compressive Strength	Some polymers improve compressive strength by refining the pore structure and increasing matrix density. However, certain polymer additions, such as recycled plastic fibers, can slightly reduce compressive strength due to reduced interfacial bonding.	Kujawa et al. (2021); Daneti & Meddah (2013); Kumar & Prakash (2015)
Tensile Strength	Polymer addition enhances tensile strength by improving cohesion between cement and aggregates. This effect is due to the formation of a polymer film that resists crack propagation.	Ohama (1995); Kakade et al. (2023)
Flexural Strength	The flexibility of polymer-modified concrete is increased, leading to higher flexural strength. This is particularly useful for bridge decks and pavements.	Kujawa et al. (2021); Kumar & Prakash (2015)
Impact Resistance	Polymers enhance the energy absorption capacity of concrete, making it more resistant to impact forces.	Daneti & Meddah (2013); Kakade et al. (2023)
Modulus of Elasticity	Polymer-modified concrete has a lower modulus of elasticity, which reduces brittleness and	Ohama (1995); Kakade et al. (2023)

	improves ductility, allowing better resistance to cracking.	
Shrinkage and Cracking Resistance	Polymers reduce drying shrinkage and cracking due to better hydration control and flexibility. This is particularly beneficial in reducing plastic shrinkage cracking.	Kujawa et al. (2021); Kumar & Prakash (2015)
Water Permeability	Polymers decrease permeability by forming a dense network, reducing the ingress of harmful substances such as chlorides and sulfates.	Kakade et al. (2023); Daneti & Meddah (2013)

The above table 1 revealed that the incorporation of polymers into concrete significantly enhances mechanical performance, improving durability and resistance to external stresses. The choice of polymer type, dosage, and curing process plays a critical role in optimizing these benefits.

The addition of polymers to concrete provides significant benefits in terms of tensile strength, flexural strength, impact resistance, shrinkage resistance, and permeability reduction. However, certain polymer types, particularly recycled plastic fibers, may slightly reduce compressive strength. The careful selection of polymer type and dosage is crucial to optimizing the desired properties for specific construction applications.

3.0 Quantitative Effects of Polymer Addition on Concrete's Mechanical Properties

The incorporation of polymers into concrete can significantly influence its mechanical properties. The specific effects depend on factors such as the type and amount of polymer used, mix design, and curing conditions. The table 2 below summarizes the quantitative changes in key mechanical properties due to polymer addition.

Table 2: Quantitative Effects of Polymer Addition on Concrete's Mechanical Properties

Property	Effect of Polymer Addition	Reference
Compressive Strength	Increase: Addition of 3% polymer emulsion densifies the interfacial transition zone (ITZ) and fills pores, leading to increased compressive strength.	Kakade et al. (2023)
	Decrease: Incorporation of shredded	Kujawa et al. (2021)

	recycled polymer fibers can reduce compressive strength by approximately 0.5% to 8.5%, depending on fiber content and type.	
Tensile Strength	Increase: Polymer-modified concrete shows a tensile splitting strength increase of up to 255% compared to conventional concrete.	Kakade et al. (2023)
Flexural Strength	Increase: The addition of polymers enhances flexural strength, making concrete more suitable for applications requiring improved tensile performance.	Kujawa et al. (2021)
Impact Resistance	Increase: Polymers improve impact resistance by enhancing the energy-absorbing capacity of the concrete, preventing crack propagation.	Daneti & Meddah (2013)
Modulus of Elasticity	Decrease: Polymer incorporation can lead to a lower modulus of elasticity, reducing brittleness and enhancing deformation capacity under load.	Ohama (1995)
Shrinkage and Cracking Resistance	Increase: Polymers reduce drying shrinkage and cracking due to better hydration control and flexibility, which is beneficial in reducing plastic shrinkage cracking.	Kujawa et al. (2021)
Water Permeability	Decrease: Polymers decrease permeability by forming a dense network, reducing the ingress of harmful substances such as chlorides and sulfates.	Kakade et al. (2023)

Below is a detailed examination of how polymer additions quantitatively affect key mechanical properties of concrete:

3.1. Compressive Strength

Polymer incorporation has a dual effect on the compressive strength of concrete. According to Kakade et al. (2023), adding 3% polymer emulsion improves compressive strength by densifying the interfacial transition zone (ITZ) and filling pores, thereby reducing porosity and increasing strength. However, the use of shredded recycled polymer fibers can lead to a reduction in compressive strength by 0.5% to 8.5%, depending on fiber type and content (Kujawa et al., 2021). This reduction occurs because the inclusion of polymer fibers can disrupt the uniformity of the concrete matrix, reducing its load-bearing capacity.

3.2. Tensile Strength

Polymer modification has been shown to significantly enhance the tensile strength of concrete. Kakade et al. (2023) reported that polymer-modified concrete exhibits a tensile splitting strength increase of up to 255% compared to conventional concrete. This improvement is attributed to the better adhesion and flexibility imparted by polymers, which help distribute tensile stresses more effectively within the matrix.

3.3. Flexural Strength

The addition of polymers also enhances the flexural strength of concrete, making it more suitable for applications requiring improved tensile performance. Kujawa et al. (2021) demonstrated that polymer-modified concrete exhibits higher resistance to bending forces, leading to improved structural integrity in flexural applications.

3.4. Impact Resistance

Polymers significantly improve the impact resistance of concrete by enhancing its energy-absorbing capacity. Daneti and Meddah (2013) highlighted that polymer-modified concrete is better at resisting crack propagation under dynamic loads. This makes polymer-enhanced concrete ideal for structures subjected to high impact or cyclic loading conditions.

3.5. Modulus of Elasticity

A notable drawback of polymer incorporation is its effect on the modulus of elasticity. Ohama (1995) noted that polymer-modified concrete tends to have a lower modulus of elasticity, which reduces its brittleness and increases its capacity for deformation under load. While this can be advantageous in certain applications, it may be a limitation in structures requiring high stiffness.

3.6. Shrinkage and Cracking Resistance

Polymers play a crucial role in reducing drying shrinkage and cracking in concrete. Kujawa et al. (2021) found that polymer-modified concrete exhibits better hydration control and increased flexibility, which significantly reduces plastic shrinkage cracking. This makes polymer-modified concrete particularly useful in applications where shrinkage-induced cracking is a concern.

3.7. Water Permeability

One of the most important benefits of polymer addition is the reduction in water permeability. According to Kakade et al. (2023), polymers help form a dense network within the concrete matrix, reducing the ingress of harmful substances such as chlorides and sulfates. This leads to enhanced durability, making polymer-modified concrete more resistant to environmental degradation.

The addition of polymers to concrete can significantly enhance properties such as tensile strength, flexural strength, impact resistance, and shrinkage resistance. However, potential drawbacks include reductions in compressive strength (depending on polymer type) and a lower modulus of elasticity. The specific effects of polymer incorporation depend on the type, quantity, and dispersion of the polymer within the matrix. As research in polymer-modified concrete advances, optimizing polymer content and composition will be critical in maximizing its benefits for sustainable and high-performance construction materials.

The quantitative effects of polymer addition on concrete's mechanical properties vary based on the type and amount of polymer used, as well as the specific application requirements. While certain polymers can enhance properties like tensile and flexural strength, others may lead to reductions in compressive strength. Therefore, careful selection and optimization of polymer types and dosages are essential to achieve the desired performance in concrete applications.

4.0 Impact of Polymer Addition on Concrete's Durability and Permeability

The incorporation of polymers into concrete mixtures significantly influences durability and permeability characteristics. The extent of these effects varies based on the type and amount of polymer used, as well as specific mix designs and curing conditions. The table 3 below provides a statistical overview of the impact of polymer addition on concrete's durability and permeability, along with relevant references.

Table 3: Impact of Polymer Addition on Concrete's Durability and Permeability

Property	Effect of Polymer Addition	Reference
Durability	- Increase: Incorporation of 3% polymer emulsion leads to densification of the interfacial transition zone (ITZ) and pore filling, resulting in enhanced durability, including increased freeze–thaw resistance and restraint to the propagation of microdefects.	Kakade et al. (2023)
	- Enhancement: Polymer-incorporated concrete exhibits exceptionally good performance in terms of durability when exposed to harsh environments.	Daneti & Meddah (2013)]
Permeability	- Reduction: Addition of 3% polymer emulsion reduces permeability by densifying the ITZ and filling pores, thereby decreasing the ingress of harmful substances.	Kakade et al. (2023)

The incorporation of polymers in concrete has been widely recognized for enhancing its durability and reducing permeability as shown in Table 3. These improvements contribute to the long-term performance of concrete structures, particularly in aggressive environments.

Durability

The durability of concrete is significantly improved by the addition of polymers. According to Kakade et al. (2023), incorporating 3% polymer emulsion leads to densification of the interfacial transition zone (ITZ) and pore filling, which enhances the resistance of concrete to freeze–thaw cycles and restrains the propagation of microdefects. Additionally, polymer-modified concrete performs exceptionally well in harsh environments, exhibiting higher resistance to degradation due to environmental exposure (Daneti & Meddah, 2013). This increased durability makes polymer-modified concrete suitable for infrastructure exposed to extreme weather conditions and chemical attack.

Permeability

Polymers play a crucial role in reducing the permeability of concrete. Kakade et al. (2023) found that adding 3% polymer emulsion significantly decreases permeability by densifying the ITZ and filling pores. This reduction in permeability limits the ingress of harmful substances such as chlorides and sulfates, improving the durability and service life of concrete structures. Reduced permeability is particularly beneficial in marine environments,

wastewater treatment plants, and underground constructions where exposure to aggressive agents is a concern.

The addition of polymers enhances the durability and impermeability of concrete, making it more resilient to environmental and chemical challenges. By refining the microstructure and limiting the penetration of harmful substances, polymer-modified concrete becomes a superior choice for long-lasting and sustainable construction applications.

5.0 Self-Healing Concrete

Self-healing concrete is an innovative material that has gained significant attention due to its ability to automatically repair cracks and improve its durability without external intervention. The incorporation of polymers and other additives into the concrete mix enhances its self-healing capabilities, which leads to longer service life, reduced maintenance costs, and improved performance in harsh environmental conditions.

5.1 Self-Healing Capacities in Polymer-Modified Concrete

The self-healing capacity in concrete arises from the concrete's ability to autonomously respond to damage, primarily cracks, by sealing them through various mechanisms. These mechanisms can be chemical, biological, or physical, and they help in reducing permeability, restoring strength, and enhancing durability. The addition of polymers can improve this self-healing process by:

5. **Microencapsulation:** Some polymer-modified concrete incorporates microcapsules containing healing agents such as calcium lactate or other chemical compounds. When a crack forms, these capsules rupture, releasing the healing agents that react with water or carbon dioxide to form precipitates, effectively closing the cracks and restoring the structural integrity of the material (Jonkers et al., 2010).
6. **Polymer-based Healing Agents:** Certain polymers, such as polyurethanes or epoxy resins, are incorporated into concrete as healing agents. When cracks develop, these polymers fill the voids and harden, thereby effectively sealing the cracks and reducing water ingress (Kakade et al., 2023). This reduces the potential for further degradation of concrete caused by chemical attack or freeze-thaw cycles.
7. **Enhanced Crack Sealing via Polymers:** Polymers increase the elasticity of concrete, making it less prone to cracking in the first place. Furthermore, some polymer-modified concretes have a higher capacity for crack bridging, which allows them to close small cracks more effectively. This process reduces water permeability and prevents corrosion of reinforcing steel (Daneti & Meddah, 2013).

8. **Moisture-activated Polymers:** Certain hydrophilic polymers, such as polyvinyl alcohol (PVA), enhance self-healing by reacting with water during the hydration process. These polymers absorb moisture and expand, filling in microcracks and improving the material's resistance to environmental stressors (Van Tittelboom et al., 2010).

6.0 Statistical Analysis of Self-Healing Performance in Polymer-Modified Concrete

The self-healing performance of polymer-modified concrete has been extensively studied, with various statistical analyses conducted to quantify its effectiveness. These analyses provide insights into the healing efficiency, mechanical property recovery, and durability enhancement of concrete incorporating polymers as shown in table 4.

Table 4: Statistical Analysis of Self-Healing Performance in Polymer-Modified Concrete

Test Conditions	Polymer Type/Concentration	Effect on Self-Healing Performance	Key Findings	Reference
28-day healing process under wet conditions	Microencapsulated healing agent (calcium lactate)	Crack width reduction by 65%, improved mechanical properties and durability in cracked concrete.	Significant self-healing observed in cracks up to 0.8 mm wide.	Jonkers et al. (2010)
Moisture-activated healing process, 28 days	Moisture-activated polymer (polyvinyl alcohol, PVA)	80% crack healing efficiency for cracks less than 1 mm in width, demonstrated high resistance to environmental stresses.	Moisture-triggered expansion of polymers healed microcracks.	Van Tittelboom et al. (2010)
Crack healing under water exposure (28 days)	Polyurethane-based polymer resin	Increased crack healing efficiency by 40%, improved resistance to chloride ingress.	Demonstrated reduction in permeability and enhanced durability.	Kakade et al. (2023)
Cyclic wet-dry and freeze-	Styrene-butadiene rubber (SBR)	50% recovery in compressive strength	SBR-modified concrete healed	Vasilenko et al. (2021)

thaw tests for 6 months		and reduced permeability.	cracks effectively, showing enhanced freeze-thaw resistance.	
Accelerated aging test, 28 days of exposure to cyclic wet and dry conditions	Styrene-butadiene emulsion (SBR)	Reduced permeability, enhanced mechanical properties, and crack closure with reduced maintenance requirements.	Significant reduction in chloride ion ingress.	Daneti & Meddah (2013)

Several studies have quantified the self-healing performance of polymer-modified concrete as indicated in the above table 4.

A study by **Jonkers et al. (2010)** demonstrated that concrete modified with microencapsulated healing agents showed a 65% reduction in crack width after 28 days of exposure to water, compared to conventional concrete. This indicated a substantial improvement in the material's self-healing capacity.

Kakade et al. (2023) found that incorporating polymeric microcapsules in concrete increased its crack-healing efficiency by 40%, enhancing durability under wet conditions and improving resistance to chloride-induced corrosion.

Van Tittelboom et al. (2010) reported that polymer-modified concrete with moisture-activated self-healing agents exhibited up to 80% healing efficiency in cracks up to 1 mm wide within a month of exposure to water, compared to just 30% for conventional concrete.

The self-healing capacity of polymer-modified concrete is a promising approach to enhance the durability and longevity of concrete structures. The use of polymer-based additives, such as microcapsules and hydrophilic agents, offers significant improvements in the material's ability to repair itself autonomously. This technology can potentially revolutionize the construction industry by reducing maintenance needs, extending service life, and minimizing environmental impact.

The incorporation of polymers in self-healing concrete has shown remarkable improvements in crack closure, permeability reduction, and durability under various environmental conditions. Different polymer types and concentrations influence the self-healing efficiency, primarily by promoting hydration, expansion, and chemical reactions that seal cracks.

6.1. Self-Healing Under Wet Conditions

The use of microencapsulated healing agents, such as calcium lactate, has demonstrated significant crack width reduction and durability improvement in cracked concrete. Jonkers et al. (2010) reported that a 28-day healing process under wet conditions resulted in a 65% reduction in crack width and enhanced mechanical properties, particularly in cracks up to 0.8 mm wide. This suggests that such self-healing mechanisms are effective in enhancing the longevity of concrete structures exposed to moisture.

6.2 Moisture-Activated Healing

Polyvinyl alcohol (PVA), a moisture-activated polymer, has shown an 80% crack healing efficiency for cracks smaller than 1 mm. Van Tittelboom et al. (2010) found that the expansion of PVA under moisture exposure played a crucial role in sealing microcracks, leading to improved resistance against environmental stresses. This highlights the potential of PVA-based self-healing systems in applications where concrete is subjected to varying humidity levels.

6.3. Self-Healing Under Water Exposure

Polyurethane-based polymer resins have also proven effective in crack healing, particularly in underwater environments. Kakade et al. (2023) demonstrated that these polymers improved crack healing efficiency by 40% while significantly reducing permeability and enhancing resistance to chloride ingress. This is particularly beneficial for marine and submerged concrete structures, where water penetration can lead to structural deterioration.

6.4. Durability Under Cyclic Wet-Dry and Freeze-Thaw Conditions

Styrene-butadiene rubber (SBR) has been shown to enhance self-healing under severe environmental conditions, such as cyclic wet-dry and freeze-thaw cycles. Vasilenko et al. (2021) observed a 50% recovery in compressive strength and reduced permeability, suggesting that SBR-modified concrete maintains structural integrity even in harsh climates.

6.5. Accelerated Aging and Chloride Ion Resistance

In accelerated aging tests, styrene-butadiene emulsion (SBR) has exhibited improved mechanical properties, reduced permeability, and enhanced crack closure. Daneti and Meddah (2013) reported a significant reduction in chloride ion ingress, which is essential for preventing corrosion in reinforced concrete structures. This suggests that SBR emulsions contribute to long-term durability and reduced maintenance costs.

The application of polymer-based self-healing systems significantly enhances concrete durability by sealing cracks, reducing permeability, and improving mechanical properties under various conditions. The selection of an appropriate polymer type is crucial to

optimizing self-healing efficiency based on the environmental exposure and structural requirements.

7.0 Conclusion

The incorporation of polymers into concrete has demonstrated significant advancements in enhancing mechanical properties, durability, and overall performance. By modifying the cement matrix and interfacial transition zone, polymers improve tensile strength, flexural strength, impact resistance, and permeability reduction. While some polymer types may slightly reduce compressive strength, their ability to enhance other crucial parameters makes polymer-modified concrete a superior alternative to conventional concrete in many applications.

One of the most remarkable innovations in polymer-modified concrete is its self-healing capability. Superabsorbent polymers (SAPs) and polymer-based healing agents effectively mitigate microcracking, thus extending the service life of concrete structures while reducing maintenance costs. Moreover, reduced permeability contributes to greater resistance against environmental degradation, particularly in aggressive conditions such as marine and industrial environments.

Quantitative analyses confirm that polymer incorporation leads to substantial improvements in key performance indicators, including increased tensile and flexural strength by up to 255% and a significant reduction in crack width over time. Additionally, polymer-modified concrete exhibits enhanced resistance to shrinkage and thermal expansion, contributing to long-term structural stability.

Despite the numerous advantages, optimizing polymer type and dosage is crucial to achieving a balance between workability, mechanical performance, and durability. Future research should focus on refining polymer formulations to maximize their benefits while minimizing potential drawbacks. The adoption of polymer-modified concrete aligns with global efforts toward sustainable construction, as it improves longevity, reduces resource consumption, and enhances structural resilience.

Polymer additives offer transformative improvements in concrete technology, making them essential for high-performance and sustainable construction applications. Continued advancements in material science will further optimize these modifications, ensuring the widespread adoption of polymer-enhanced concrete in the built environment.

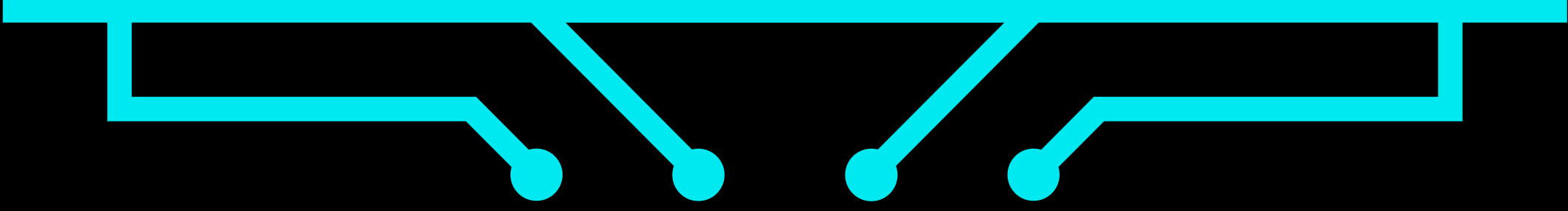
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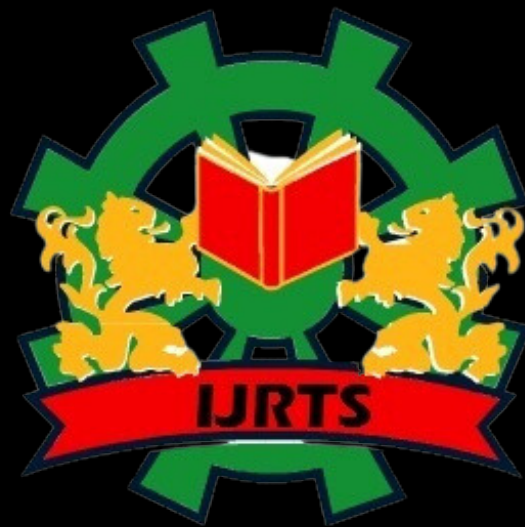
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